

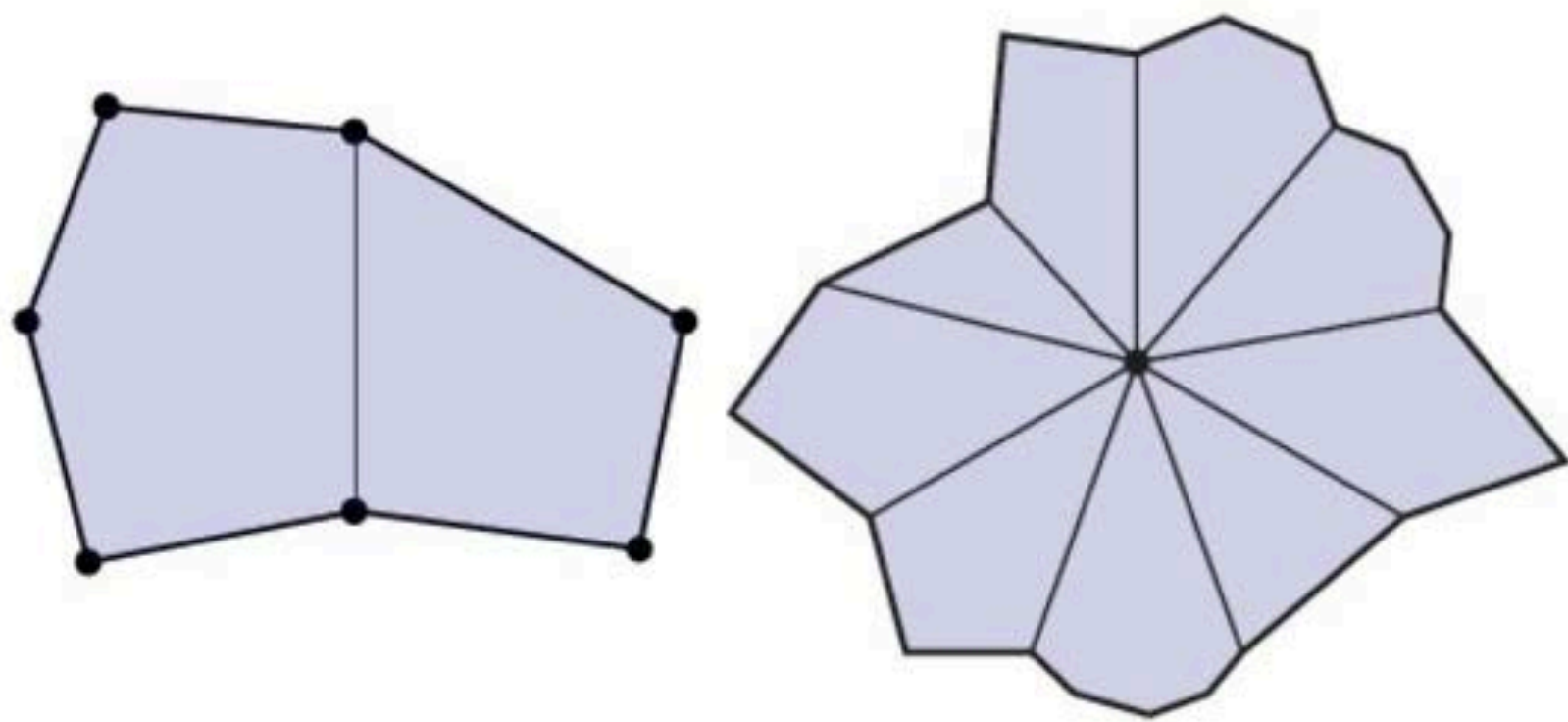
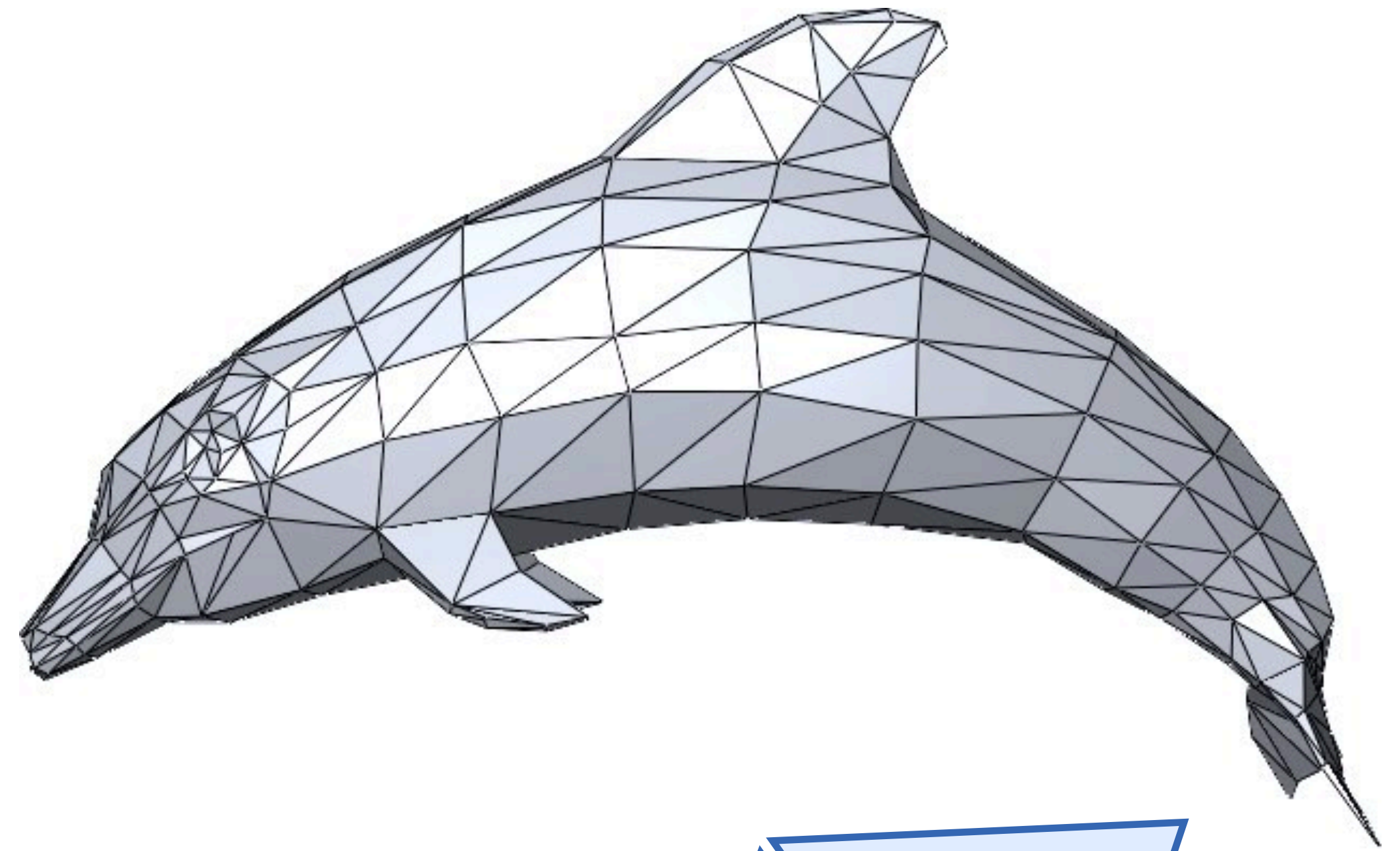
**CSE 167 (FA22)**  
**Computer Graphics:**  
**Digital Geometry Processing**

**Albert Chern**

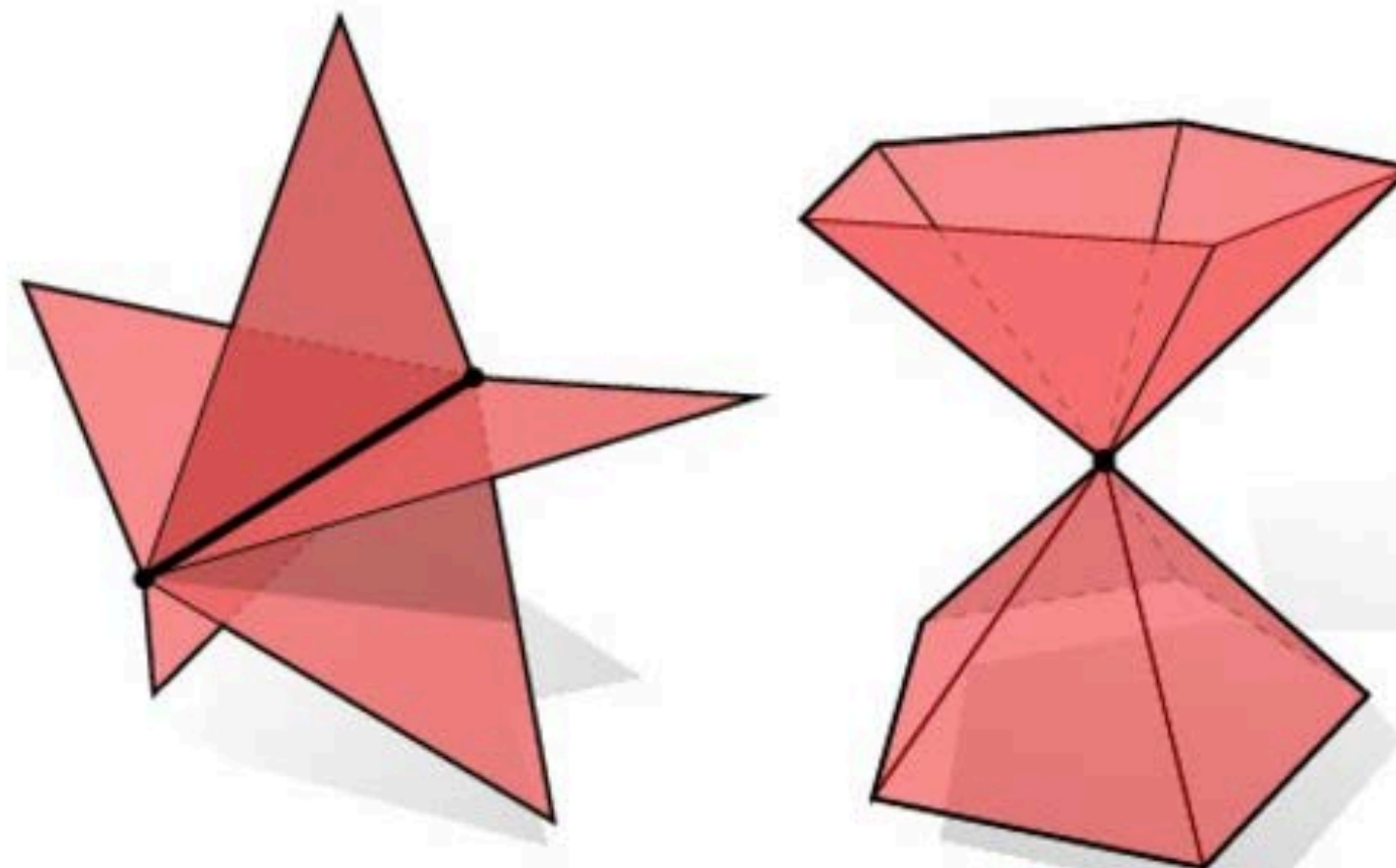


# Surfaces

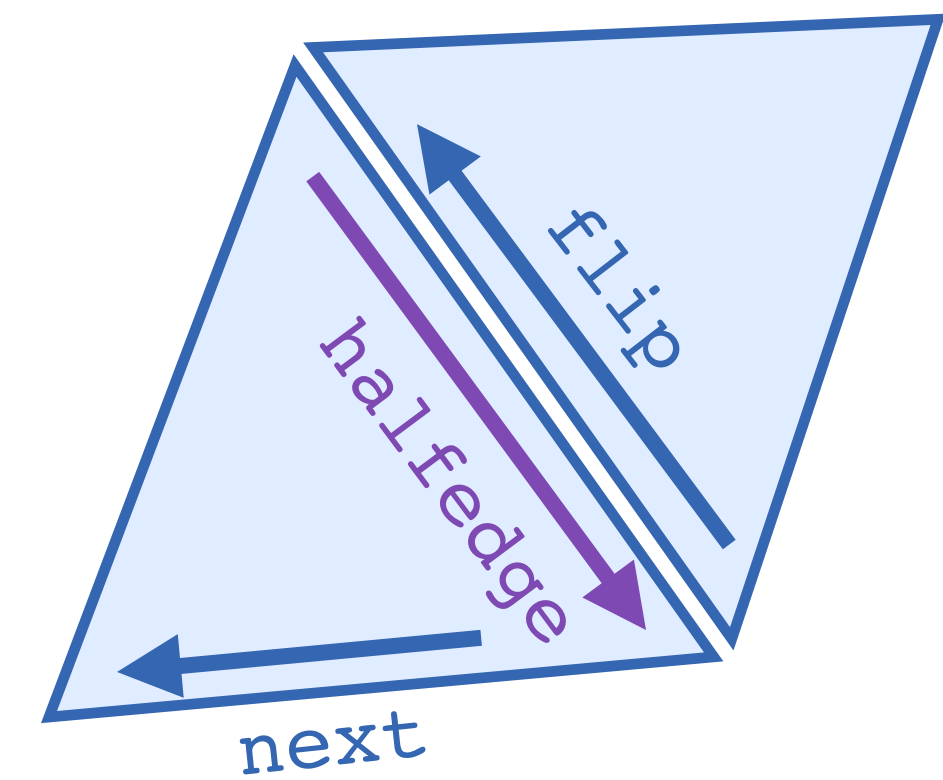
- A few weeks ago, we had a brief overview of surfaces
  - ▶ Triangle meshes
  - ▶ Modeling: generate surfaces by spline or subdivision



manifold (valid) mesh



non-manifold mesh

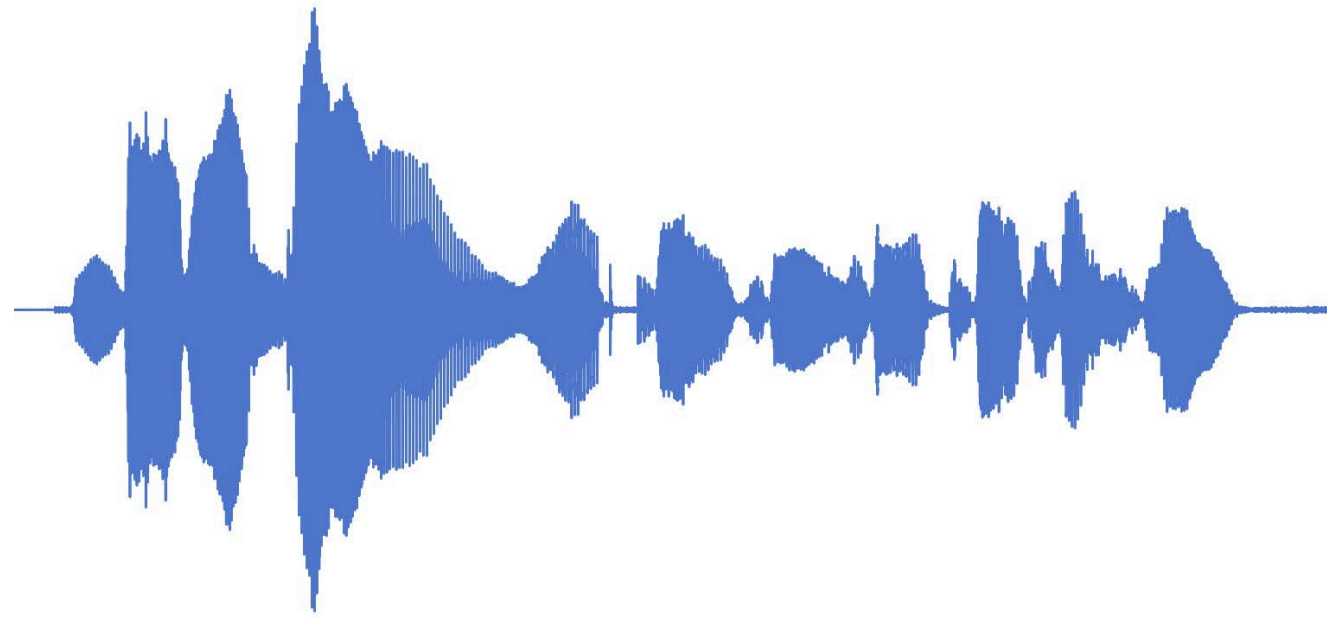


half-edge  
data structure



# Geometry processing

- View discrete surface as a form of “signal data”
  - ▶ Traditional signal data: audio & images

