

Programs Synthesis from Polymorphic Refinement Types

Nadia Polikarpova

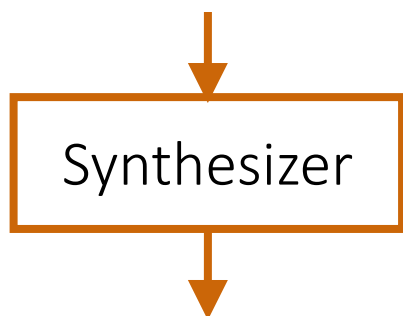
Ivan Kuraj

Armando Solar-Lezama



Program synthesis

“Make a list with n copies of x ”



```
replicate n x =  
  if n ≤ 0  
  then Nil  
  else Cons x  
    (replicate (dec n) x)
```

declarative
specification

$\Downarrow?$

2^{50}

executable
program

A stack of four orange rectangular boxes representing executable programs. A curly brace on the right side of the stack indicates the total number of programs, which is 2^{50} .

Modular verification for synthesis



Specifications for synthesis

refinement
types



Synthesizer

1. supports automatic, modular verification
2. abstract and concise
3. sufficiently expressive

```
replicate n x =  
  if n ≤ 0  
  then Nil  
  else Cons x (replicate (dec n) x)
```

Demo: replicate

-- Specification:

`replicate` :: $n: \text{Nat} \rightarrow x: \alpha \rightarrow \{v: \text{List } \alpha \mid \text{len } v = n\}$

`replicate` = ??

-- Components:

`zero` :: $\{v: \text{Int} \mid v = 0\}$

`inc` :: $x: \text{Int} \rightarrow \{v: \text{Int} \mid v = x + 1\}$

`dec` :: $x: \text{Int} \rightarrow \{v: \text{Int} \mid v = x - 1\}$

`leq` :: $x: \text{Int} \rightarrow y: \text{Int} \rightarrow \{\text{Bool} \mid v = (x \leq y)\}$

`neq` :: $x: \text{Int} \rightarrow y: \text{Int} \rightarrow \{\text{Bool} \mid v = (x \neq y)\}$

Synthesis from refinement types

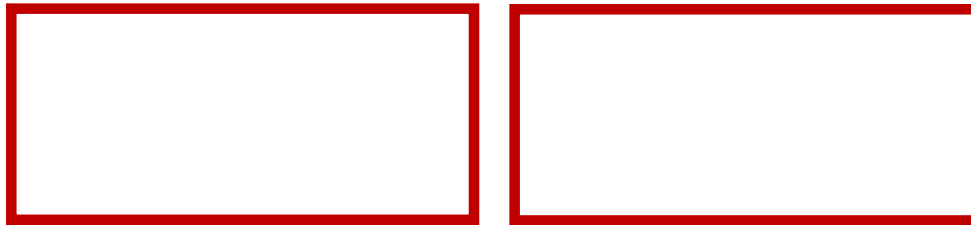
$\Gamma \vdash ?? :: T$

Synthesis from refinement types

$$\begin{array}{l} x_1 :: T_1; \dots \\ \phi_1; \dots \end{array} \vdash \text{??} :: T$$

