

# Intro to online learning: Learning from expert advice.

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# Review of optimization

$$\min_{\theta \in \mathcal{C}} f(\theta) \quad (1)$$

Here  $\mathcal{C}$  is a convex set and  $f(\cdot)$  is a convex function

We care about the complexity

$$f(\theta^k) - f(\theta^*) \leq \epsilon \quad (2)$$

If  $k = \log \frac{1}{\epsilon}$ , it is linear convergence

# Complexity Table

	convex	+ smooth	+ strong convex
gradient	$\frac{1}{\epsilon^2}$	$\frac{1}{\epsilon} \rightarrow \frac{1}{\sqrt{\epsilon}}$	$\frac{L}{m} \log \frac{1}{\epsilon} \rightarrow \sqrt{\frac{L}{m}} \log \frac{1}{\epsilon}$
SGD	$\frac{1}{\epsilon^2}$	$\frac{1}{\epsilon^2}$	$\frac{1}{m\epsilon}$
finite sum + SGD	$\frac{1}{\epsilon^2}$	$n + \frac{1}{\epsilon} \rightarrow \frac{1}{\sqrt{\epsilon}}$	$(n + \frac{L}{m}) \log \frac{1}{\epsilon} \rightarrow (n + \sqrt{\frac{L}{m}}) \log \frac{1}{\epsilon}$

Table: Complexity of first order methods

- Second order method: LBFGS, quasi-newton
- Non-convex optimization: have to convexify or adding noise to escape from local solutions
- How about DNN: too many local/global solutions!

- Problem statement for statistical learning: Given any dataset  $(x_1, y_1), \dots, (x_n, y_n)$  iid from  $\mathcal{D}$ . To find  $\mathcal{H} : X \rightarrow Y$
  - Reliable setting:  $\exists h^* \in \mathcal{H}$  s.t.  $\mathbf{P}(h^*(x) = y) = 1$
  - Error  $error(f) = \mathbb{E}\mathbf{1}(h(x) \neq y) \approx \frac{1}{n} \sum_{i=1}^n \mathbf{1}(h(x_i) \neq y_i)$
  - Online learning:
    - 1 set  $x_1$ , choose  $h_1 \in \mathcal{H}$ , such as  $\hat{y}_1 = h_1(x_1)$ ,
    - 2 set  $x_2$ , choose  $h_2 \in \mathcal{H}$ , such as  $\hat{y}_2 = h_2(x_2)$ , ...
- The cumulative loss  $\frac{1}{n} \sum_{i=1}^n \mathbf{1}(\hat{y}_i \neq y_i)$
- Design algorithm such that  $M(A) \leq \mathcal{O}(n)$
  - Example:  $X \in \{0, 1\}^d$ ,  $Y \in \{0, 1\}$ ,  $h = x(1)$  or  $x(4)$  or  $x(16)$

# Algorithm 1: Online ERM/FTL

- $V_1 = H$
- for  $t = 1, 2, \dots, n$ :
  - 1 receive  $x_t$ , pick any  $h \in V_t$
  - 2 predict  $\hat{y}_t = h(x_t)$
  - 3 Receive  $y_t$
  - 4 loss  $\mathbf{1}(\hat{y}_t \neq y_t)$  Update  $V_{t+1} = \{h \in V_t, h(x_t) = y_t\}$
- Convergence speed:  $1 \leq |V_t| \leq |H| - M_t$ , so  $M_t \leq |H| - 1$

## Algorithm 2: Majority voting (Halving)

- $V_1 = H$
- for  $t = 1, \dots, n$ 
  - 1 Receive  $x_t$
  - 2 Majority voting: for any  $h \in V_t$ ,  $\hat{y}_t = \arg \max \sum_h \mathbf{1}(h(x_t) = y_t)$
  - 3 Receive  $y_t$
  - 4 Update  $V_{t+1} = \{h \in V_t, h(x_t) = y_t\}$
- $1 \leq |V_t| \leq |H| \frac{1}{2}^{m_t}$ , so  $m_t \leq \log_2(|H|)$