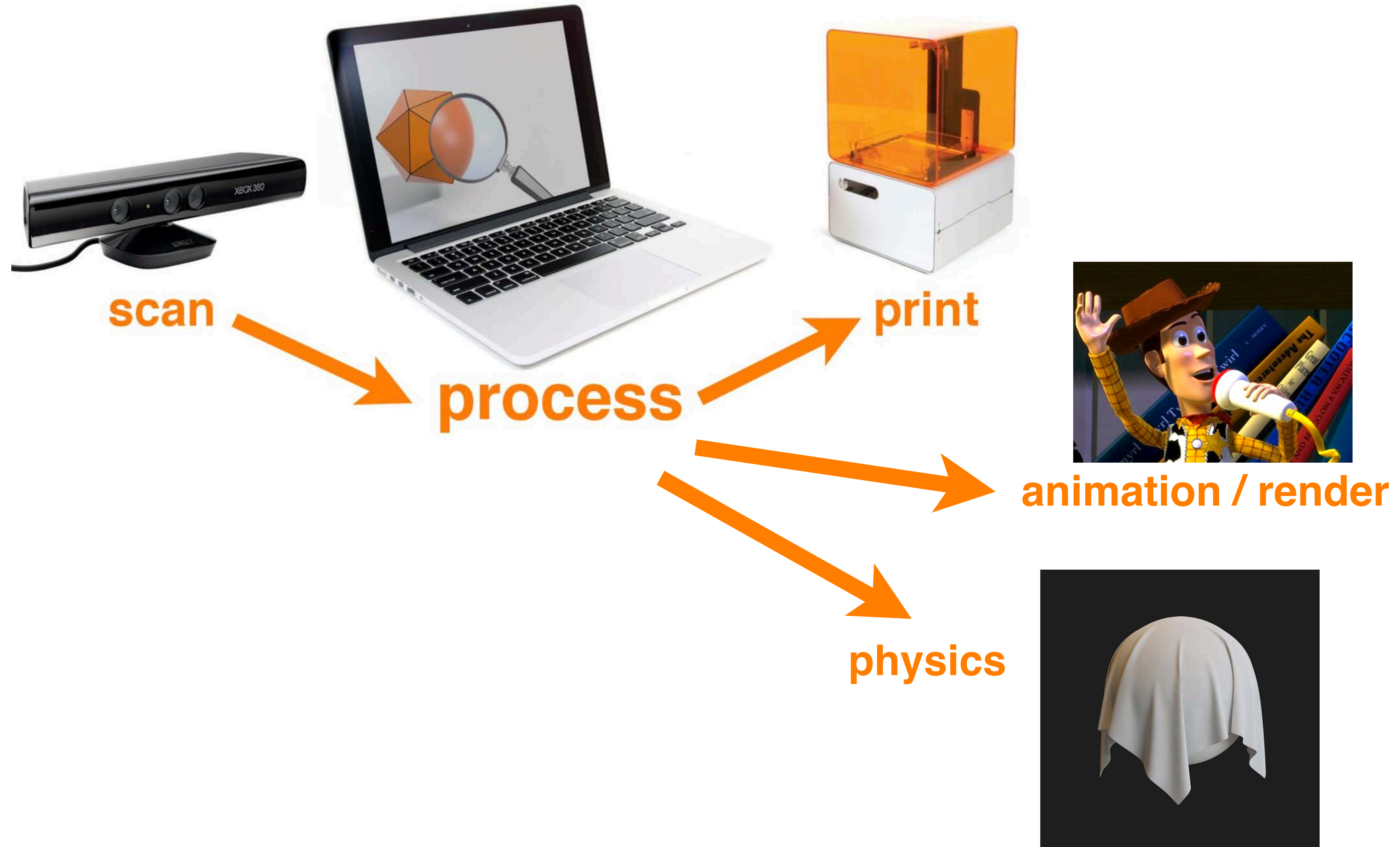


- ▶ Geometric signal
- ▶ Upsampling / downsampling / filtering / aliasing



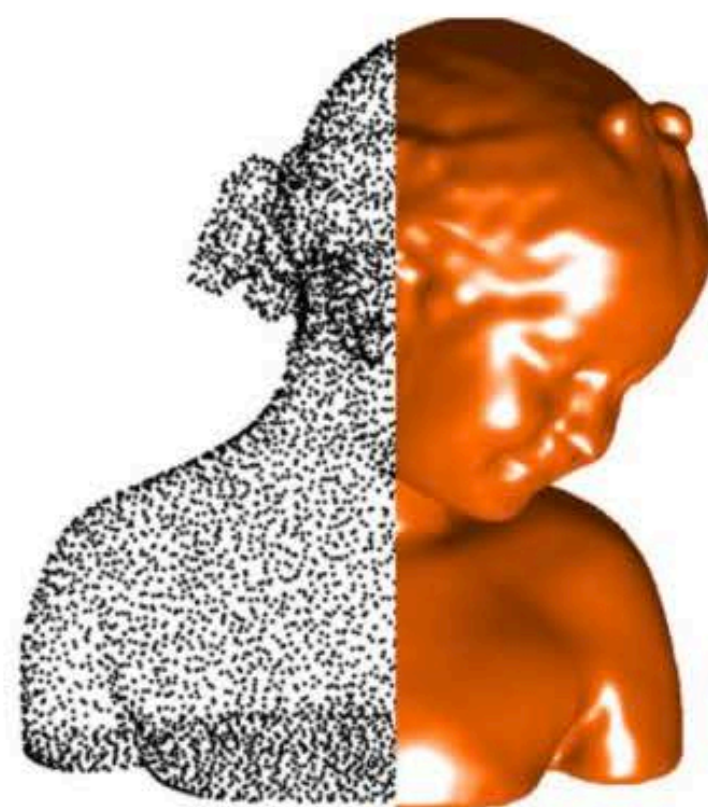


# Geometry processing

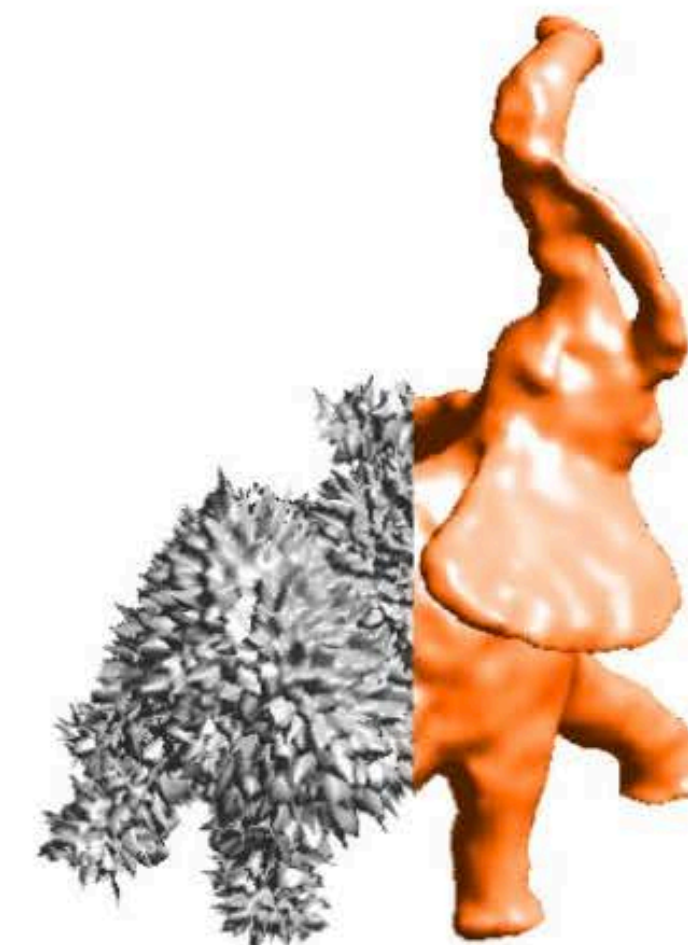




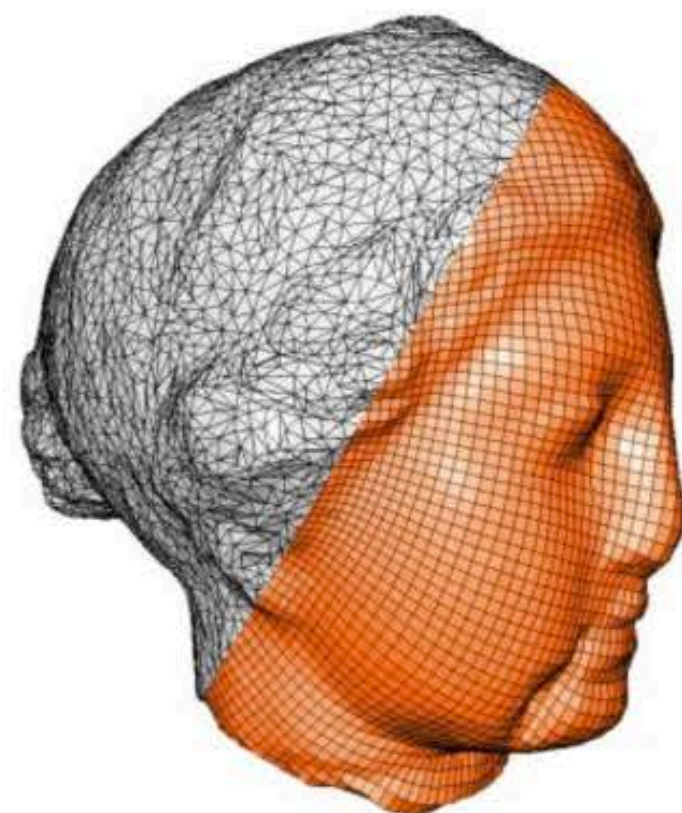
# Tasks of geometry processing



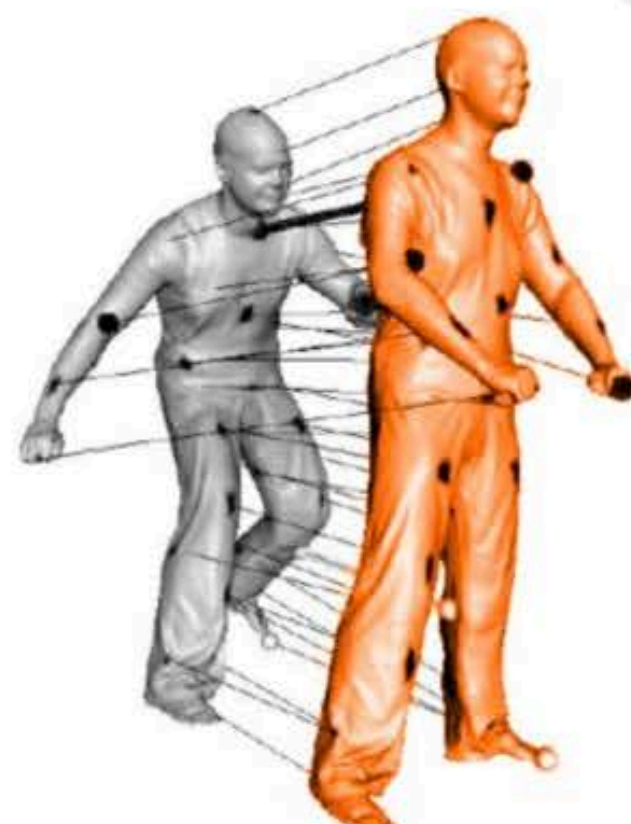
reconstruction



filtering



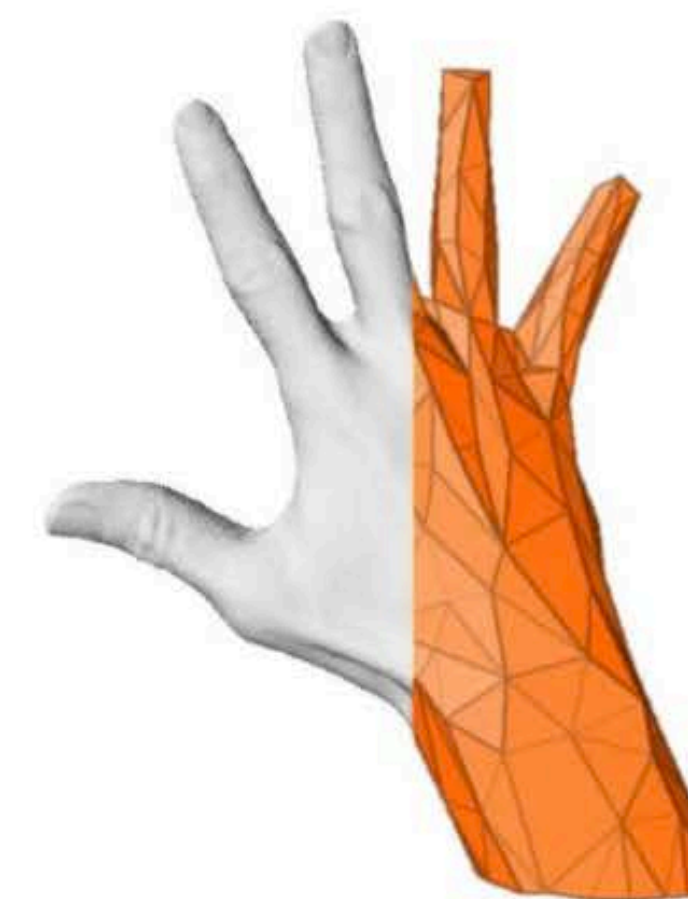
remeshing



shape analysis



parameterization

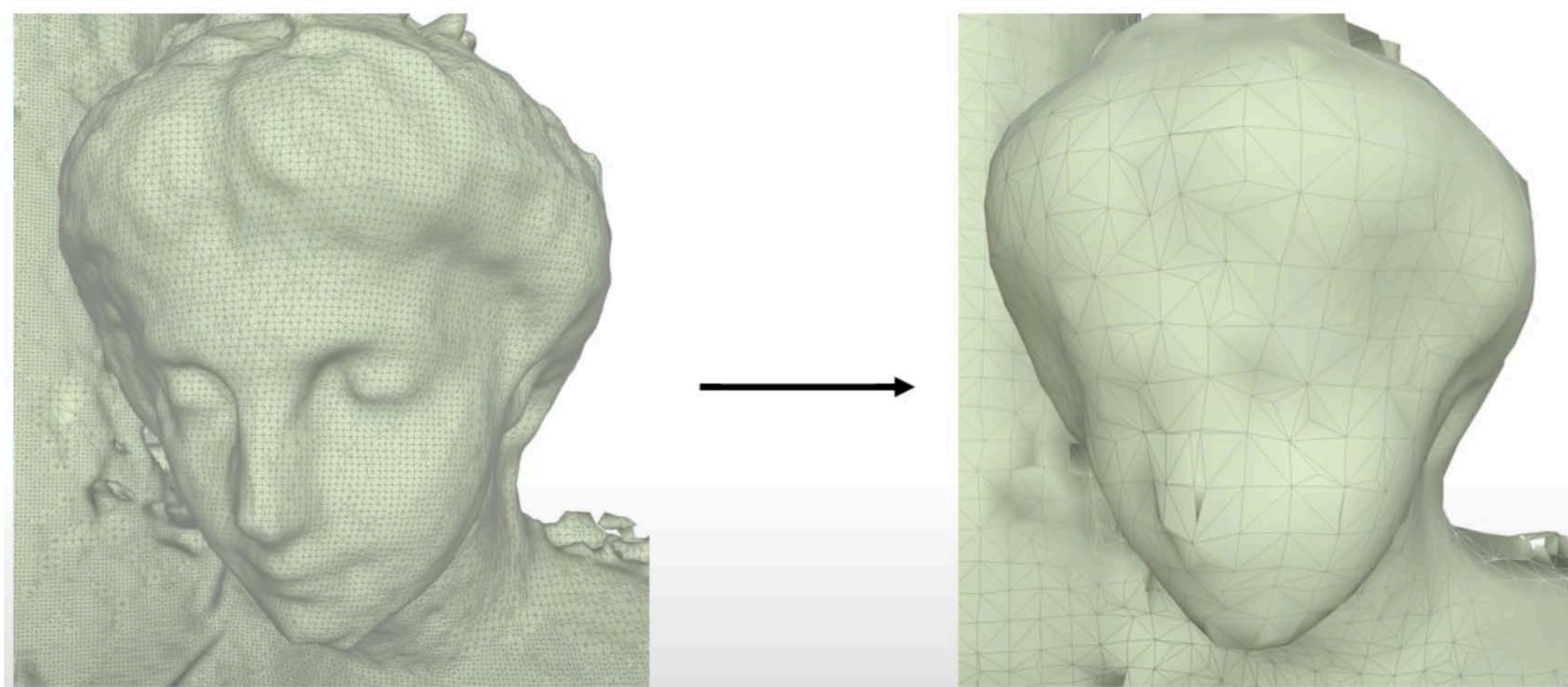


compression

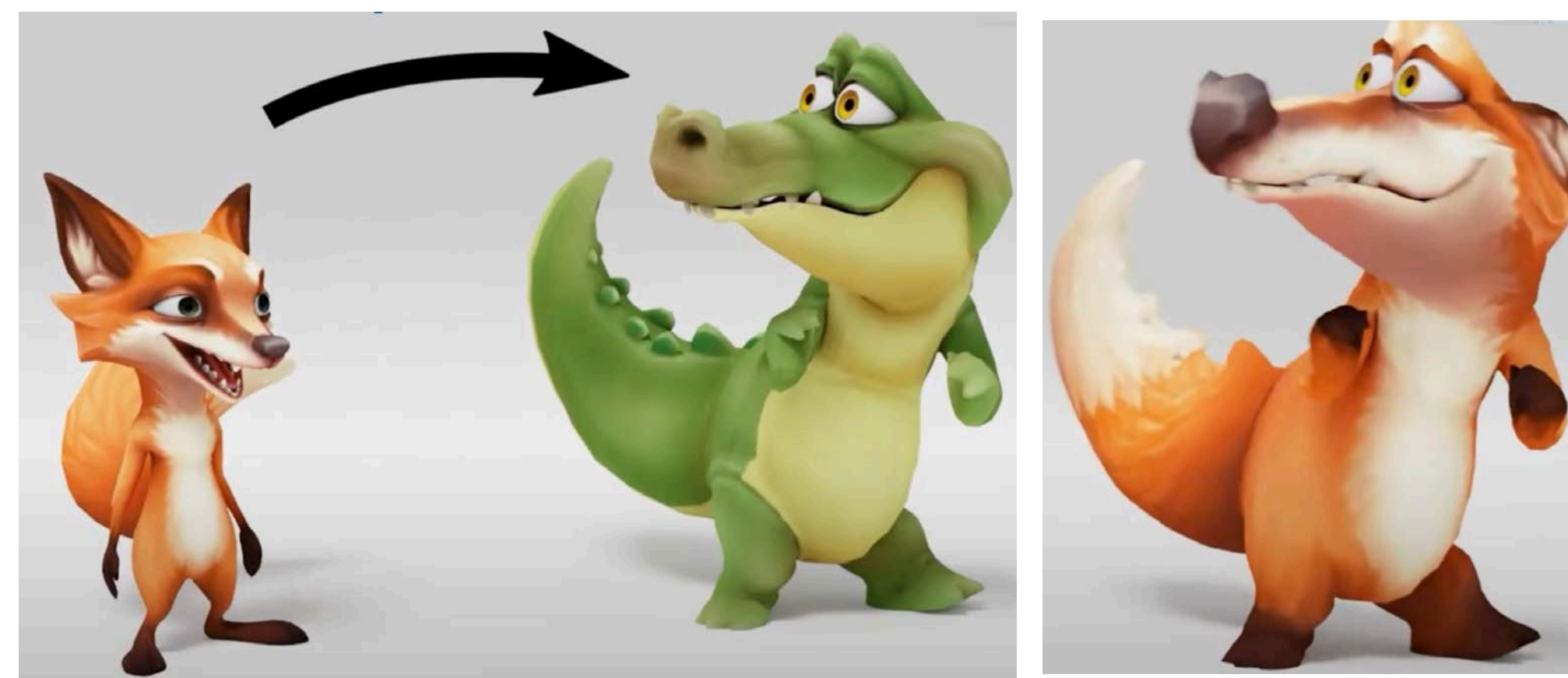


# Symposium on Geometry Processing

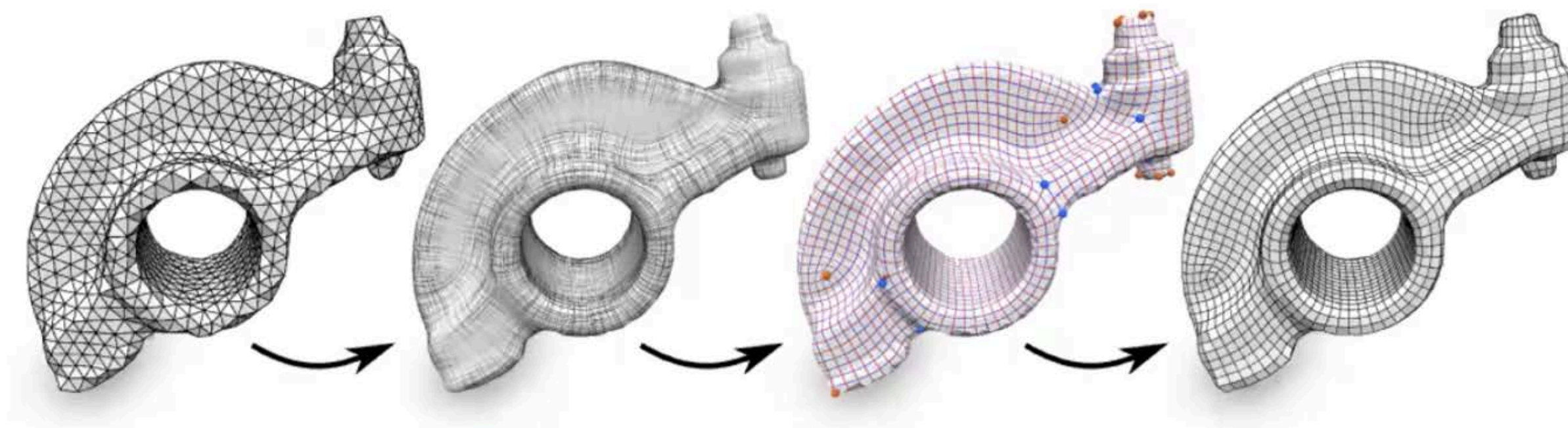
- SGP summer school  
<http://school.geometryprocessing.org/>



