

On the Computational Power of Neural Nets

by Hava Siegelmann and Eduardo Sontag (1995)[3]

Presented by Jiaman Wu

1

Outline

- Problem Statement
- Model with I/O
- Results of Model with I/O
- Model without I/O
- Results of Model without I/O
- Model of 2-stack Turing Machine
- Highlight of Proof

2

Problem Statement

- Recurrent Neural Net (RNN) can theoretically simulate any Turing Machines (TM) in real time.
 - Turing complete: compute any algorithm which our computer can compute
- Many consequences of Turing completeness:
 - Undecidable \Leftrightarrow not Turing computable: no algorithm can compute it and TM loops forever.
 - Undecidable to determine if a given neuron ever assumes the value "0".
 - Undecidable to determine if a dynamical system of the particular form $x(t+1)=\sigma(Ax(t)+c)$ ever reaches an equilibrium point, from a given initial state. [2]

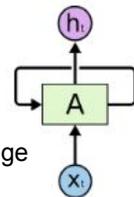


Figure 1. RNN diagram. Image from [4]

3

Model of RNN with I/O

- Basic structure of RNN
 - **Saturated-linear function** is as activation of finite neurons in the **hidden** layer.
 - **2 infinite binary** input lines are in **input** layer (TBD in the later).
 - **Partially defined function** maps to **output** (TBD in the later).
 - Coefficients are **rationals**.

$$x(t + 1) = \sigma(Ax(t) + bu(t) + c): \text{ update rule}$$

$$x(t) \in \mathbb{Q}^N: \text{ the state at time } t$$

$$u(t) = [u_D(t), u_V(t)]^T \in \{0, 1\}^2: \text{ inputs (TBD in the later)}$$

$$A \in \mathbb{Q}^{N \times N}: \text{ weights}$$

$$b \in \mathbb{Q}^{N \times 2}: \text{ weights}$$

$$c \in \mathbb{Q}^N: \text{ bias}$$

$$y(t) = [y_D(t), y_V(t)]^T \in \{0, 1\}^2: \text{ outputs (TBD in the later)}$$

$$\sigma(x) := \begin{cases} 0 & \text{if } x < 0 \\ x & \text{if } 0 \leq x \leq 1 \\ 1 & \text{if } x > 1. \end{cases}$$

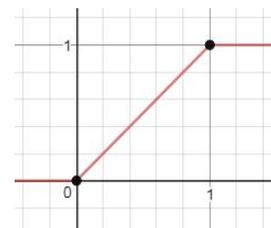


Figure 2. Saturated-linear function graph

4

Model of RNN with I/O

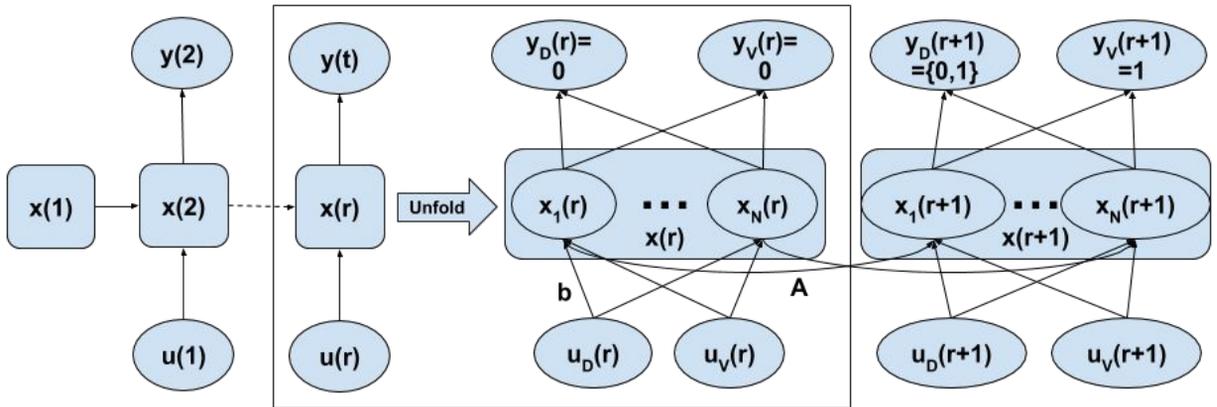


Figure 3. Diagram of RNN with I/O

Model of RNN with I/O

- Two infinite binary input lines for reading $w \in \{0, 1\}^k$

$$u(t) = \begin{cases} [w_t \in \{0, 1\}, 1]^T, & \text{if } t = 1, \dots, k \\ [0, 0]^T, & \text{otherwise} \end{cases}$$

t	1	2	3	4	5	...
data line u_D	$w_1=1$	$w_2=0$	$w_3=1$	0	0	...
valid line u_V	1	1	1	0	0	...
$u(t)$	$u(1)=[1, 1]^T$	$u(2)=[0, 1]^T$	$u(3)=[1, 1]^T$	$u(4)=[0, 0]^T$	$u(5)=[0, 0]^T$...