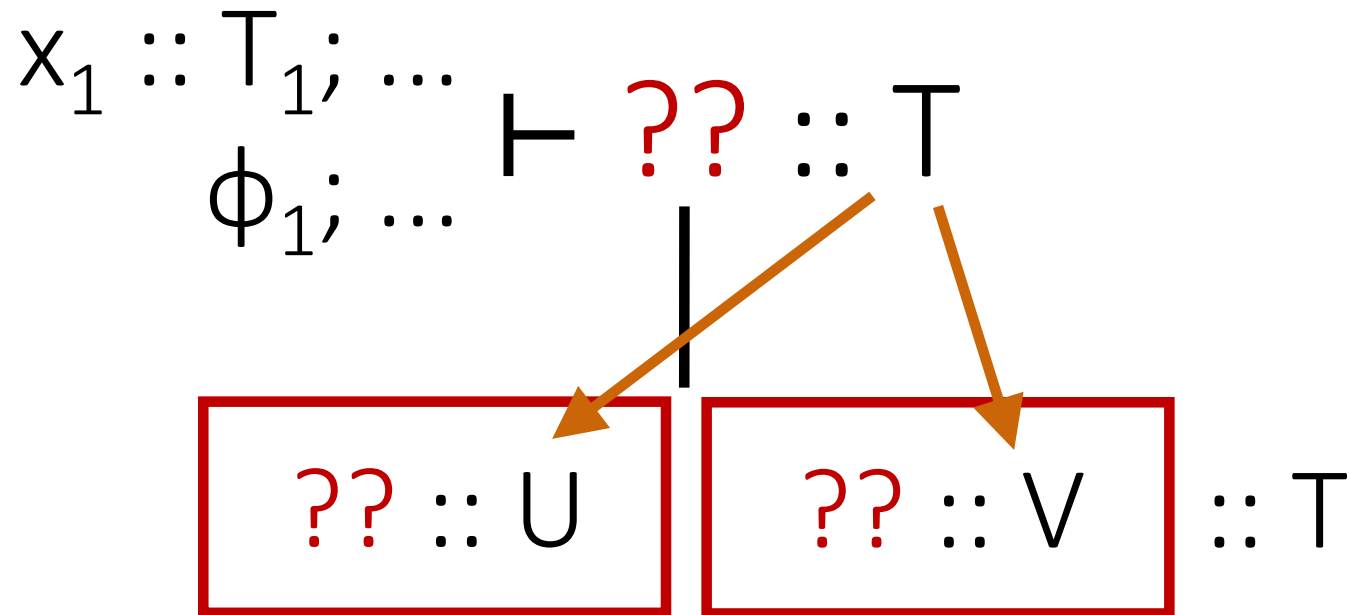
Synthesis from refinement types

$$x_1 :: T_1; \dots \vdash \text{??} :: T$$
$$\phi_1; \dots$$



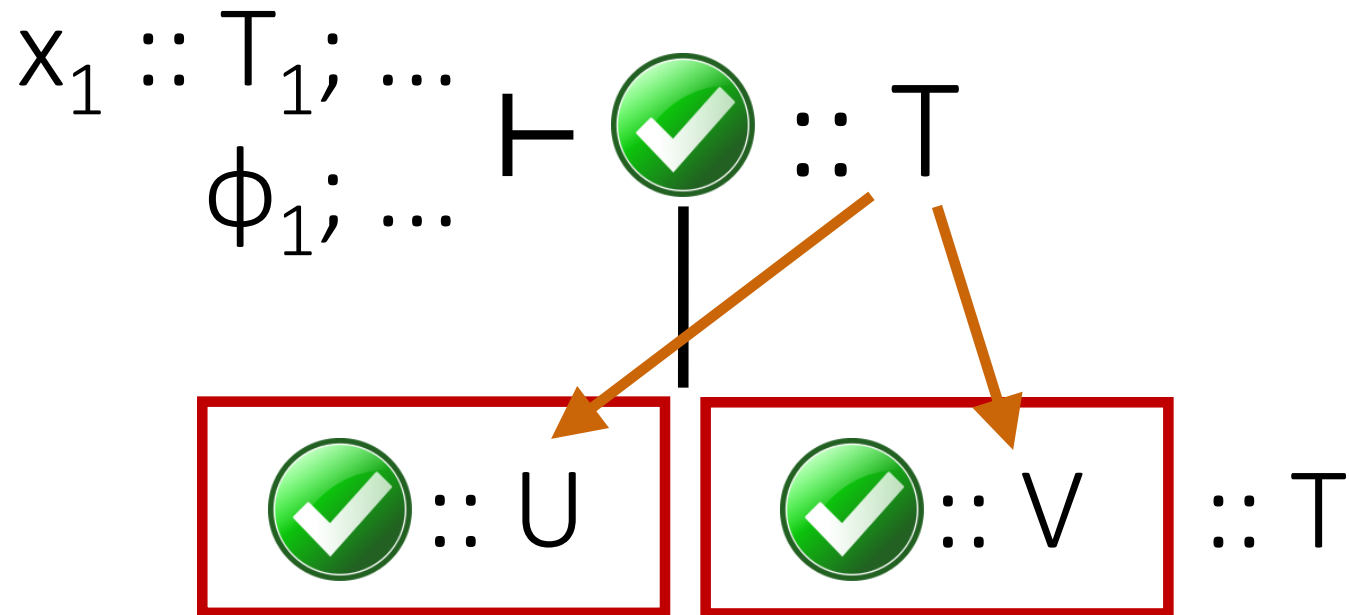
I. top-down enumerative search

Synthesis from refinement types



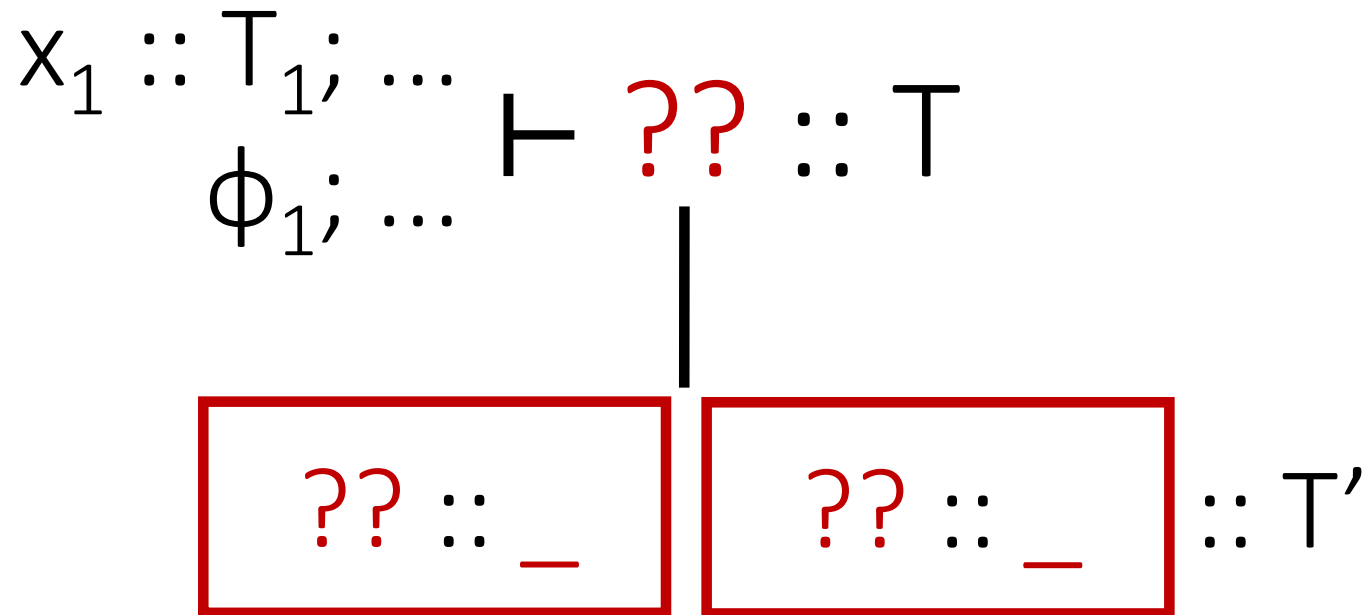
I. top-down enumerative search

Synthesis from refinement types



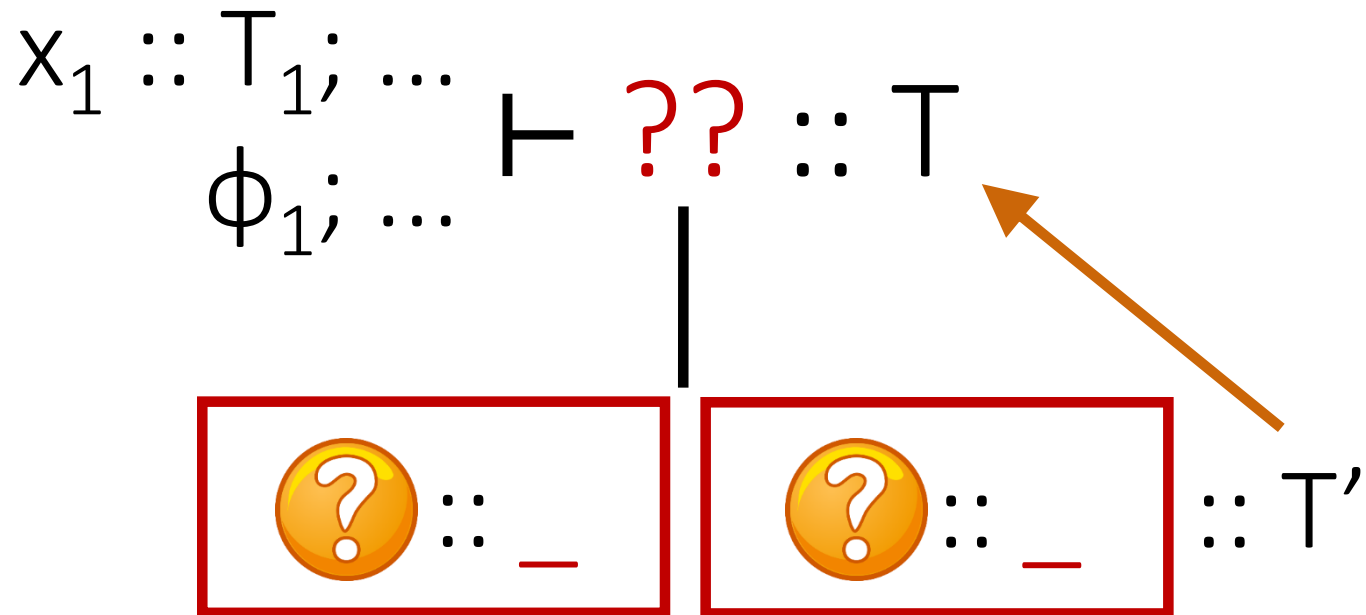
I. top-down enumerative search

Synthesis from refinement types