Shape approximation



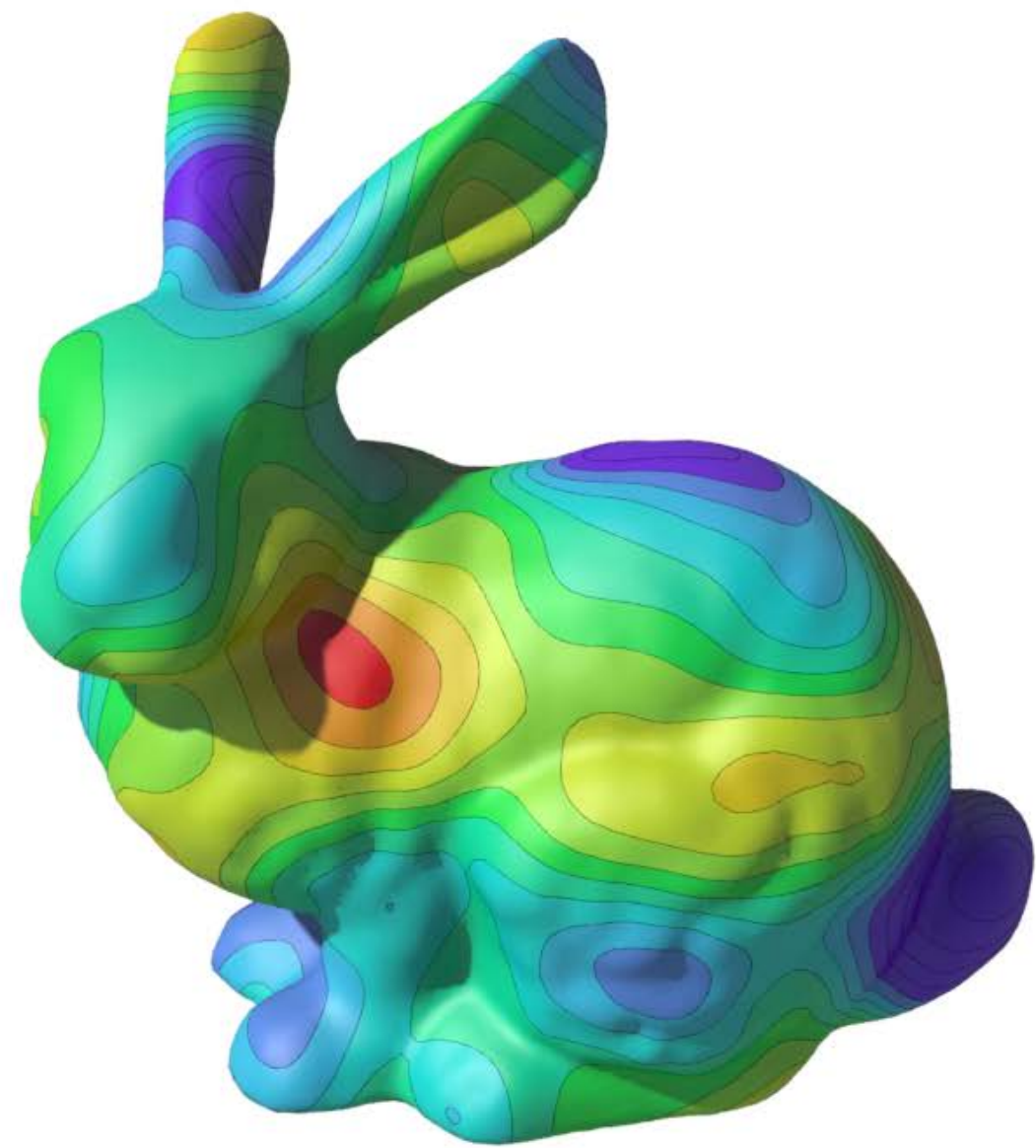
Maps between surfaces



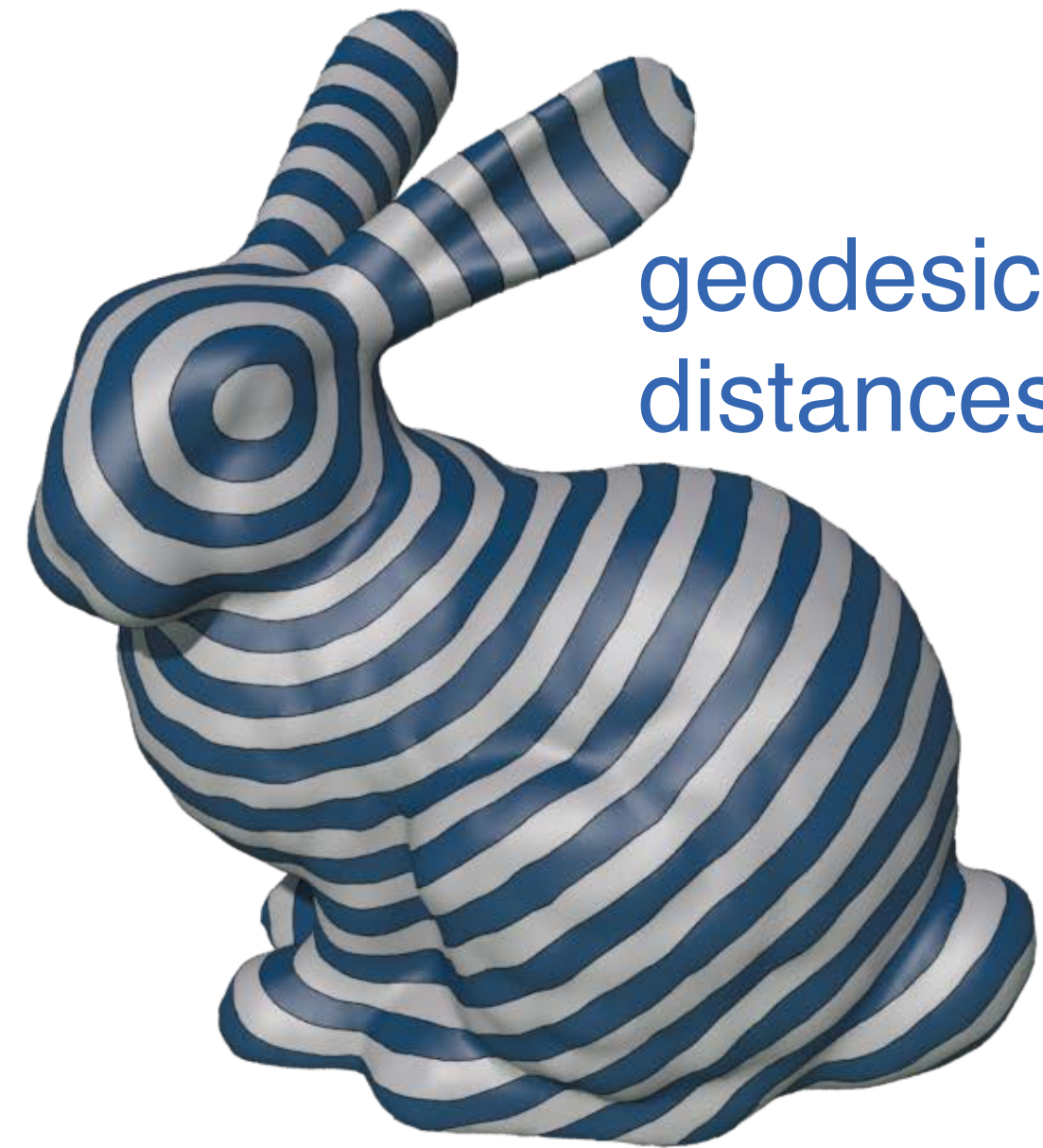
Directional field



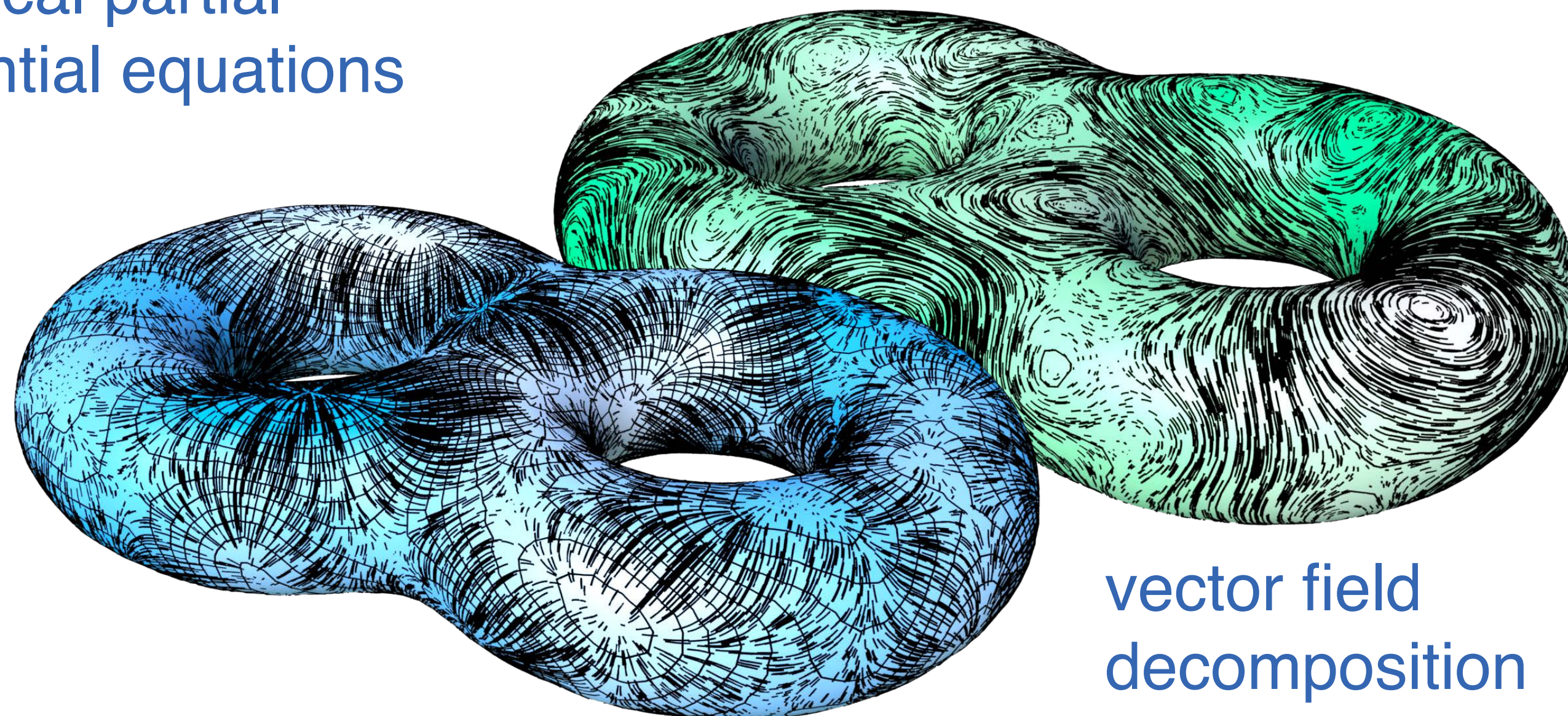
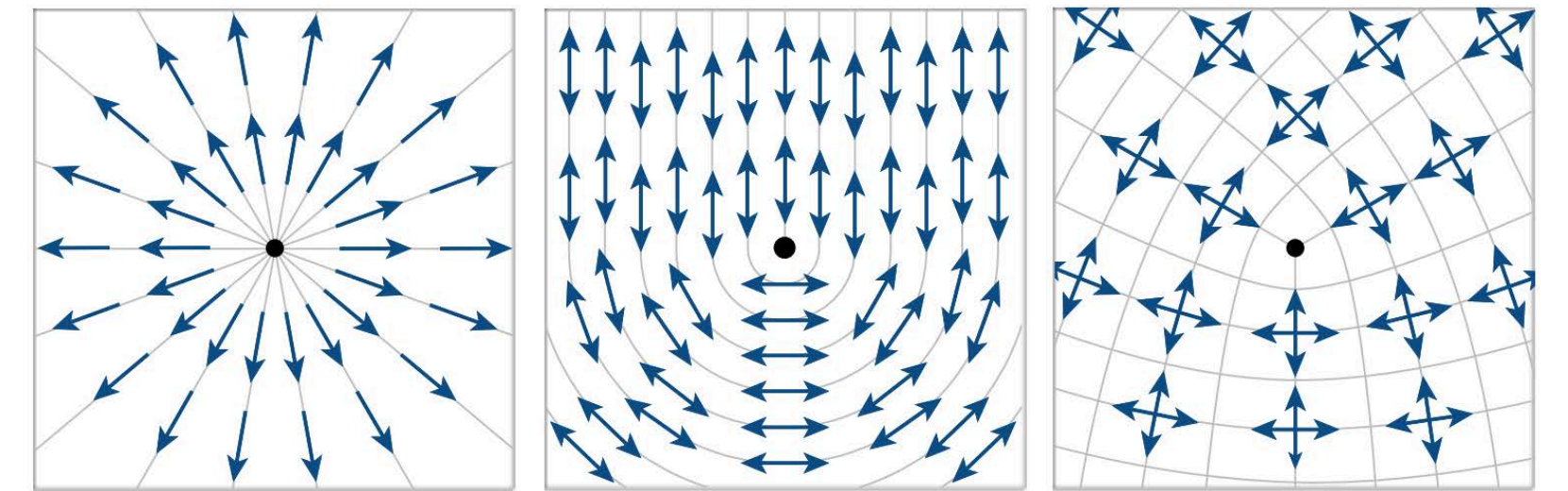
# CSE274 (discrete differential geometry)



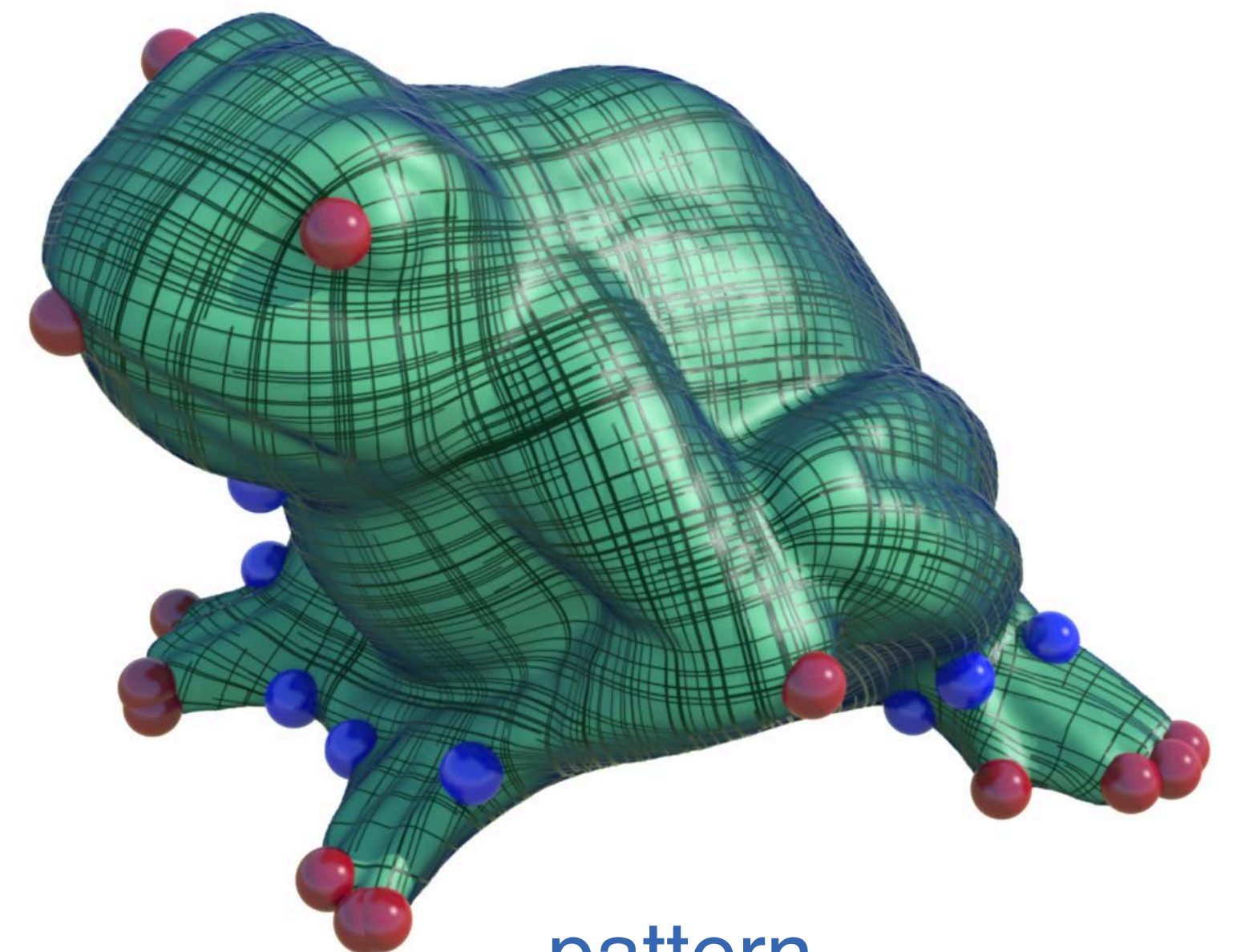
numerical partial differential equations