If there does not exist a  $h$  such that  $h(x_i) = y_i, \forall i = 1, \dots, n$

$$\text{regret}(h) = \sum_{i=1}^n \mathbf{1}(y_t = h_t(x_t)) - \min_{h \in \mathcal{H}} \sum_{i=1}^n \mathbf{1}(y_t \neq h(x_t)) \quad (3)$$

## Example: stock prediction, Google

	Dearaj	Omid	Yuqing	Paul the Octopus	Truth
Day 1	Down	Up	Up	Down	Down
Day 2	Up	Up	Down	Down	Down
Day 3	Up	Down	Up	Up	Up
Day 4					

Table: Choices of expert

	Dearaj	Omid	Yuqing	Paul the Octopus
Day 1	1	1	1	1
Day 2	1	$\frac{1}{2}$	$\frac{1}{2}$	1
Day 3	$\frac{1}{2}$	$\frac{1}{4}$	$\frac{1}{2}$	1
Day 4	$\frac{1}{2}$	$\frac{1}{8}$	$\frac{1}{2}$	1

Table: Weights of expert

## Algorithm 3: Weighted majority

- $M$ : Number of mistakes
- $m$ : Number of mistakes of the best expert
- $n$ : Number of expert
- $W_t = \sum_{i=1}^n w_{it}$ ,  $W_1 = n$ , and  $W_{t+1} \leq W_t(1 - \frac{1}{4})$

$$\left(\frac{1}{2}\right)^m \leq W \leq n\left(\frac{3}{4}\right)^m \quad (4)$$

- $M \leq -\frac{\log 1/2}{\log 3/4}m + \frac{\log n}{\log 3/4}$

$$W_{i_{t+1}} = W_{it}(1 - \epsilon) \text{ if } \hat{y}_{it} \neq y_t \quad (5)$$

Then

$$(1 - \epsilon)^m \leq W \leq m(1 - \frac{1}{2}(1 - \epsilon))^m \quad (6)$$

$$m \log(1 - \epsilon) \leq \log m + m \log(\frac{1}{2} + \frac{1}{2}\epsilon) \quad (7)$$

$$M \leq \frac{-\log(1 - \epsilon)}{-\log(\frac{1}{2} + \frac{1}{2}\epsilon)} m + \frac{\log n}{-\log(\frac{1}{2} + \frac{1}{2}\epsilon)} \leq 2(1 + \epsilon)m + \mathcal{O}(\log m) \quad (8)$$

## Algorithm 4: randomized weighted majority (RWM)

- Set  $W_1^{(i)} = 1$  for all  $i$
- for  $t = 1, \dots, T$ ,

$$\text{Output} = \begin{cases} \text{Up} & \text{with probability } \frac{\sum_i W^i \mathbf{1}(y_t^i = \text{up})}{W} \\ \text{Down} & \text{Otherwise} \end{cases}$$

- $F_t = \frac{\sum_{i=1}^n W_t^i \mathbf{1}(\hat{y}_t^i \neq y^t)}{W_t}$ ,  $W_t = n(1 - \epsilon F_1) \dots (1 - \epsilon F_T)$
- $m \log(1 - \epsilon) \leq \log(W_{t+1}) \leq \log n + \sum_{i=1}^n \log(1 - \epsilon F_t)$   
 $\leq \log n - \epsilon \sum_{t=1}^T F_t = \log n - \mathbb{E}(M)$
- $\mathbb{E}(M) \leq \frac{\log n}{\epsilon} + \frac{-\log(1-\epsilon)}{\epsilon} m \approx (1 + \frac{\epsilon}{2})m + \frac{\log n}{\epsilon} \leq m + \sqrt{\frac{m \log n}{2}}$
- The last equality holds when  $\epsilon = \sqrt{\frac{2 \log n}{m}}$

# Relationship with convex optimization

- Learning with expert:  $\min \sum_i f_i(\theta_i)$
- $f_i(\theta_i) = \langle \theta_i, \ell \rangle = \mathbb{E}_{\theta_i}[l_i]$
- $C_i \sum \theta_i = 1, \theta_i \geq 0$