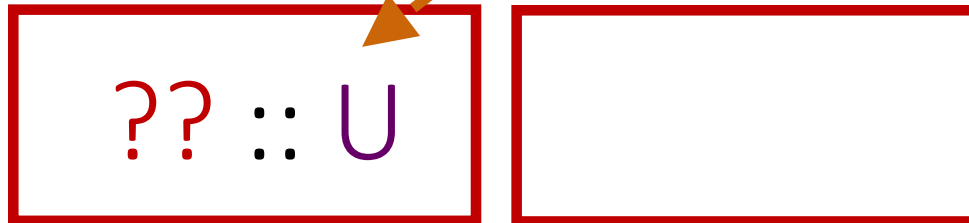
I. top-down enumerative search

Synthesis from refinement types



I. top-down enumerative search

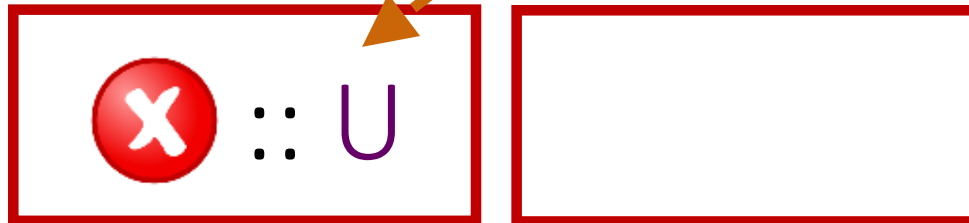
Synthesis from refinement types

$$x_1 :: T_1; \dots \quad \vdash \quad ?? :: T$$
$$\phi_1; \dots$$


I. top-down enumerative search

II. round-trip type checking

Synthesis from refinement types

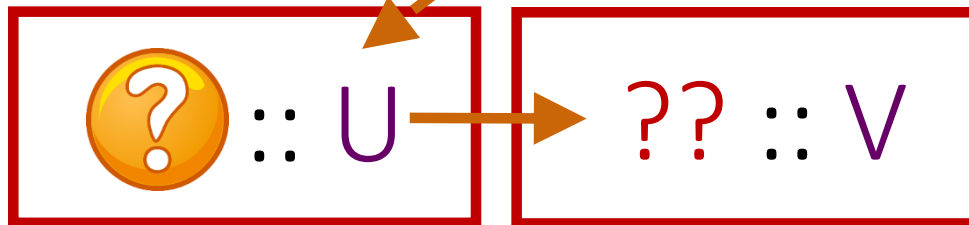
$$x_1 :: T_1; \dots \quad \vdash \quad ?? \quad :: T$$
$$\phi_1; \dots$$