geodesic distances



vector field decomposition



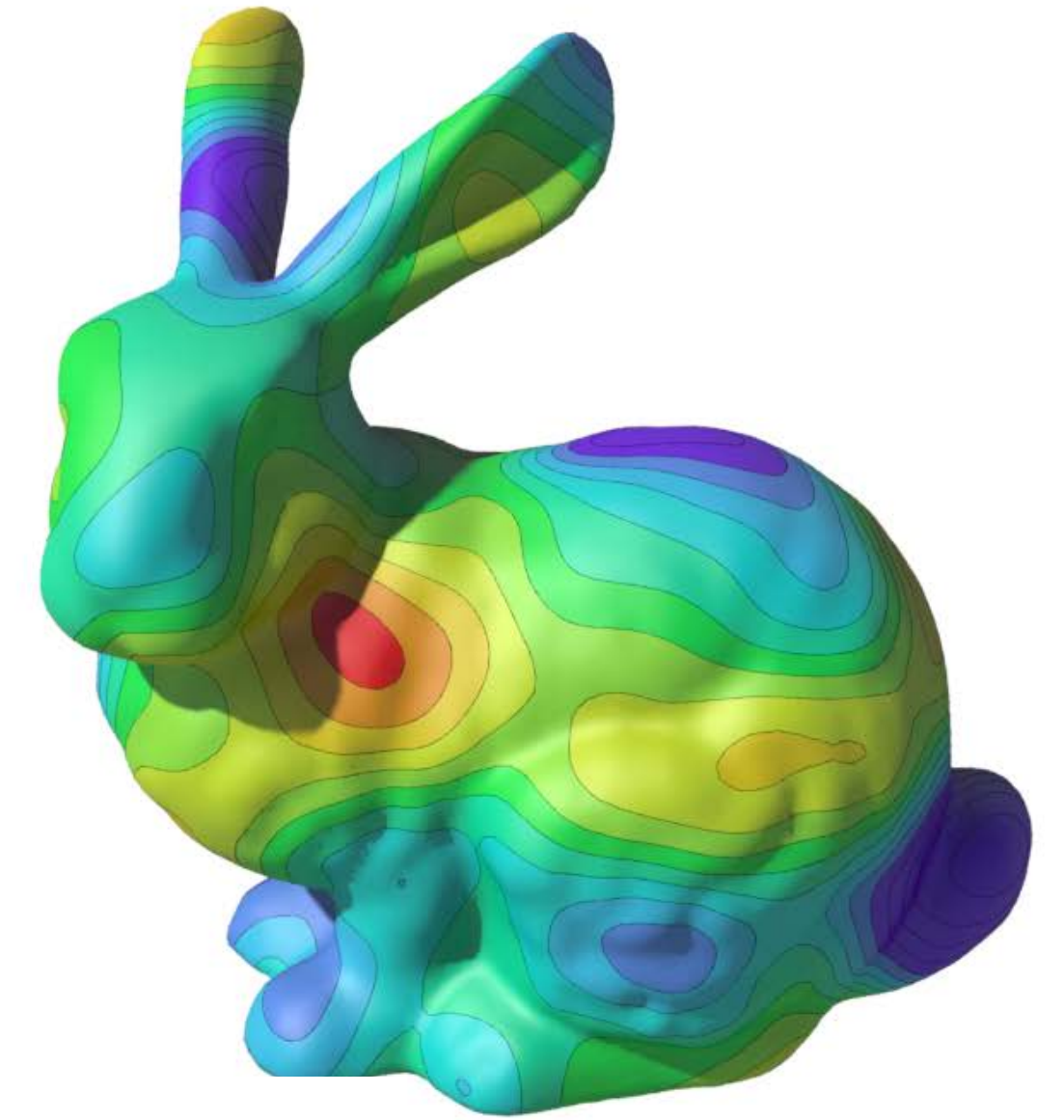
pattern design



# Today

- Surface processing using **Laplacian**

	1	
1	-4	1
	1	



- Another topic: **Remeshing**





# Laplacian

- Laplacian
- Remeshing



# 2nd derivative

- For a function of 1 variable  $f(t)$   
the **Laplacian** of the function is its 2nd derivative

$$(\Delta f)(t) := \frac{d^2 f}{dt^2}(t)$$

- Laplacian is usually denoted by  $\Delta$  or  $\nabla^2$  or  $L$
- 2nd derivative on 1D measures the difference between the function value at a point and the averaged function value around that point.
- Laplacian is a 2nd derivative on a 2D or 3D or surface domain measuring the deviation of value from the neighbor average



# A simple discrete Laplacian

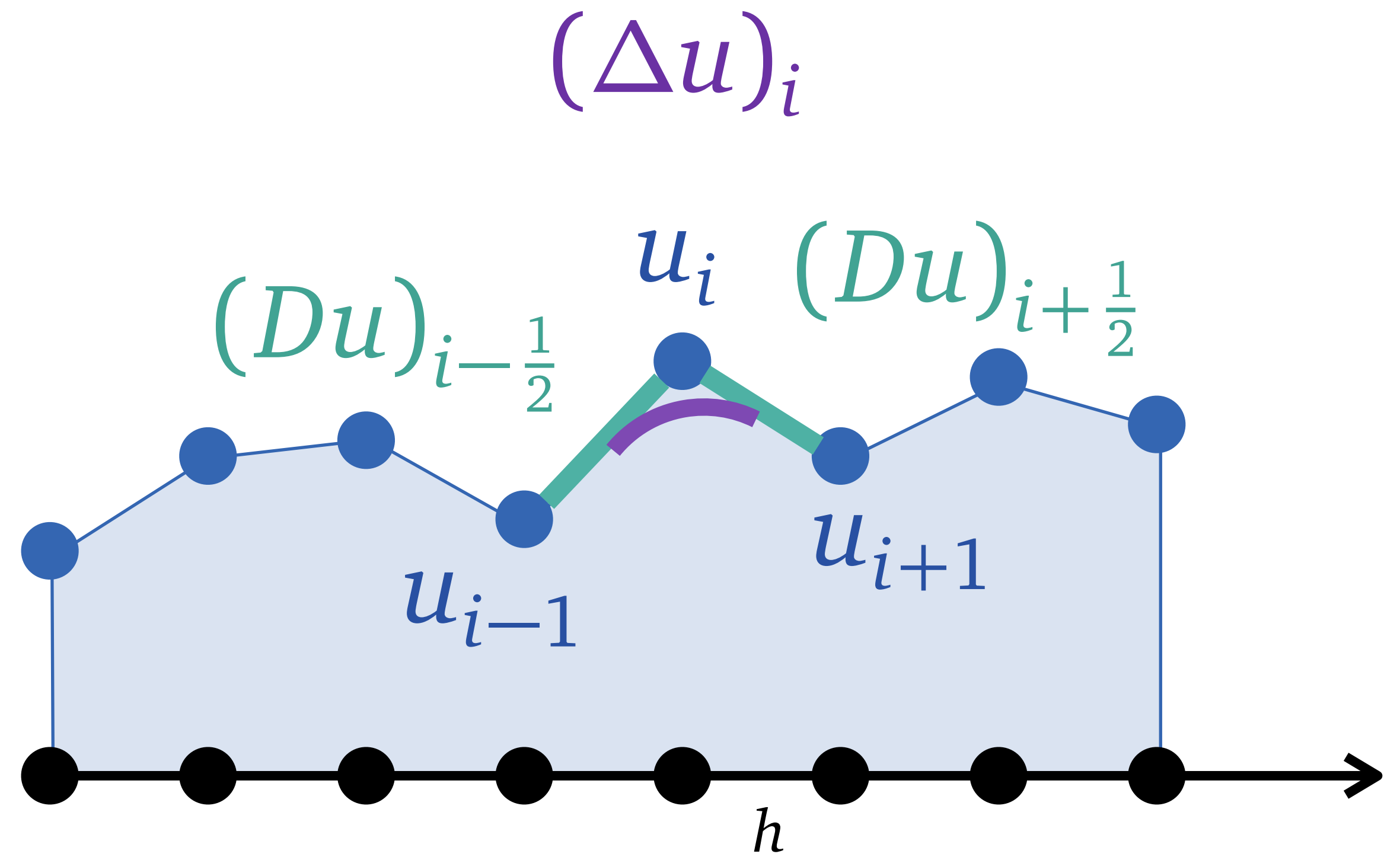
- In 1D  $\Delta u(x) = u''(x)$
- Discretize 1D space into uniform grid
- grid size  $h$
- Discrete 1st derivative

$$(Du)_{i-\frac{1}{2}} = \frac{u_i - u_{i-1}}{h}$$

$$(Du)_{i+\frac{1}{2}} = \frac{u_{i+1} - u_i}{h}$$

- Discrete 2nd derivative

$$(\Delta u)_i = \frac{(Du)_{i+\frac{1}{2}} - (Du)_{i-\frac{1}{2}}}{h} = \frac{u_{i-1} - 2u_i + u_{i+1}}{h^2} = \frac{2}{h^2} \left( \frac{u_{i-1} + u_{i+1}}{2} - u_i \right)$$



average of  
neighbor      function  
value



# In image processing (2D domain)

- Laplace filter

	1	
1	-4	1
	1	

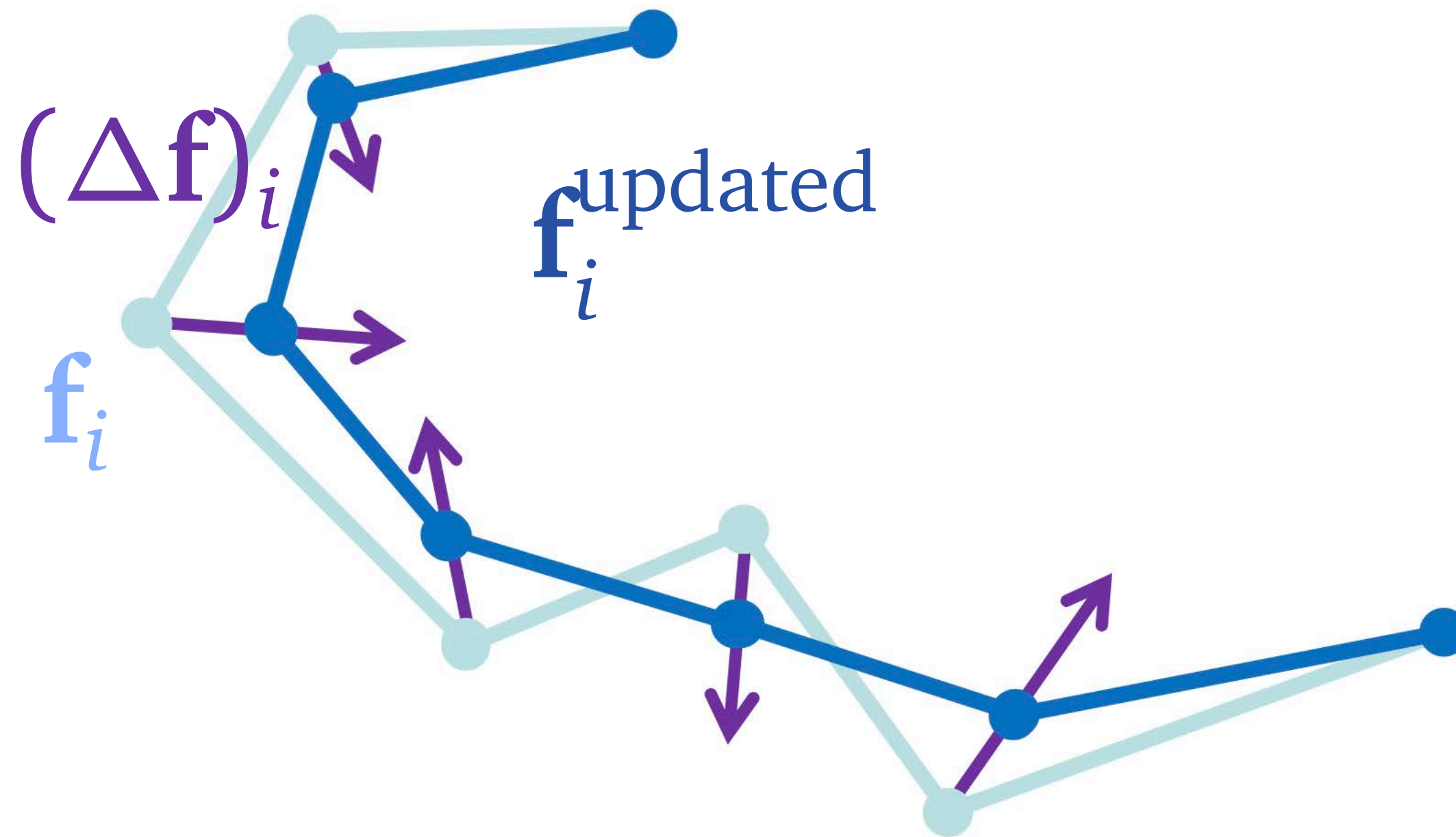
- In multivariable calculus  $\Delta = \frac{\partial^2}{\partial x^2} + \frac{\partial^2}{\partial y^2}$
- Useful for edge detection





# For smoothing

- Laplacian measures how much the neighbor deviates from center



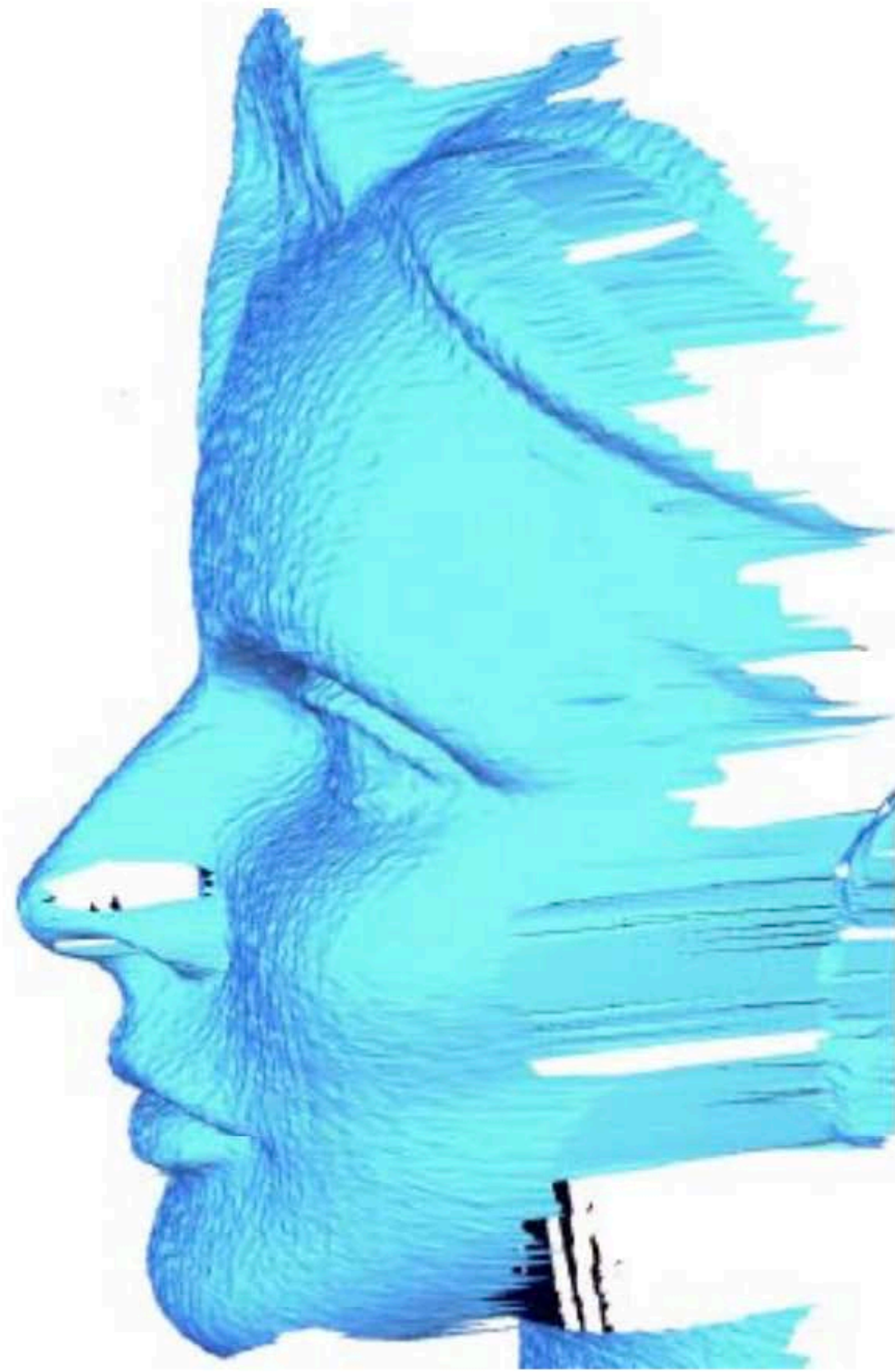
$$\frac{\partial}{\partial t} \mathbf{f}_i = (\Delta \mathbf{f})_i$$

*This is also the heat diffusion equation*

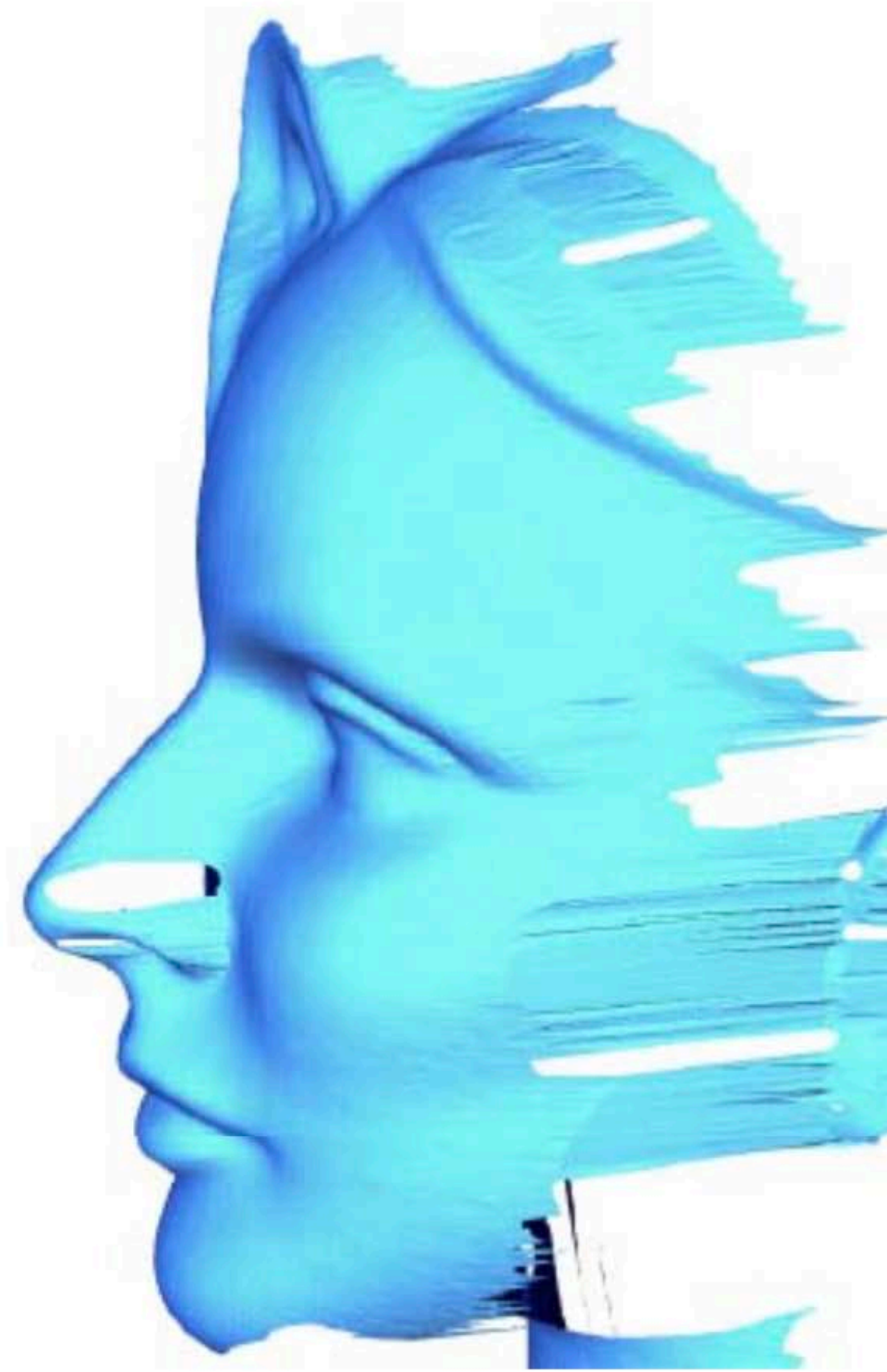
- Laplacian of the vertex position is proportional to the (mean) curvature of the curve/surface