Model of RNN with I/O

- Partially defined function maps input w to output $\phi(w) \in \{0,1\}^l$ or $\phi(w)$ undefined
 - $\phi(w)$ is **partially defined function**: some w 's are not mapped and their $\phi(w)$ is undefined.
 - $\phi(w)$ is undefined \Leftrightarrow TM loops forever and RNN gives all 0's on both lines.
 - $\phi(w)$ is defined \Leftrightarrow TM halts and RNN gives output on data line and 1's on validation line
 - Output **takes time** $r \in \mathbb{N}$, the response time.

$$y_r(t) = \begin{cases} [\phi_{t-r+1}(w), 1]^T, & \text{if } \phi(w) \text{ defined and } t = r, \dots, r + l - 1 \\ [0, 0]^T, & \text{otherwise} \end{cases}$$

7

Results of Model with I/O

- Theorem 1. Let $\phi: \{0,1\}^+ \rightarrow \{0,1\}^+$ be any recursively computable partial function. Then, there exists a net N with the following property: If N is started from the 0 initial state, and the input sequence u is applied, N produces the output y_r for some r . Furthermore, if a p -stack Turing Machine M computes $\phi(w)$ in time $T(w)$, then one may take $r(w) = O(|w|) + T(w)$.
- **Any function computable by a Turing Machine can be computed by such a processor net.**
- **886 neurons** are sufficient for computing such a universal partial recursive function to be Turing complete.

8

Model of RNN without I/O

- A model analogous to the previous one but without input and output stream.
- It is more convenient to encode the **binary input string into Cantor set** and to the initial state.
 - The number ranges in $[0,1)$ with gaps.
 - Encoding converts binary into rationals.

$$\delta(w) = \sum_{i=1}^k \frac{2w_i + 1}{4^i}$$

$$x^1(w) = [\delta(w), 0, \dots, 0]^T \in \mathbb{Q}^N$$

$$x^{t+1}(w) = \sigma(Ax^t(w) + c)$$

Example of encoding with Cantor set:

empty $\rightarrow 0$

0 $\rightarrow 1/4$

1 $\rightarrow 3/4$

01 $\rightarrow 1/4 + 3/4^2 = 0.4375$

10 $\rightarrow 3/4 + 1/4^2 = 0.8125$

11 $\rightarrow 3/4 + 3/4^2 = 0.9375$



Figure 4. Graph of Cantor set [1]

9

Results of RNN without I/O

Theorem 2. Let M be a p -stack Turing Machine computing $\phi: \{0,1\}^+ \rightarrow \{0,1\}^+$ in time T . Then there exists a processor net N without inputs so that properties below hold. In both cases, for each $w \in \{0,1\}^+$, we consider the corresponding **initial state** $x^1(w) = [\delta[w], 0, \dots, 0]^T \in \mathbb{Q}^N$.

$\phi(w)$ is **undefined**, if and only if $x_2^j(w)$ is identically equal to 0 , for all j . If instead $\phi(w)$ is **defined** and computed in time T , there exists a $J=O(T)$ so that: $x_2^j(w)=0$, $j \in [0, J-1]$ and $x_2^J(w)=1$, $x_1^j(w)=\delta[\phi(w)]$. ($x^j(w)=[\delta[\phi(w)], 1, \dots]^T$)

Model of 2-stack TM

- 2-tape TM
 - 2 tapes
 - Operations: read, write, move
- 2-stack TM is a common 2-tape TM, but has different operations.
 - Read the top element
 - Push 0 or 1
 - Pop
 - Judge if empty

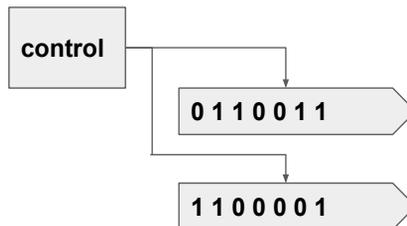


Figure 5. Diagram of 2-tape TM

11

Highlights of Proof

- Rational weighted RNN is proven to be computationally equivalent to TMs
- Encode binary string into Cantor set of number $\delta(w) = \sum_{i=1}^k \frac{2w_i+1}{4^i}$
- Simulate TM stack operation:
 - No operations: $q = q$
 - Read the top element: $\text{top}(q) = \sigma(4q-2)$
 - e.g. 1000 $\rightarrow 0.832$, $\sigma(4*0.832-2= 1.328) = 1$
 - e.g. 0001 $\rightarrow 0.3398$, $\sigma(4*0.3398-2= -0.6408) = 0$
 - Push 0: $q/4+1/4$; Push 1: $q/4+3/4$
 - Pop: $q = 4q-2\text{top}(q)-1 = 4q-1$ (if pop 0) or $4q-3$ (if pop 1)
 - e.g. pop 1000 $\rightarrow 4*0.832-3=0.328$; 000 $\rightarrow 0.328$
 - Judge if empty: $\sigma(4q)=0 \Leftrightarrow \text{empty}$
- 3 neurons for encoding q , $\text{top}(q)$, empty indicator per stack + finite neurons for no operators, push, pop + finite neurons for reading and computing the next step.

12

Reference

[1] A little finding regarding the middle-half Cantor set. Retrieved from <https://drchristiansalas.files.wordpress.com/2013/02/n04.pdf>

[2] Hopfield, J. J., & Tank, D. W. (1985). "Neural" computation of decisions in optimization problems. *Biological cybernetics*, 52(3), 141-152.

[3] Siegelmann, H. T., & Sontag, E. D. (1995). On the computational power of neural nets. *Journal of computer and system sciences*, 50(1), 132-150.

[4] Wang, T. Recurrent Neural Network. Retrieved from https://www.cs.toronto.edu/~tingwuwang/rnn_tutorial.pdf