I. top-down enumerative search

II. round-trip type checking

Synthesis from refinement types

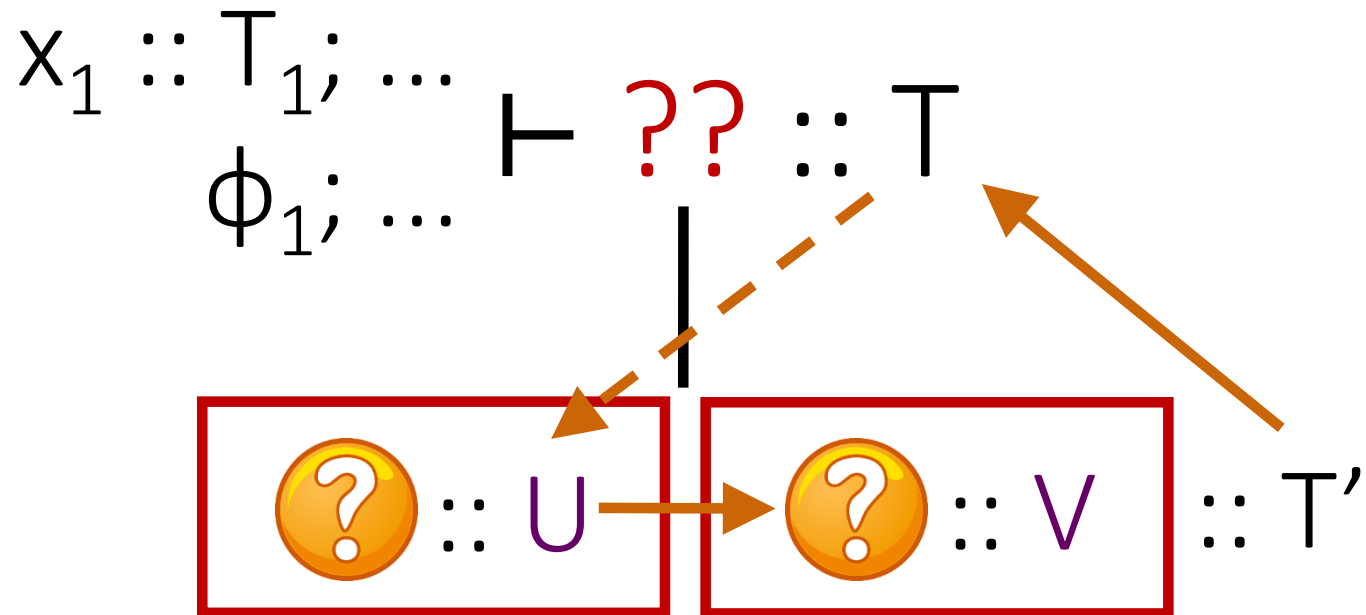
$x_1 :: T_1; \dots$
 $\phi_1; \dots \vdash \text{??} :: T$



I. top-down enumerative search

II. round-trip type checking

Synthesis from refinement types



I. top-down enumerative search

II. round-trip type checking

Synthesis from refinement types

$$x_1 :: T_1; \dots \vdash \text{??} :: T$$
$$\phi_1; \dots$$

if $\text{??} :: \text{Bool}$ then \square else \square

I. top-down enumerative search

II. round-trip type checking

Synthesis from refinement types

$$x_1 :: T_1; \dots \vdash \text{??} :: T$$
$$\phi_1; \dots$$

if then $P \vdash \text{??} :: T$ else

I. top-down enumerative search

II. round-trip type checking

III. condition abduction

Synthesis from refinement types

$$\begin{array}{l} x_1 :: T_1; \dots \\ \vdash \text{??} :: T \\ \phi_1; \dots \end{array}$$

if $\text{??} :: \{\text{Bool} \mid v=P\}$ then $P \vdash \checkmark :: T$ else $\neg P \vdash \text{??} :: T$