# For smoothing



raw 3D scanning data



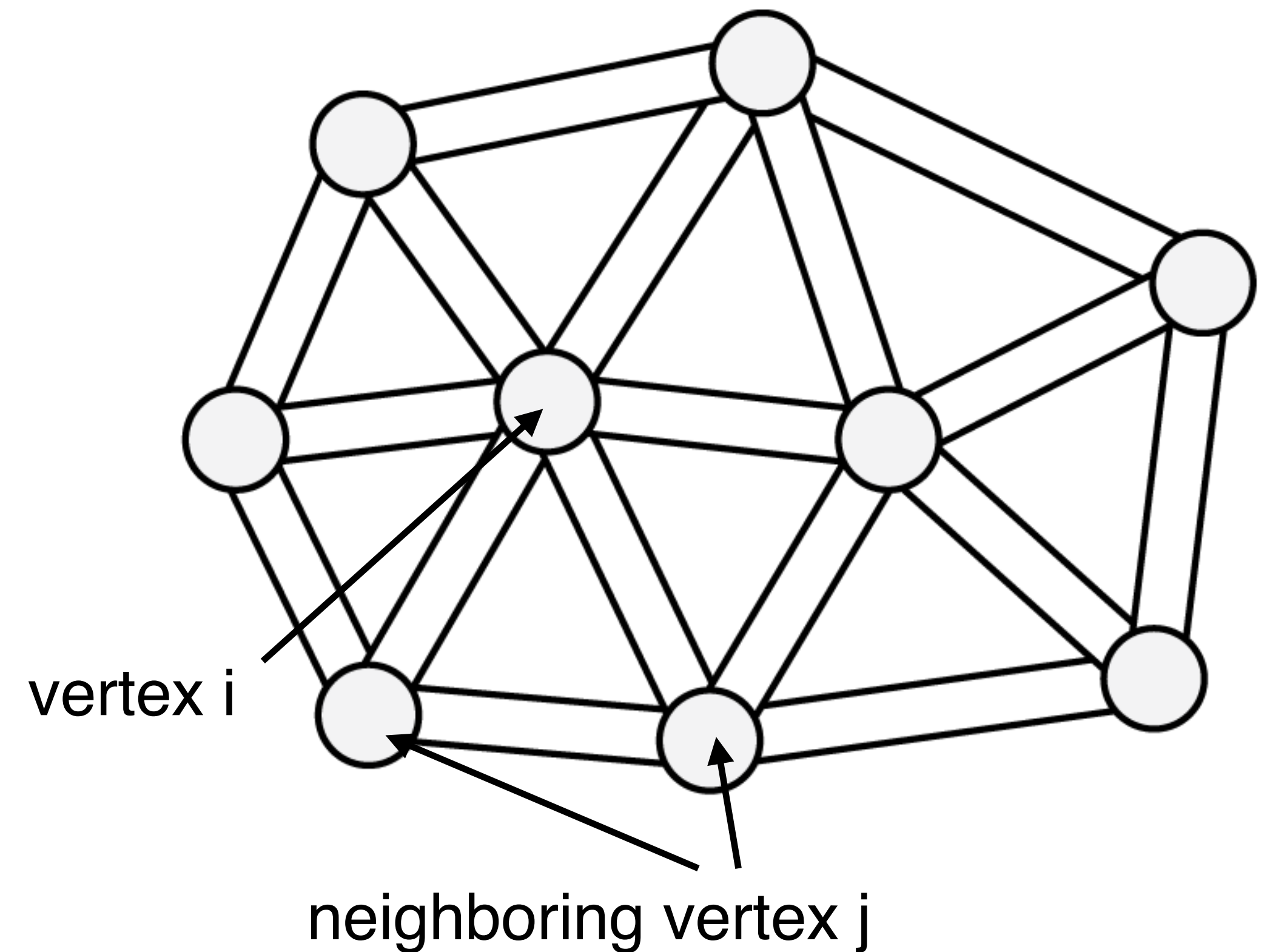
after a few steps of Laplacian smoothing  
(mean curvature flow)



# Laplacian on mesh or graph

$$(\Delta u)_i = \sum_{j \in \text{Neighbor}(i)} \overset{\text{edge weights}}{w_{ij}} (u_j - u_i)$$

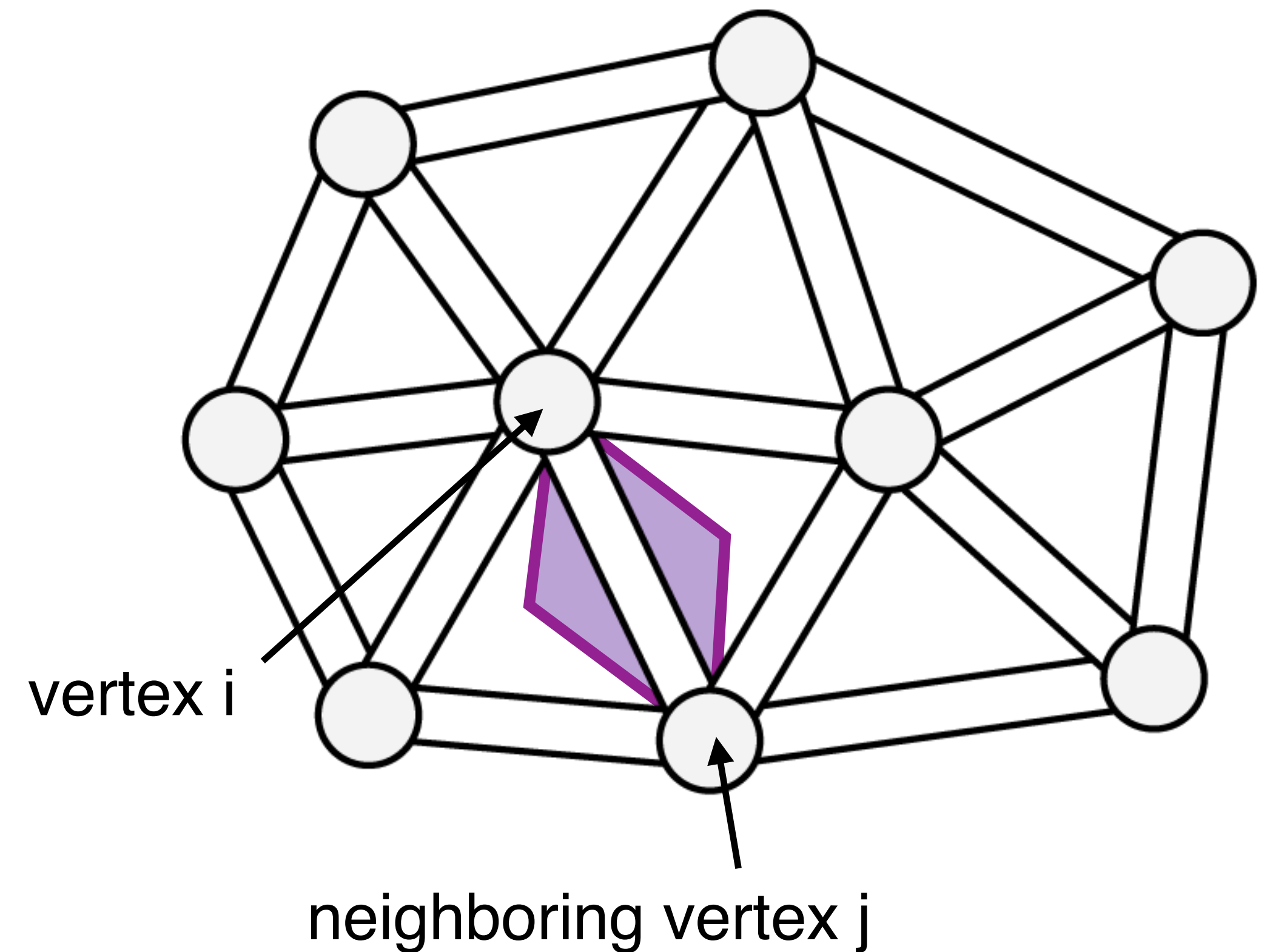
- In graph theory, people usually take edge weights to be all 1 (graph laplacian)
- In geometry processing, edge weights are chosen to mimic the effective *conductivity* on the edge





# Laplacian on mesh or graph

$$w_{ij} = \frac{\text{width of the kite}}{\text{edge length}}$$

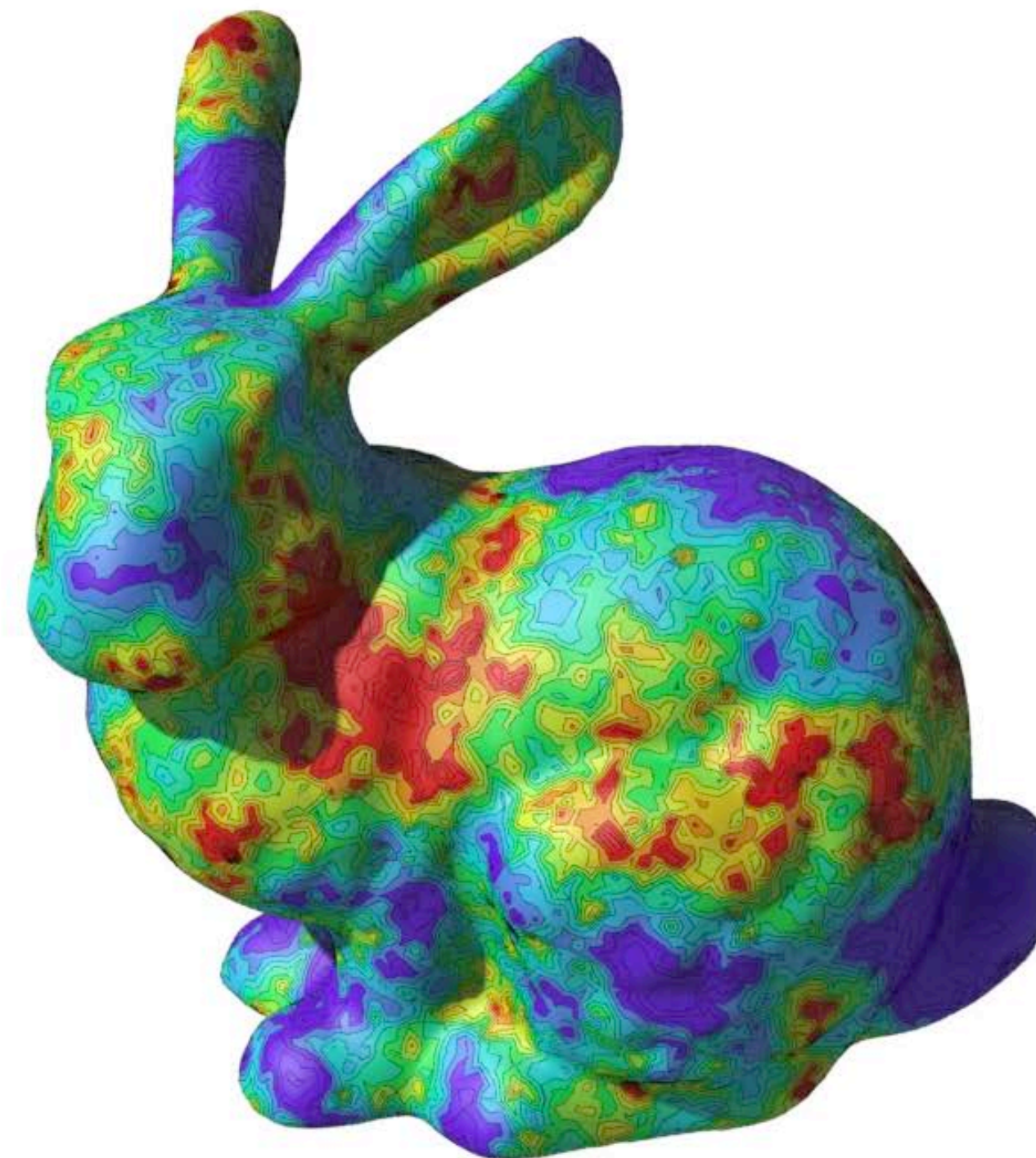


- If the mesh quality is “good” then the discrete Laplacian approximates the continuous Laplacian well



# Laplacian on mesh or graph

- With Laplacian, many image processing techniques (smoothing, curvature detection, reconstructions) can be done on surfaces.
- Many physical equation (usually only written in Cartesian coordinates) can be simulated on surfaces.





# Remeshing

- Laplacian
- Remeshing



# Today

- Let's talk about: **Remeshing**



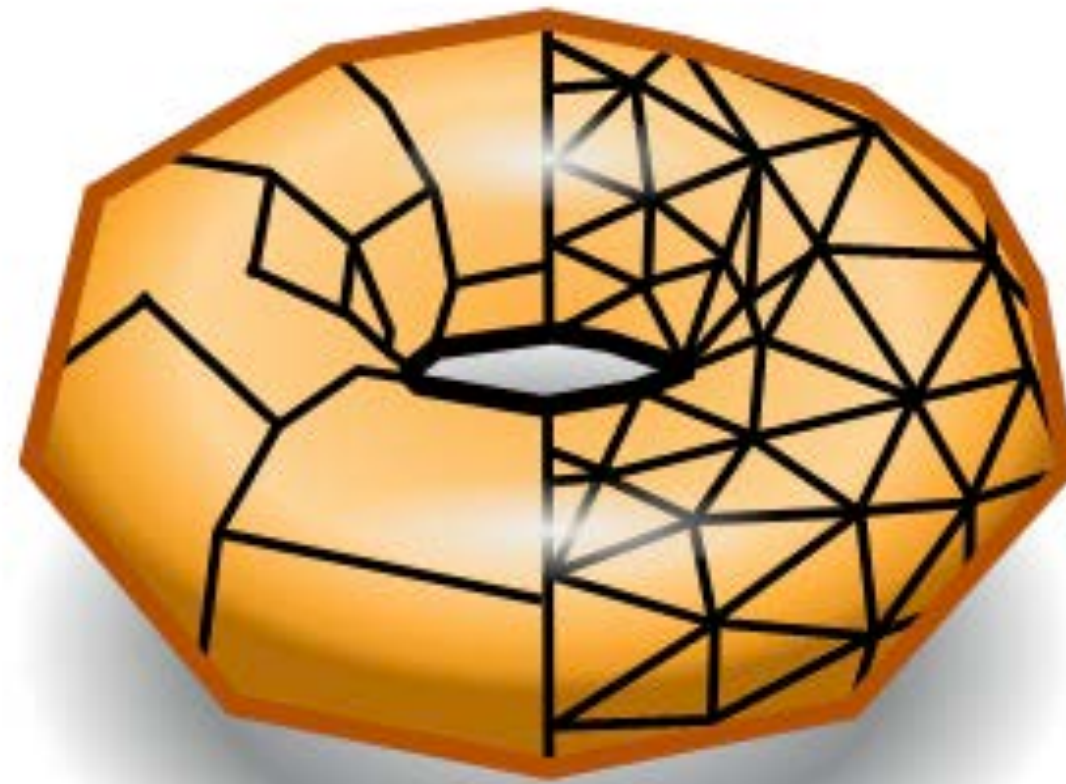


# Today

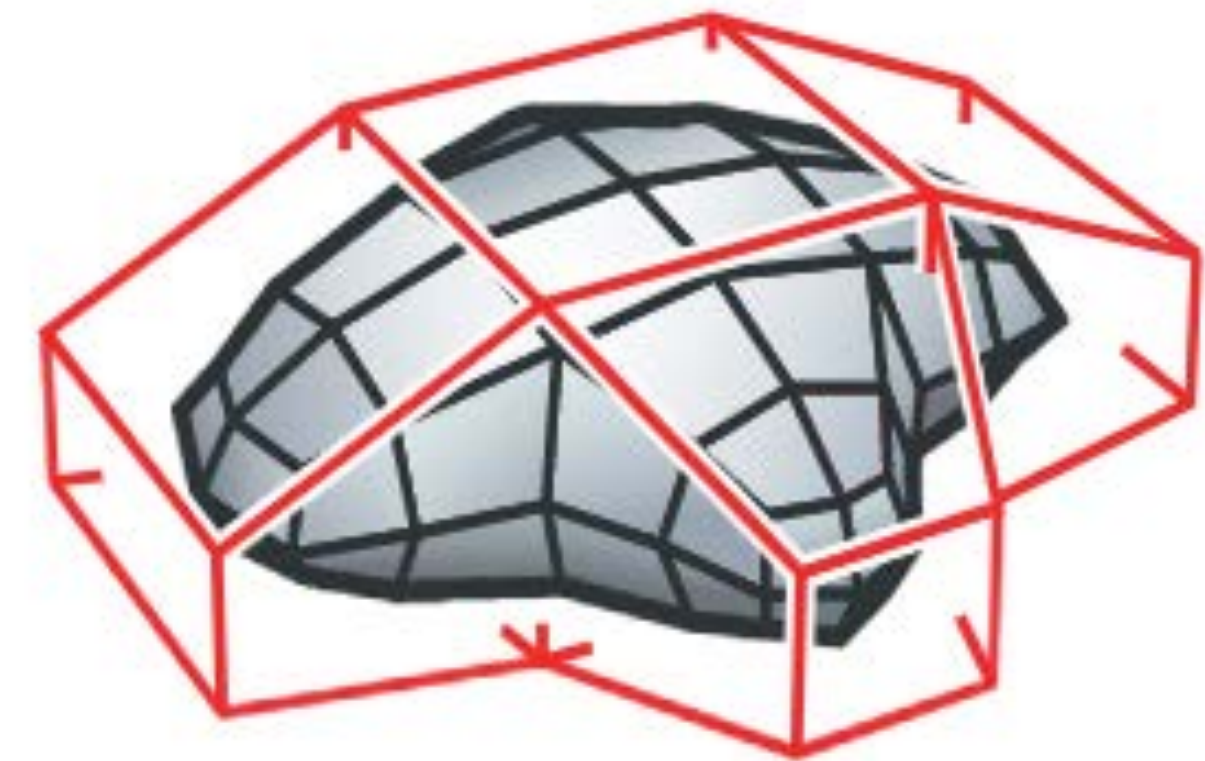
- Let's talk about: **Remeshing**



Reduce polygon



Obtain better triangle mesh

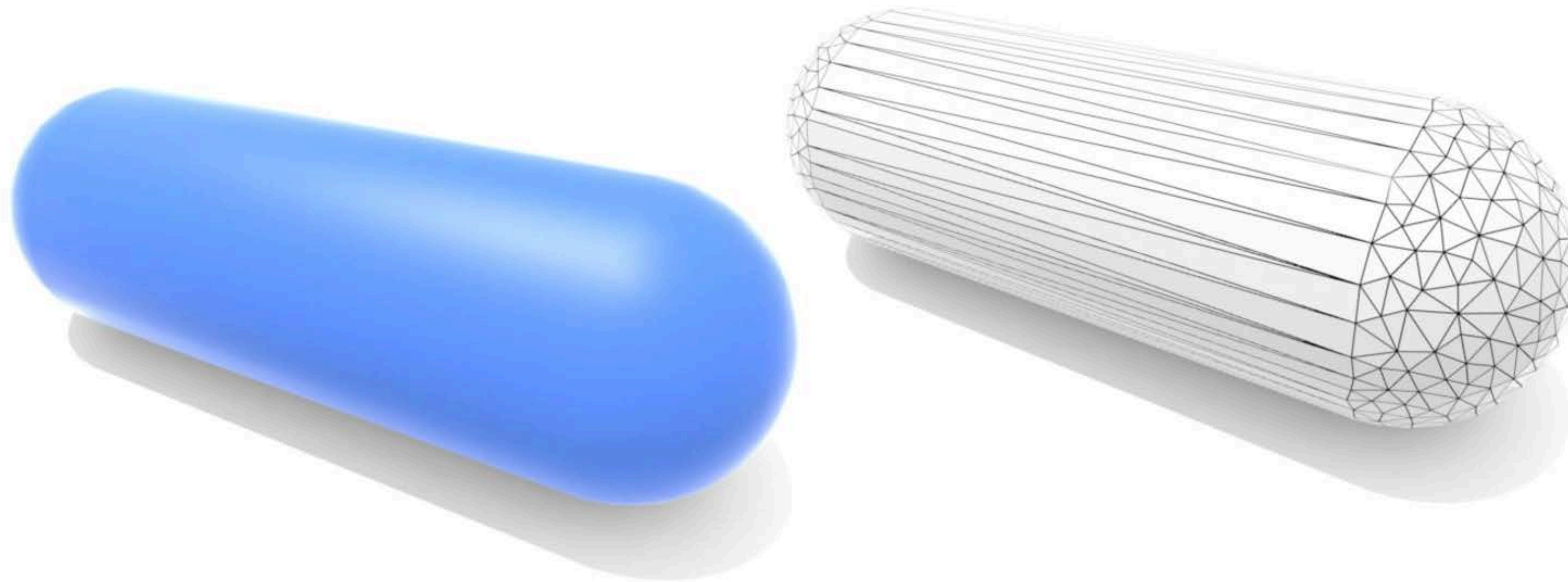


Subdivision



# What is a “good” mesh?

- One idea: good approximation of the original shape!
- Keep only elements that contribute information about shape
- Add element where, e.g., curvature is large





