Last Time

- Graph definition
 - Vertices connected by edges
- Types of graphs:
 - Multigraph (multiple edges between same vertices)
 - Pseudograph (allows self loops)
 - Simple graph (neither of the above)
 - Hypergraph (edges between more than two vertices)
 - Directed graph (edges have orientation)

Question: Graph Generalizations

Which graph types are exhibited by the graph below?



- A) Multigraph
- B) Pseudograph
- C) Simple Graph
- D) Hypergraph
- E) Directed Graph

Today

- Graph Terminology
- Handshake Lemma
- Walks, paths and Cycles

Graph Terminology I

- Two vertices u and v are *adjacent* if there is an edge connecting them.
- A vertex v is *incident* on an edge e (or is an *endpoint* of e) if v is one of the vertices e connects.



u is adjacent to v v is non-adjacent to w u is incident on e w is not incident on e

Graph Terminology II

- The *neighborhood* of a vertex v (denoted N(v)) is the set of vertices adjacent to v along with v.
- The *degree* of v (denoted d(v)) is the number of vertices adjacent to v.



Graph Terminology III

A graph is *d-regular* if all vertices have degree
d. It is *regular* if it is d-regular for some d.

This graph is 3-regular



This graph is not regular



Question: Degrees

Which vertex in this graph has the smallest degree?



The Handshake Lemma

(Theorem 1.1) For any graph G = (V, E),

$$\sum_{v \in V} d(v) = 2|E|.$$



$2 + 3 + 1 + 2 = 8 = 2 \cdot 4$

Proof I

Strategy: Counting things in two different ways.

Show both sides are equal to the number of pairs of (v,e) where v is a vertex incident on an edge e.

Proof II

Right Hand Side:

Each edges e = (u,v) has two incident vertices, u and v.



Total number of pairs is 2|E|.

Proof III

Left Hand Side:

Each vertex v is incident on d(v) edges.



Total number of pairs is $\sum_{v \in V} d(v)$.

Proof IV

Equating the two sides we find:

$$\sum_{v \in V} d(v) = \#\{\text{Incidence pairs}\} = 2|E|.$$

QED.

Question

How many edges does a 3-regular graph with 5 vertices have?

- A) 3
- B) 6
- C) 7.5
- D) 10
- E) There is no such graph

Examples of Graphs I

A complete graph on n vertices (denoted K_n) is a graph with n vertices and an edge between every pair of them



Examples of Graphs II

A cycle on n vertices (denoted C_n) is a graph with n vertices connected in a loop.



A *path* on n vertices (denoted P_n) is a graph with n vertices connected in a chain.



Examples of Graphs III

A graph H is a *subgraph* of G if V(H) \subset V(G) and E(H) \subset E(G).



A subgraph H is an *induced subgraph* if it contains *all* the edges of G connecting two vertices in V(H).

Examples of Graphs IV

A *bipartite graph* is a graph whose vertices can be split into two parts where all edges connect one part to the other.



A *complete bipartite graph* (denoted K_{n,m}) has an edge connecting every element of one part (of size n) to every element of the other (of size m).

Question: Cycle Identification

Which of the graphs below are cycles?



Question: Edge Counts

Which of these graphs has the greatest number of edges?

- A) C_{10} (10 edges)
- B) P_{12} (11 edges)
- C) K_5 (10 edges)
- D) K_{3,4} (12 edges)