I. top-down enumerative search

II. round-trip type checking

III. condition abduction

Round-trip type checking

$\Gamma \vdash ?? \text{ :: } \{\text{List Neg} \mid \text{len } v \geq 5\}$

Round-trip type checking

Nil ; 0 ; 5 ; -5
zeros
replicate ; Cons ⊢ ?? :: {List Neg | len v ≥ 5}

Round-trip type checking

Nil :: {List a | len v = 0} ; 0 ; 5 ; -5

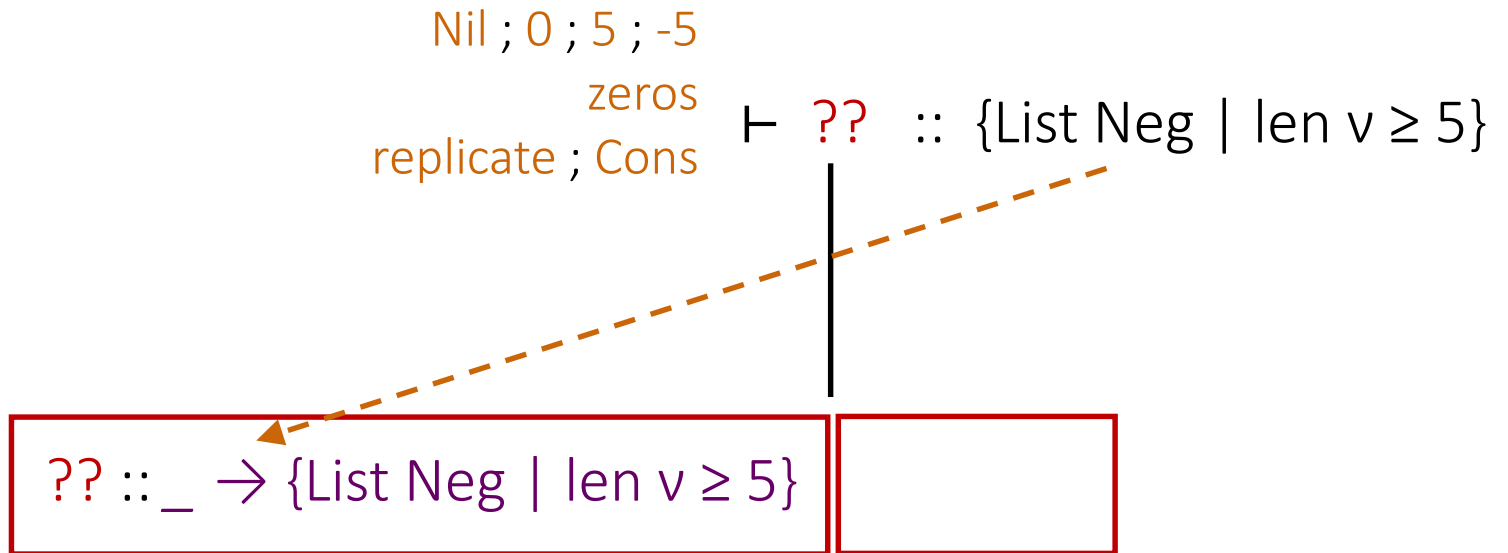
zeros
replicate ; Cons

⊢  :: {List Neg | len v ≥ 5}

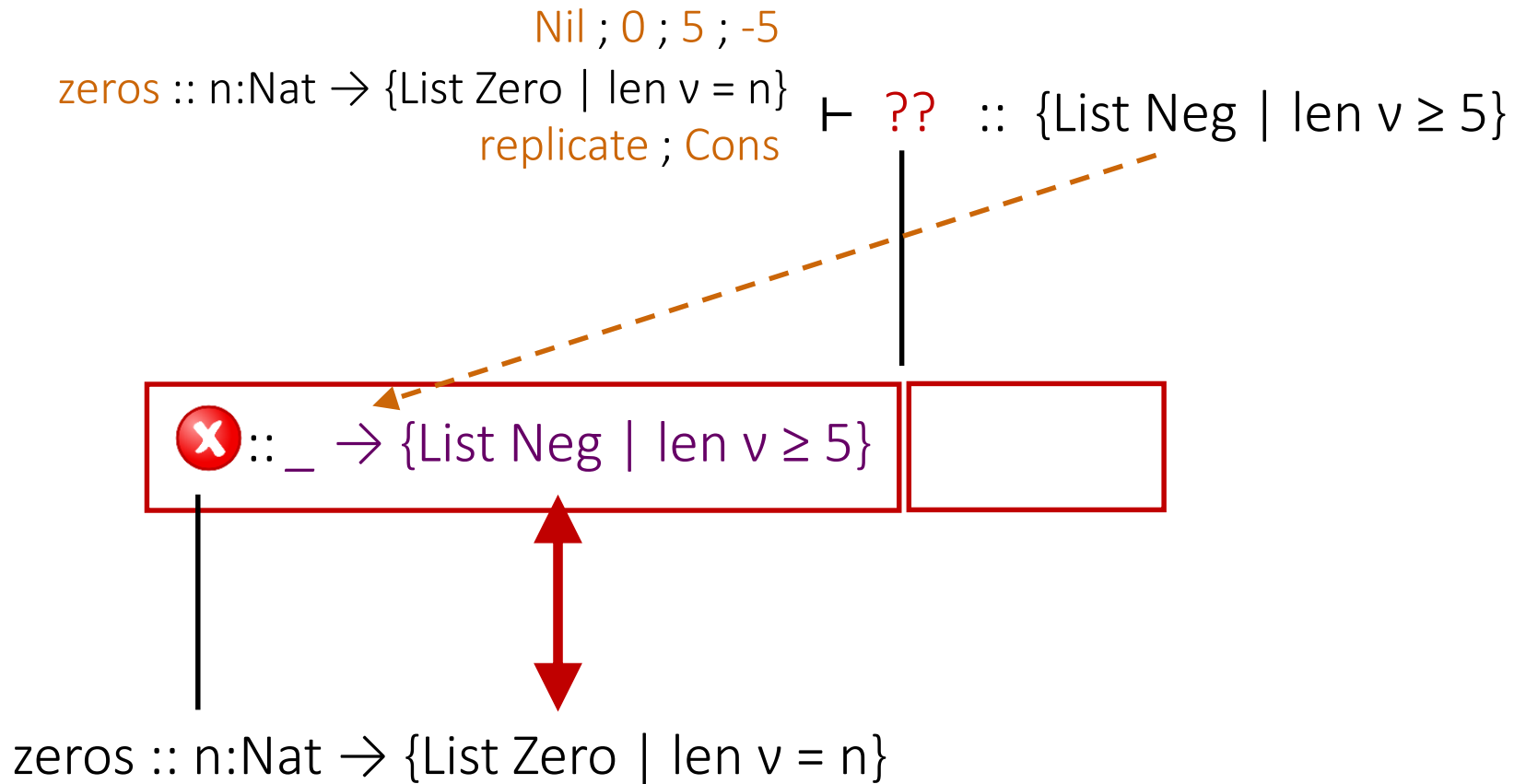
Nil :: {List Neg | len v = 0}



Round-trip type checking



Round-trip type checking



Round-trip type checking

Nil ; 0 ; 5 ; -5

zeros
replicate

Cons

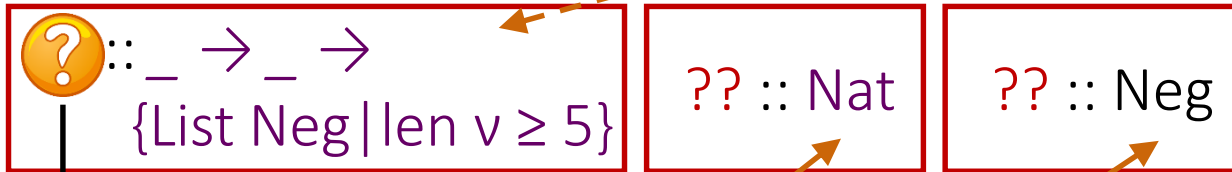
$\vdash ?? \ :: \{\text{List Neg} \mid \text{len } v \geq 5\}$

$?? \ :: _ \rightarrow _ \rightarrow$
 $\{\text{List Neg} \mid \text{len } v \geq 5\}$

Round-trip type checking

$\text{replicate} :: n: \text{Nat} \rightarrow x: a \rightarrow \{\text{List } a \mid \text{len } v = n\}$
 $\vdash \text{??} :: \{\text{List Neg} \mid \text{len } v \geq 5\}$