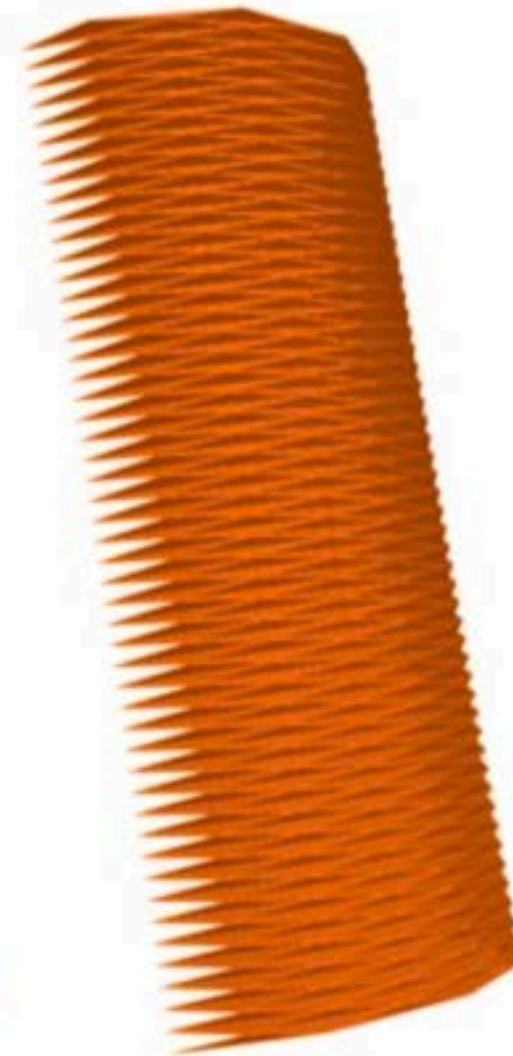
# What is a “good” mesh?

- One idea: good approximation of the original shape!
- Keep only elements that contribute information about shape
- Add element where, e.g., curvature is large
- “Good approximation” can be deceiving sometimes

vertices exactly on smooth cylinder



smooth cylinder

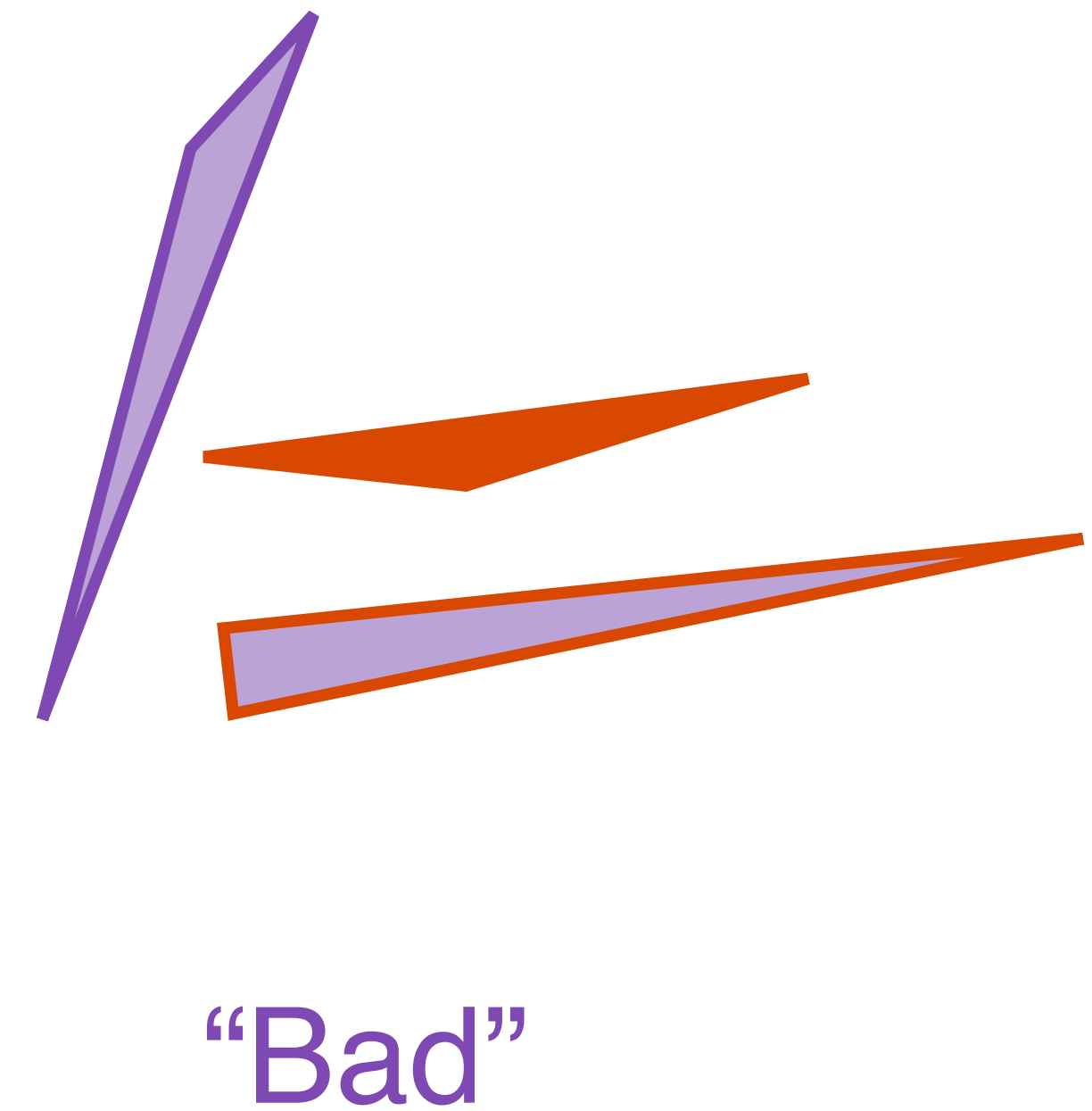
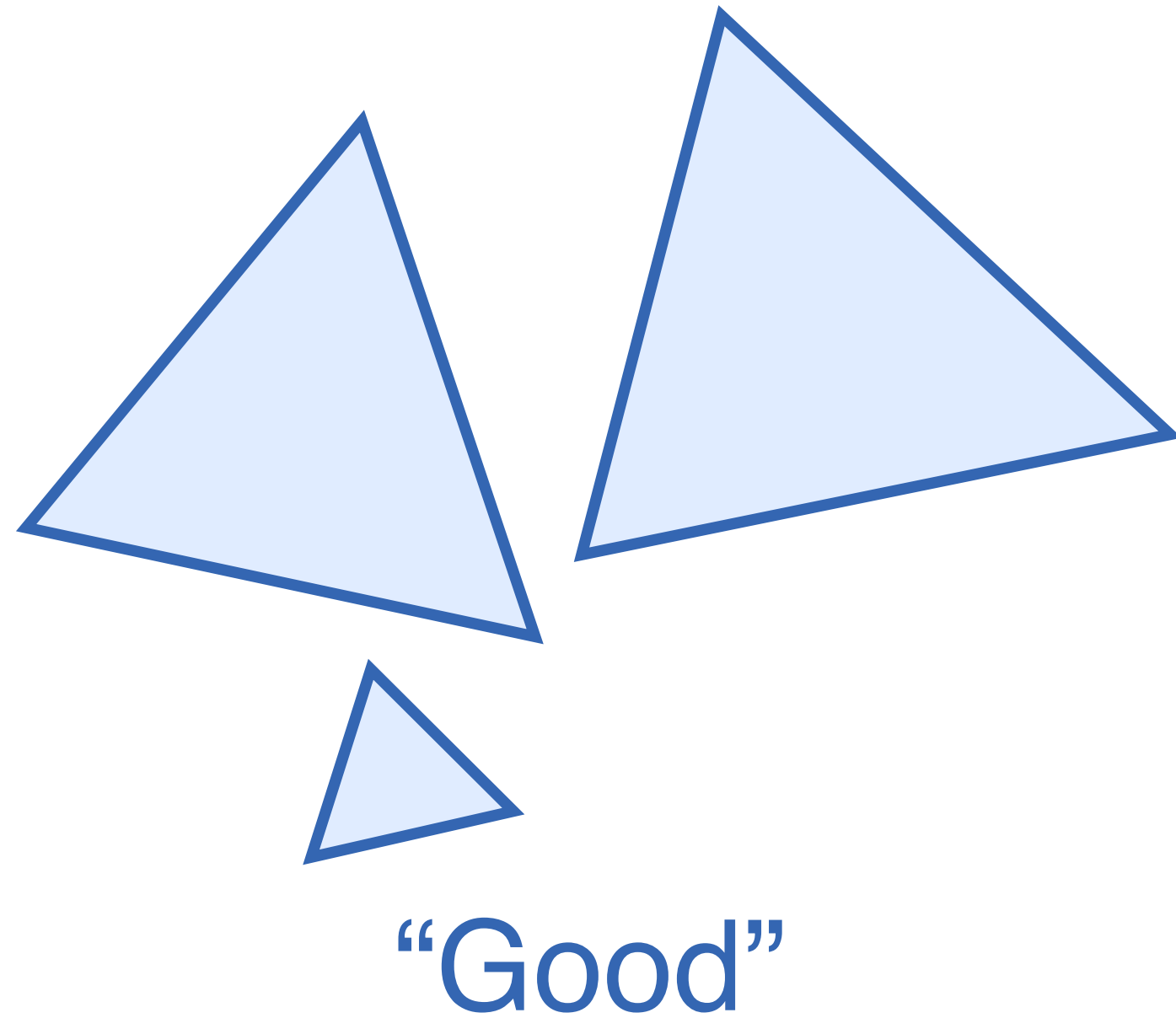


*surface area doesn't converge under refinement*



# What is a “good” mesh?

- Another rule of thumb: *triangle shape*

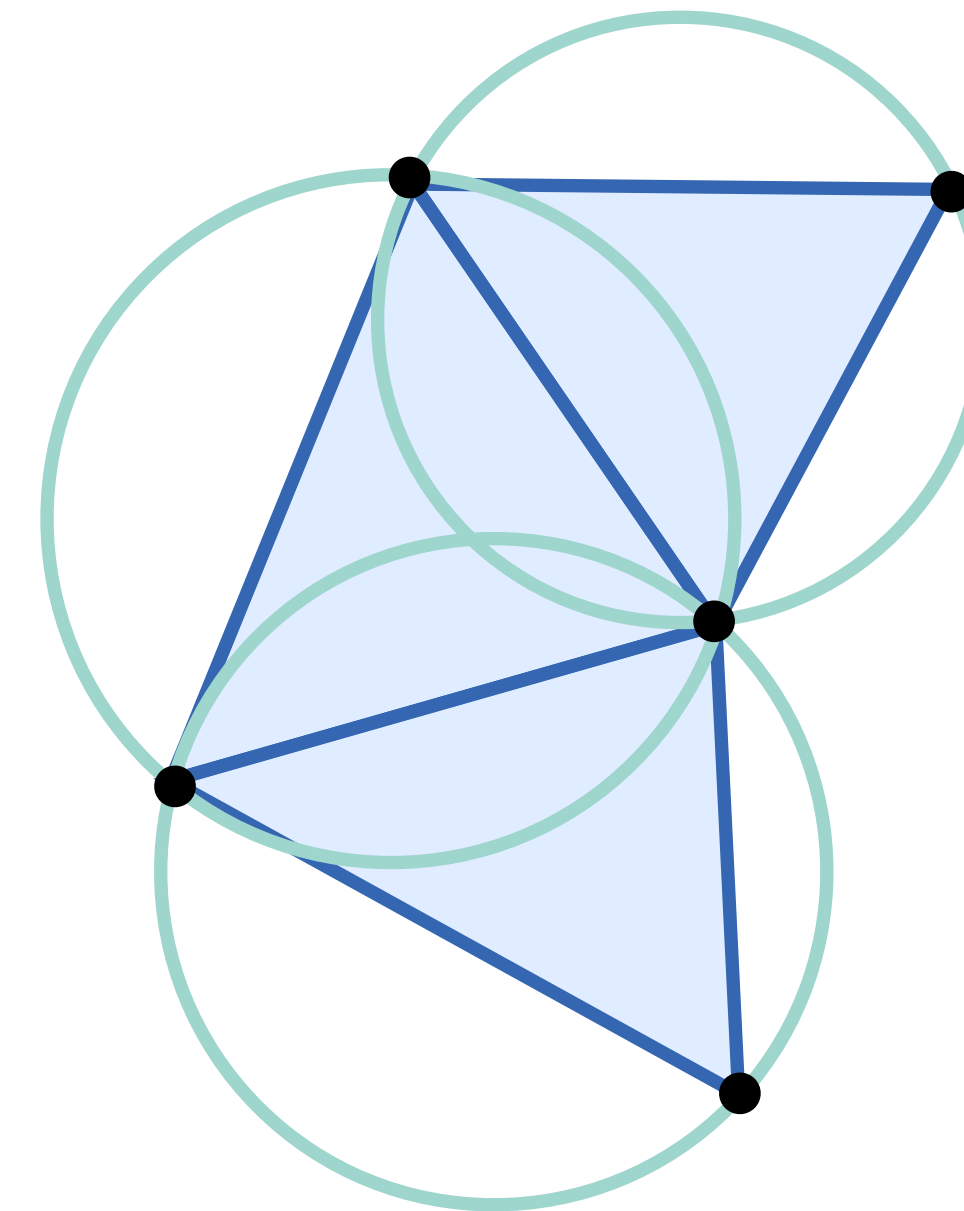


- For example, all angles close to 60 degrees
- A concrete characterization: Delaunay condition

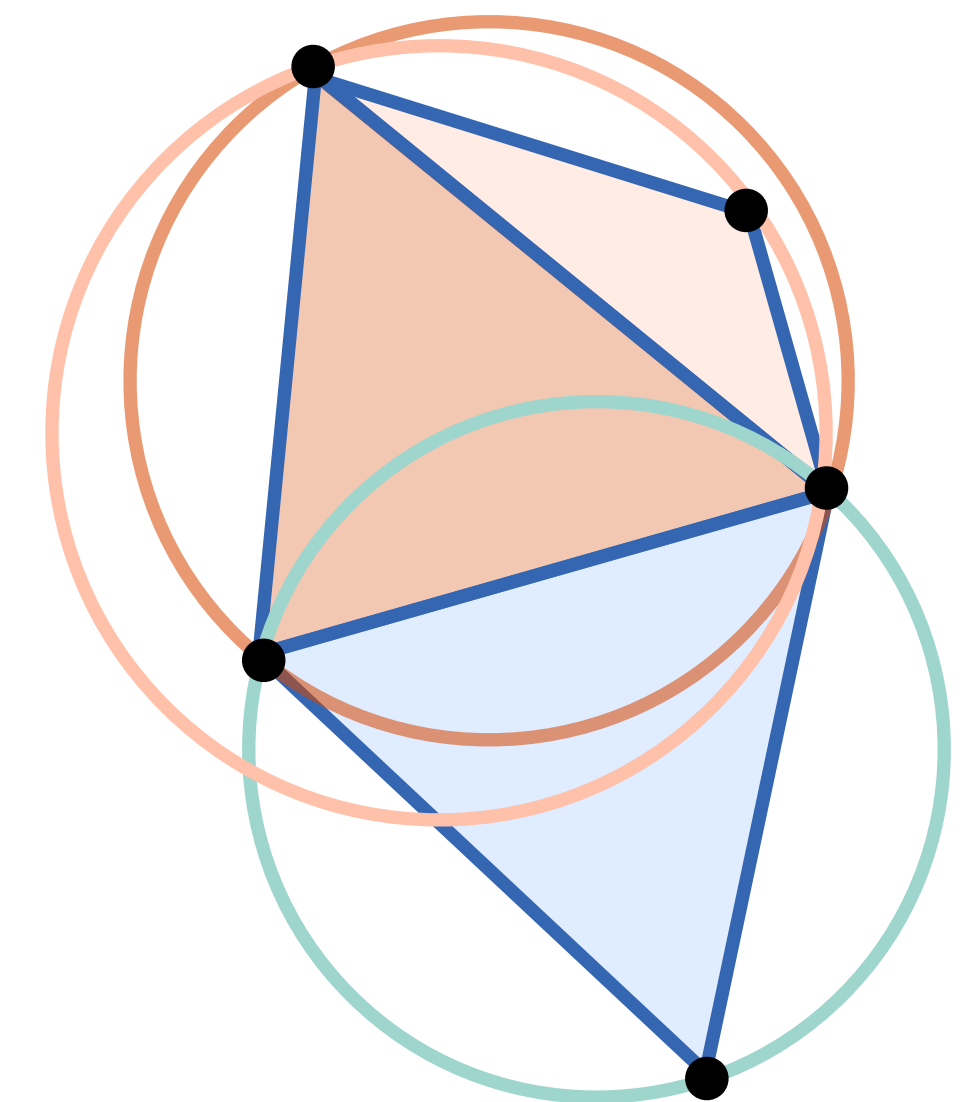


# Delaunay condition

- A triangle mesh is **Delaunay** if the circumcircle of each triangle does not contain any vertex of any adjacent triangle
- Many desirable properties
  - ▶ Helps numerical accuracy / stability
  - ▶ Maximizes minimal angle
  - ▶ Smoothest linear interpolation
- Tradeoffs with efficient shape approximation