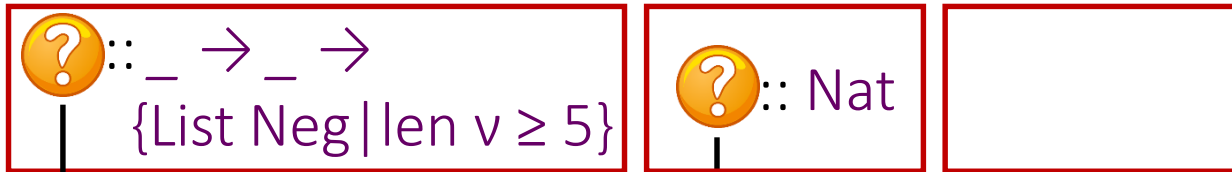
$\text{Nil}; 0; 5; -5$
 zeros
 Cons



$\text{replicate} :: n: \text{Nat} \rightarrow x: \text{Neg} \rightarrow \{\text{List Neg} \mid \text{len } v = n\}$

Round-trip type checking

$\text{replicate} :: n: \text{Nat} \rightarrow x: a \rightarrow \{\text{List } a \mid \text{len } v = n\}$
 \vdash
 $\text{Nil} ; 0 ; 5 ; -5$
 zeros
 Cons
 $??$
 $:: \{\text{List Neg} \mid \text{len } v \geq 5\}$



$0 :: \{v = 0\}$

$\text{replicate} :: n: \text{Nat} \rightarrow x: \text{Neg} \rightarrow \{\text{List Neg} \mid \text{len } v = n\}$

Round-trip type checking


$\text{replicate} :: n: \text{Nat} \rightarrow x: a \rightarrow \{\text{List } a \mid \text{len } v = n\}$
 $\vdash \text{Nil} ; 0 ; 5 ; -5$
 $\vdash \text{zeros} \text{ } ?? :: \{\text{List Neg} \mid \text{len } v \geq 5\}$
 Cons



$0 :: \{v = 0\}$

$\text{replicate} :: n: \text{Nat} \rightarrow x: \text{Neg} \rightarrow \{\text{List Neg} \mid \text{len } v = n\}$

Round-trip type checking

$\text{replicate} :: n: \text{Nat} \rightarrow x: a \rightarrow \{\text{List } a \mid \text{len } v = n\}$
 \vdash

 $:: \{\text{List } \text{Neg} \mid \text{len } v \geq 5\}$


$\text{Nil}; 0; 5; -5$
 zeros
 Cons



$5 :: \{v = 5\}$

$\text{replicate} :: n: \text{Nat} \rightarrow x: \text{Neg} \rightarrow \{\text{List } \text{Neg} \mid \text{len } v = n\}$

Round-trip type checking

$\text{replicate} :: n: \text{Nat} \rightarrow x: a \rightarrow \{\text{List } a \mid \text{len } v = n\}$
 \vdash

 $:: \{\text{List } \text{Neg} \mid \text{len } v \geq 5\}$

$\text{Nil} ; 0 ; 5 ; -5$
 zeros
 Cons



$5 :: \{v = 5\}$ $-5 :: \{v = -5\}$

$\text{replicate} :: n: \text{Nat} \rightarrow x: \text{Neg} \rightarrow \{\text{List } \text{Neg} \mid \text{len } v = n\}$

Condition abduction

$\text{Nil} ; 0 ; -5 ; n :: \text{Nat}$
 $(\leq) ; (\neq) \vdash \text{??} :: \{\text{List Neg} \mid \text{len } v = n\}$
 P

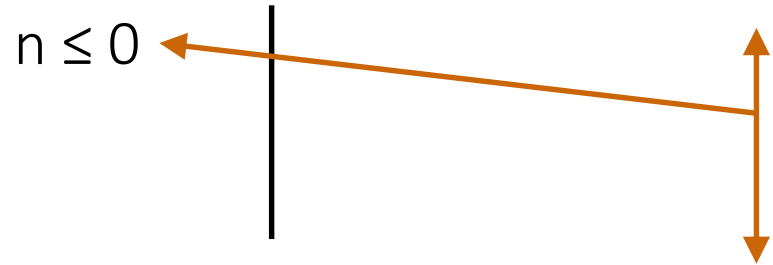
Condition abduction

$\text{Nil} ; 0 ; -5 ; n :: \text{Nat}$

$(\leq) ; (\neq) \vdash ?? :: \{\text{List Neg} \mid \text{len } v = n\}$

$n \leq 0$

$\text{Nil} :: \{\text{List Neg} \mid \text{len } v = 0\}$



Condition abduction

$\text{Nil} ; 0 ; -5 ; n :: \text{Nat}$
 $(\leq) ; (\neq) \vdash \checkmark :: \{\text{List Neg} \mid \text{len } v = n\}$
 $n \leq 0$

if $n \leq 0$ then Nil else $\Gamma ; \neg(n \leq 0) \vdash ?? :: \{\text{List Neg} \mid \text{len } v = n\}$

Liquid abduction

$$n \geq 0 \wedge \text{len } v = 0 \wedge P \Rightarrow \text{len } v = n$$

$n :: \text{Nat}$

$\text{Nil} :: \{\text{List } a \mid \text{len } v = 0\}$

Liquid abduction

$n \geq 0 \wedge \text{len } v = 0 \wedge P \wedge \neg(\text{len } v = n)$

|

★ ≤ ★

★ ≠ ★

Liquid abduction

$$n \geq 0 \wedge \text{len } v = 0 \wedge P \Rightarrow \text{len } v = n$$

|

$$n \leq 0$$

$$n \leq -5$$

$$-5 \leq n$$

$$n \neq 0$$

$$n \neq -5$$

Liquid abduction

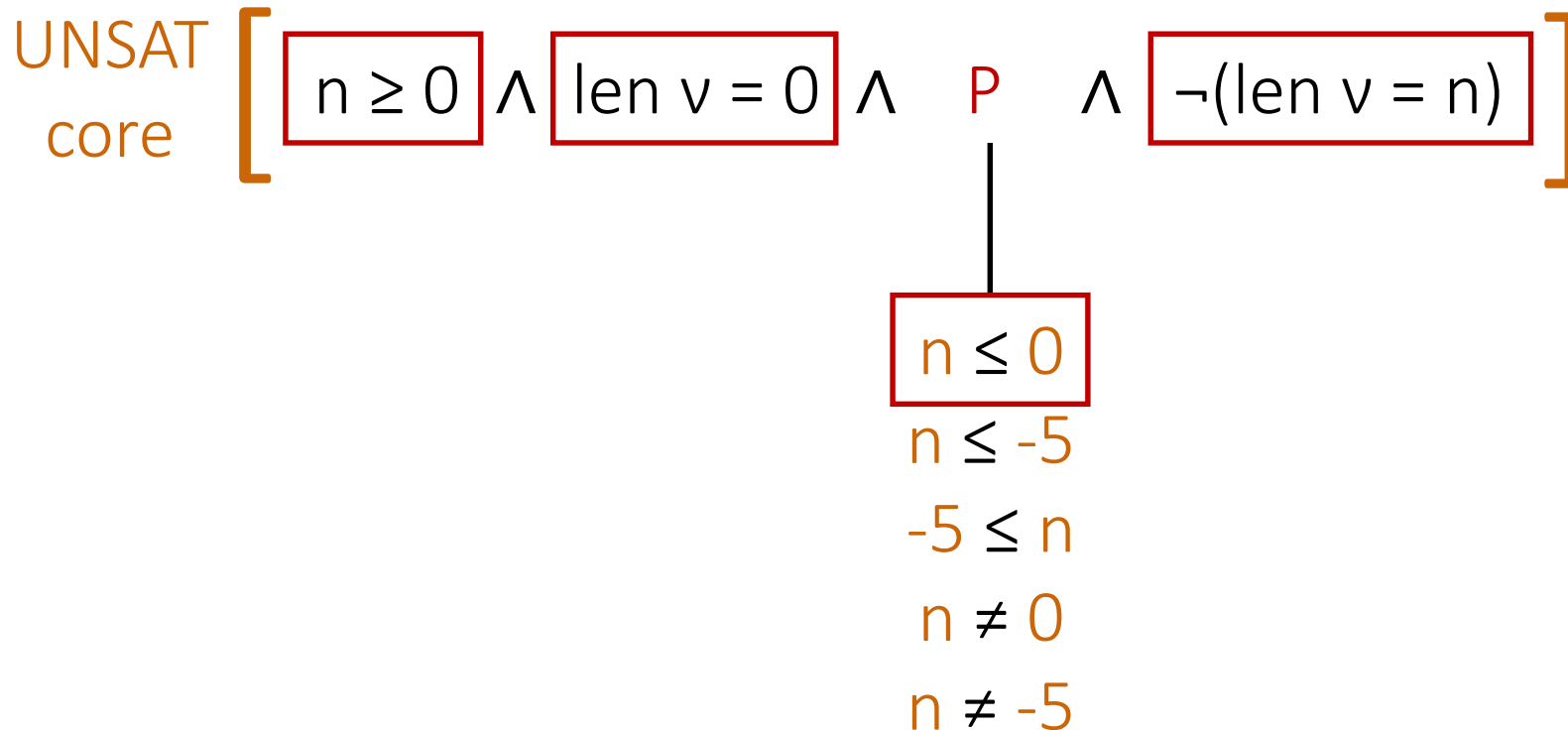
UNSAT core

$$\left[n \geq 0 \wedge \text{len } v = 0 \wedge P \wedge \neg(\text{len } v = n) \right]$$