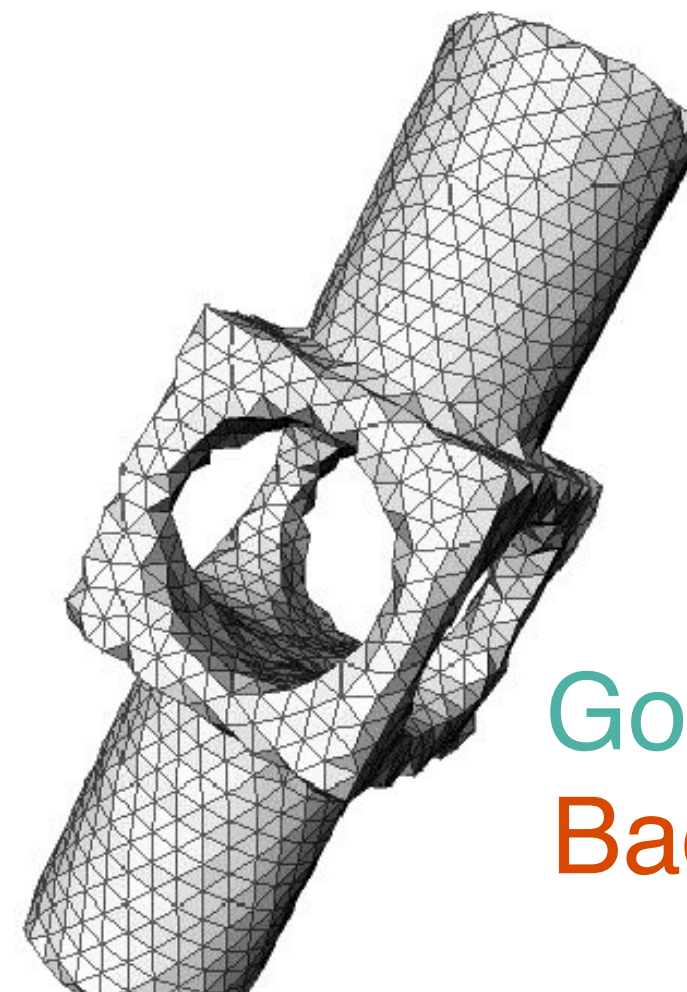
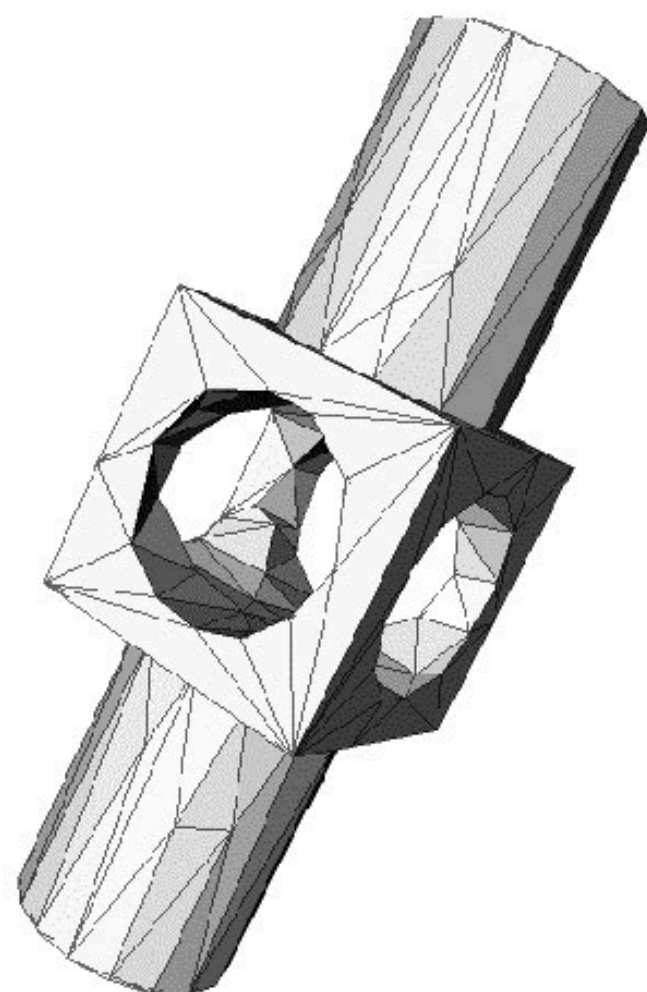


Delaunay



Non-Delaunay

Bad triangle  
Good shape

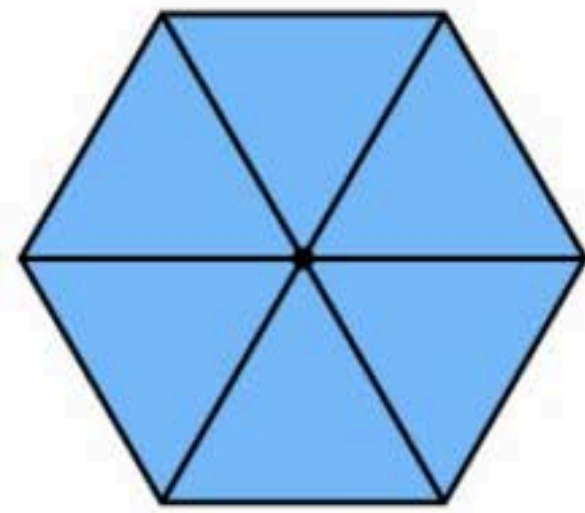


Good triangle  
Bad shape

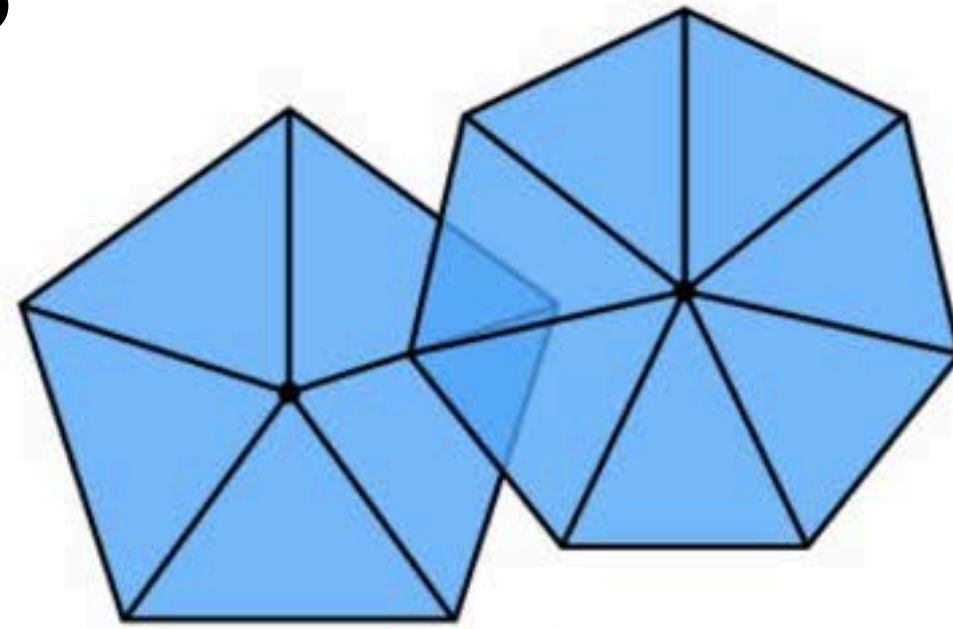


# What is a “good” mesh?

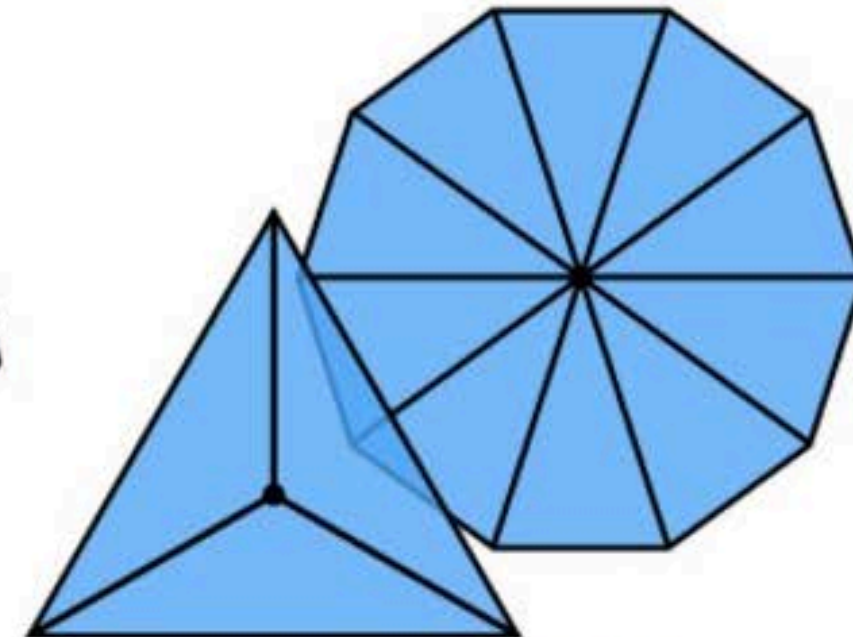
- Another rule of thumb: *regular vertex degree (valence)*
  - ▶ Regular: Vertex degree = 6



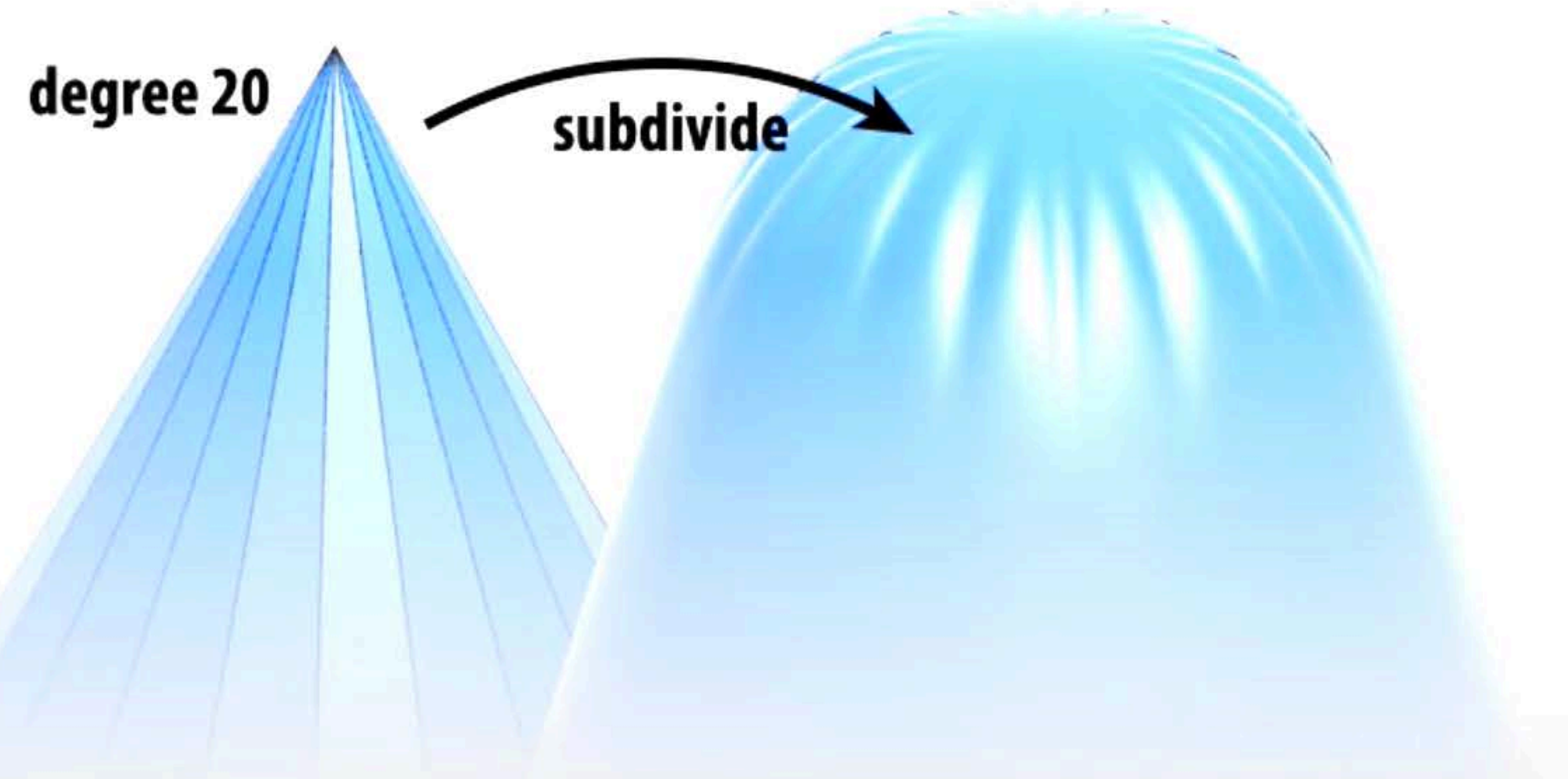
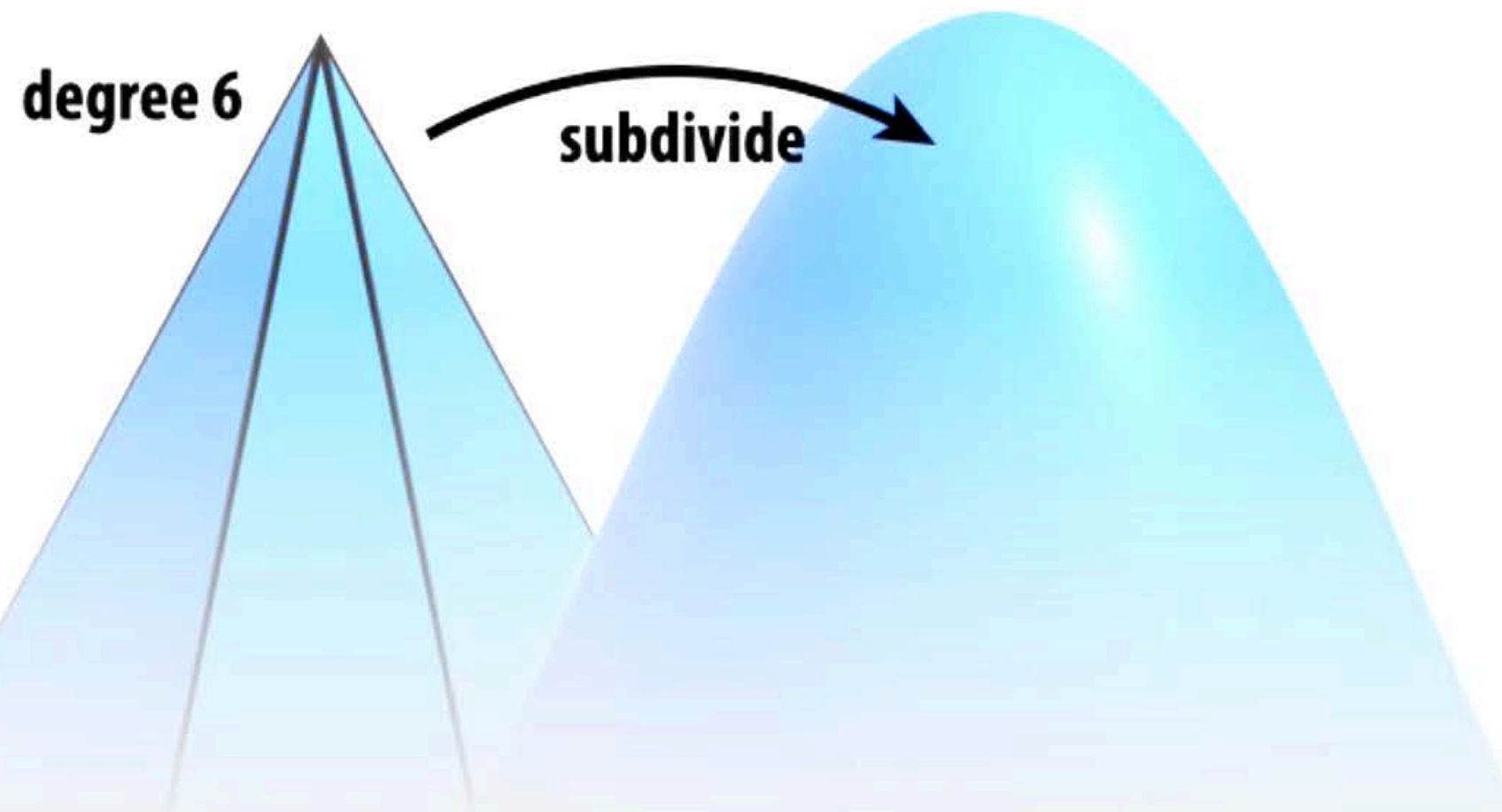
“GOOD”



“OK”



“BAD”