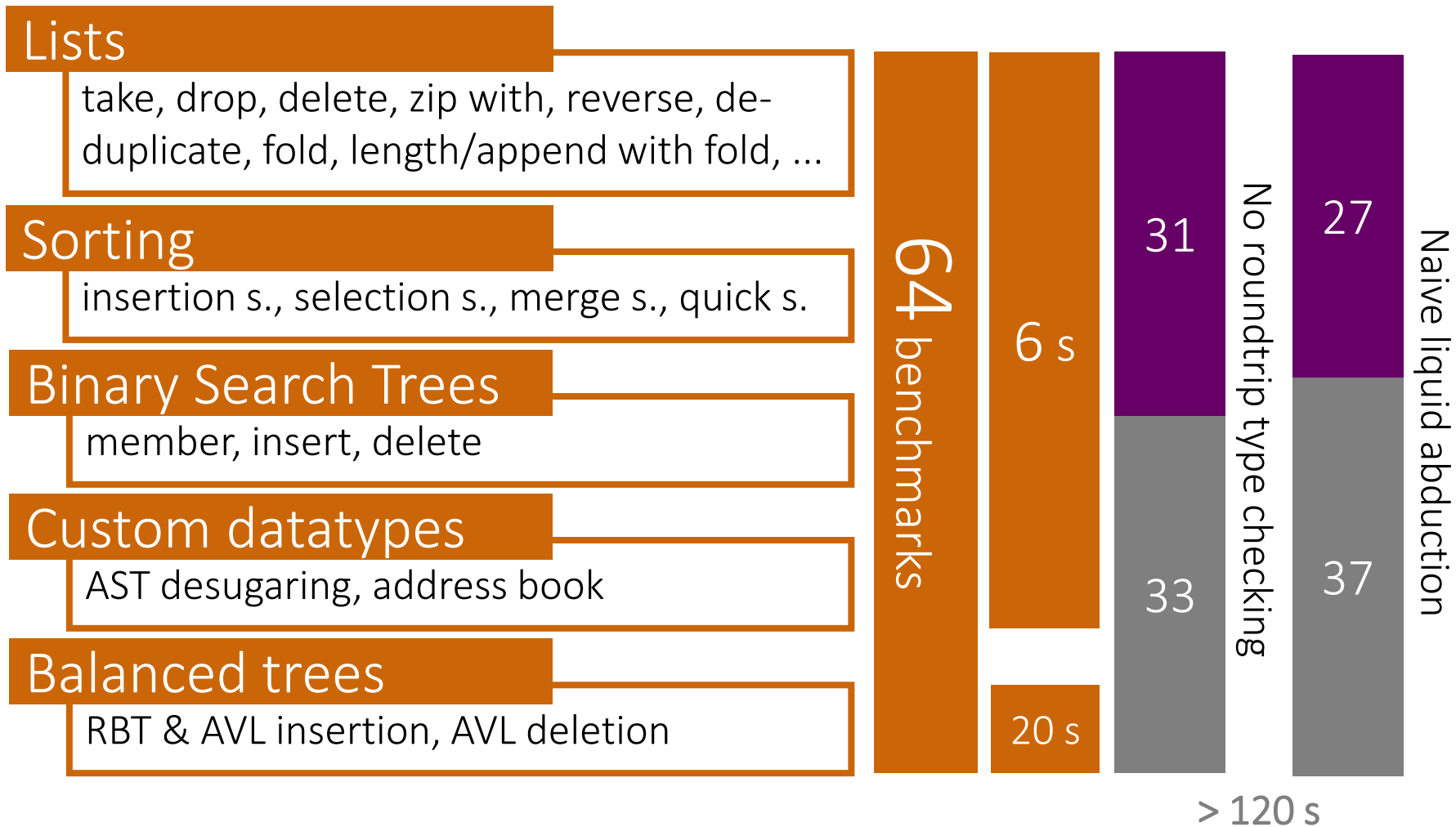
|

$n \leq 0$
 $n \leq -5$
 $-5 \leq n$
 $n \neq 0$
 $n \neq -5$

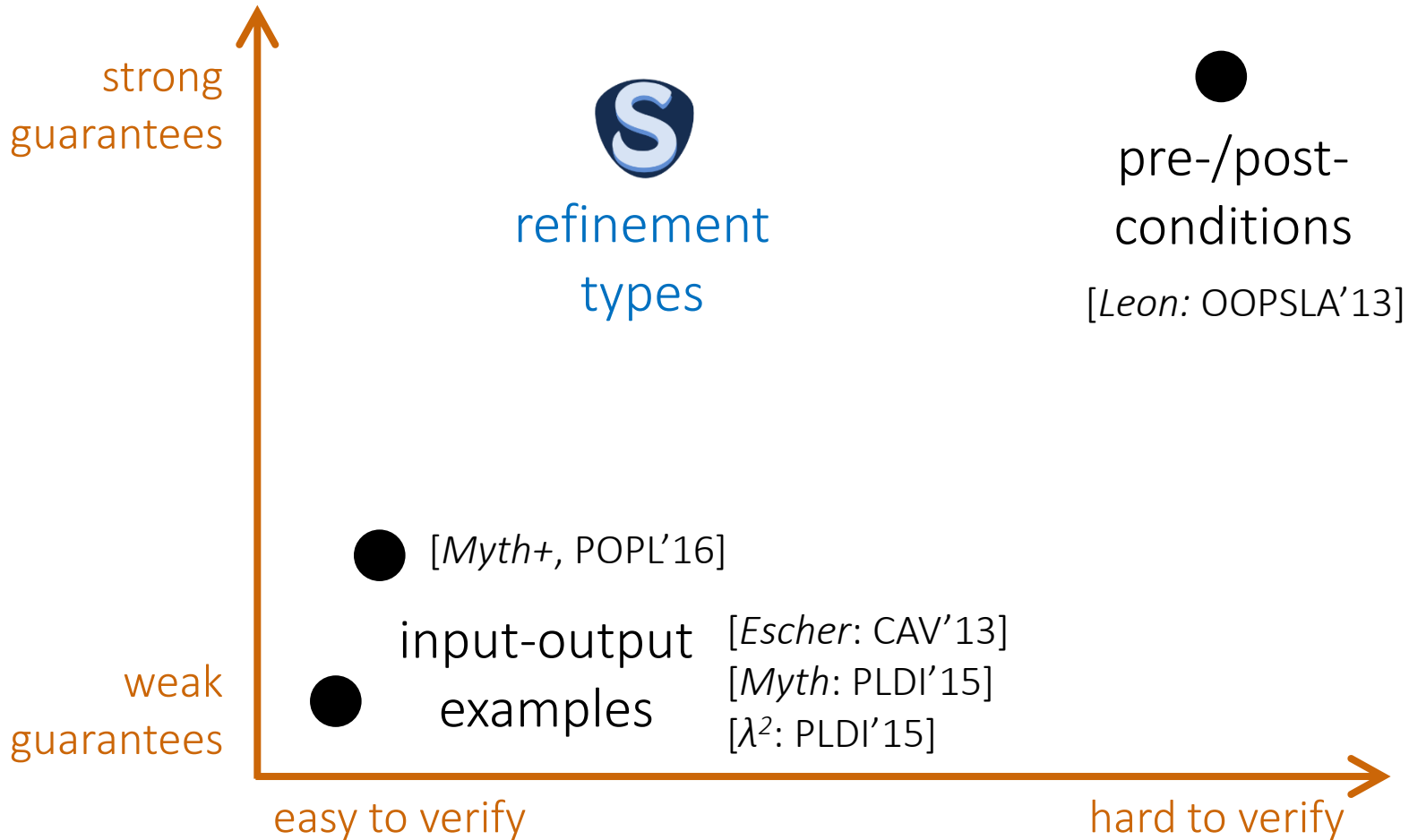
Liquid abduction



Evaluation



Synthesis of recursive programs





<http://tiny.cc/synquid>