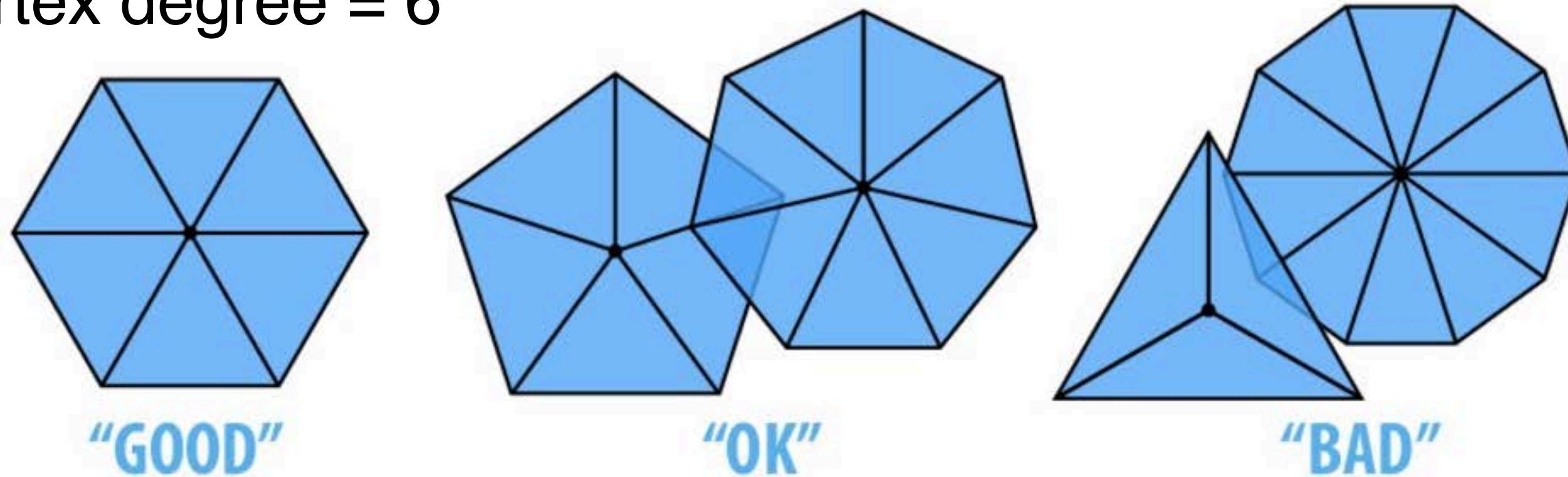




# What is a “good” mesh?

- Another rule of thumb: *regular vertex degree (valence)*

- ▶ Regular: Vertex degree = 6

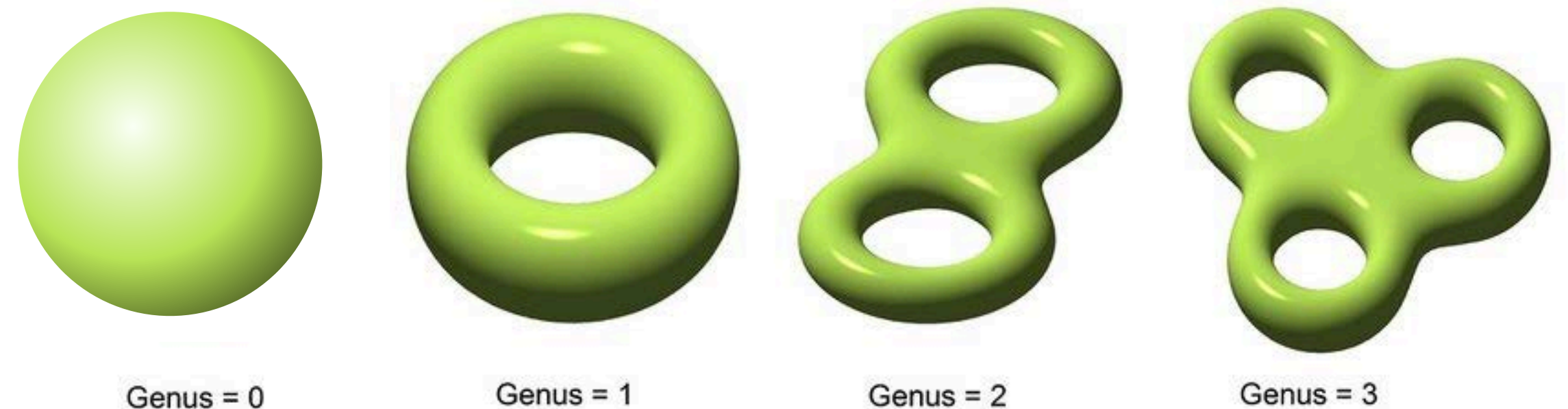


- It may be impossible to have all vertex degree = 6

## Euler–Poincaré Theorem

$$\#(\text{Vertices}) - \#(\text{Edges}) + \#(\text{Faces}) = 2 - 2 \text{ genus}$$

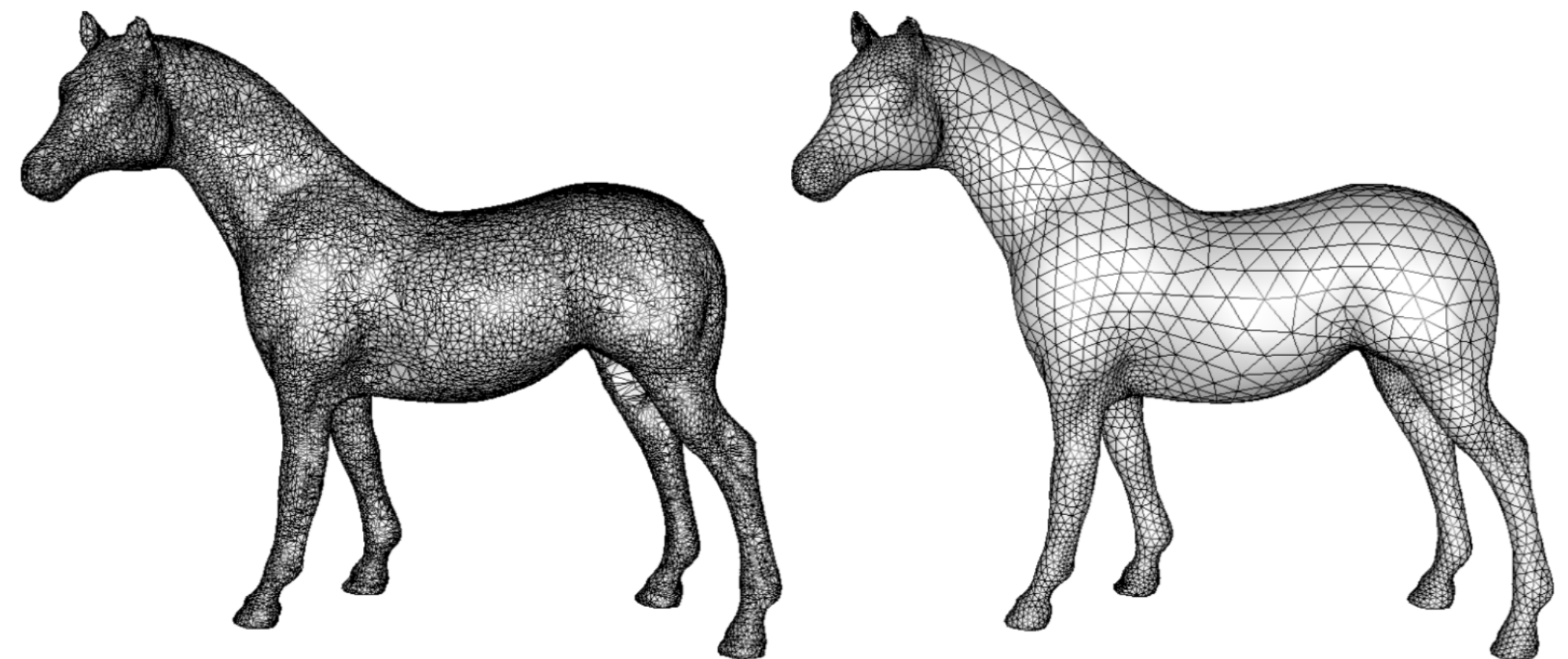
- If all vertices are regular, then genus must be 1





# Remeshing

- General objectives of re-meshing
  - ▶ Shape approximation
  - ▶ As Delaunay (or equilateral-triangle) as possible
  - ▶ Vertex degree as regular as possible
- Mesh simplification
- Improve mesh quality



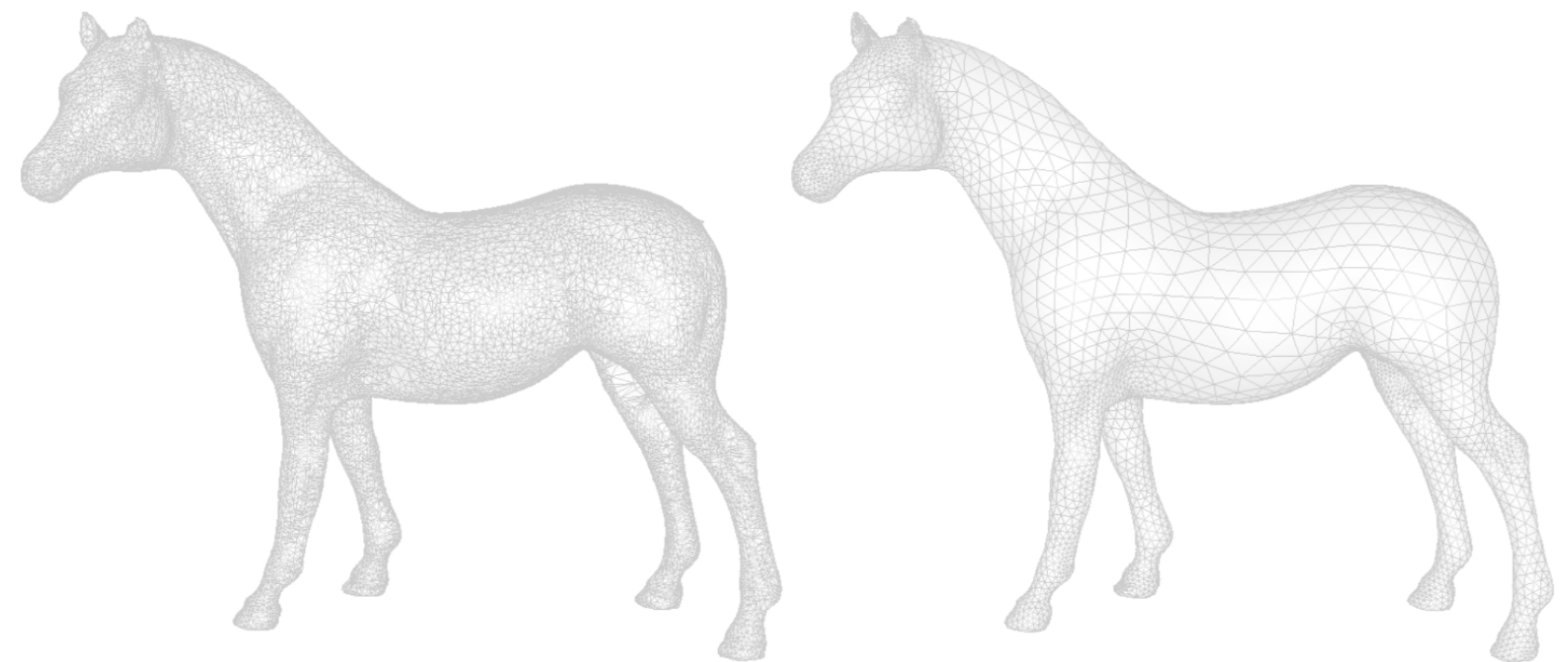


# Remeshing

- Mesh simplification



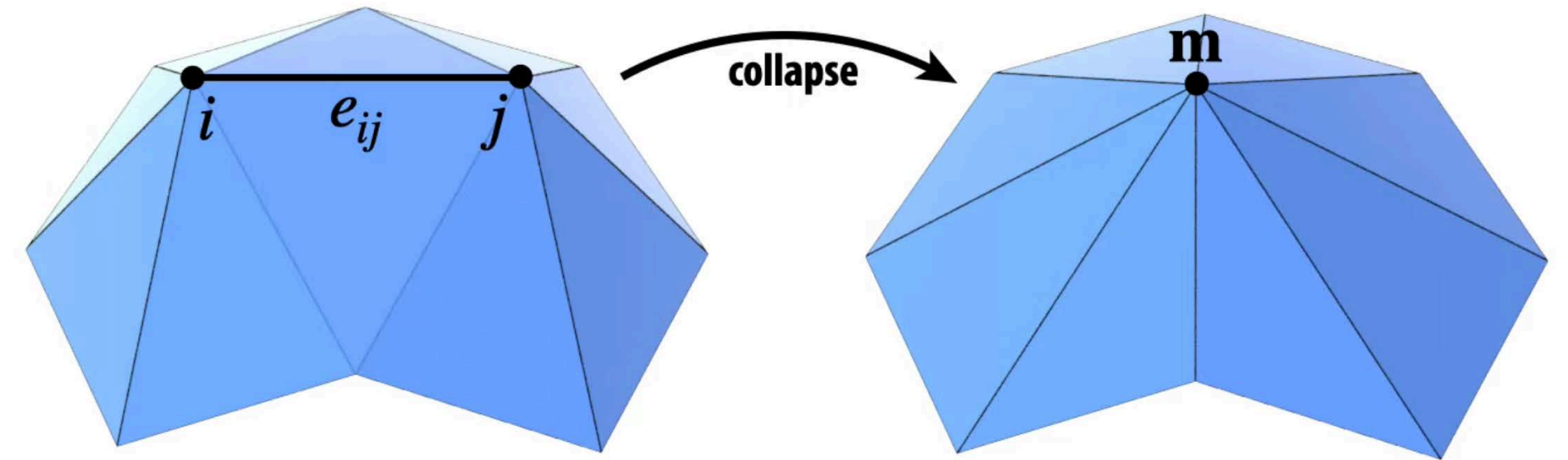
- Improve mesh quality





# Mesh simplification

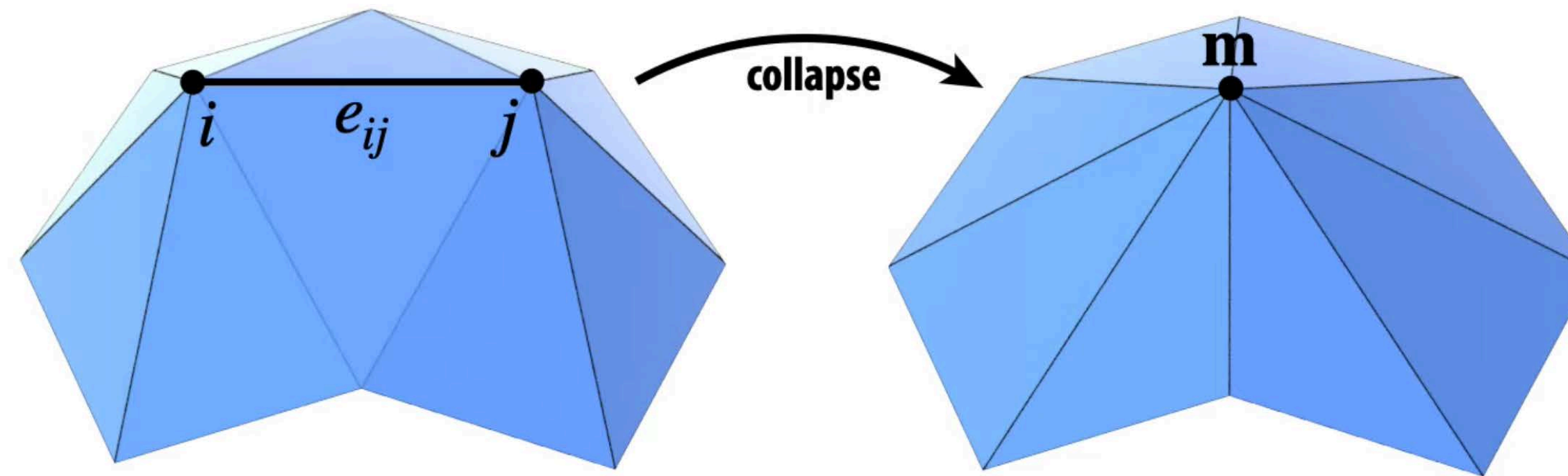
- Popular scheme: Iteratively collapse edges
- Greedy algorithm
  - ▶ Assign each edge a cost
  - ▶ Collapse edge with least cost
  - ▶ Repeat until target number of elements is reached
- Particularly effective cost function: **quadratic error metric**



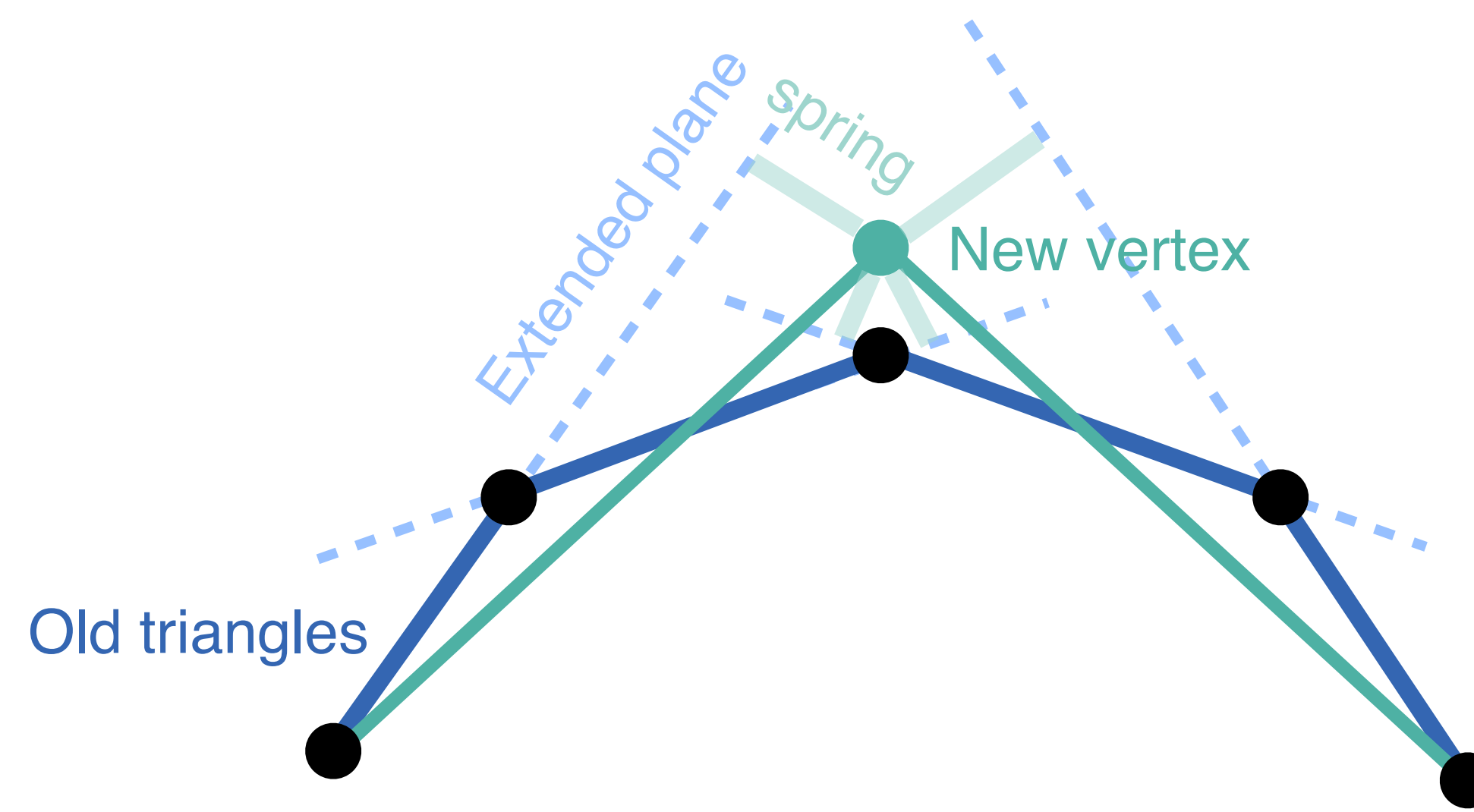


# Mesh simplification

- If we would collapse this edge, where should we set the new vertex?

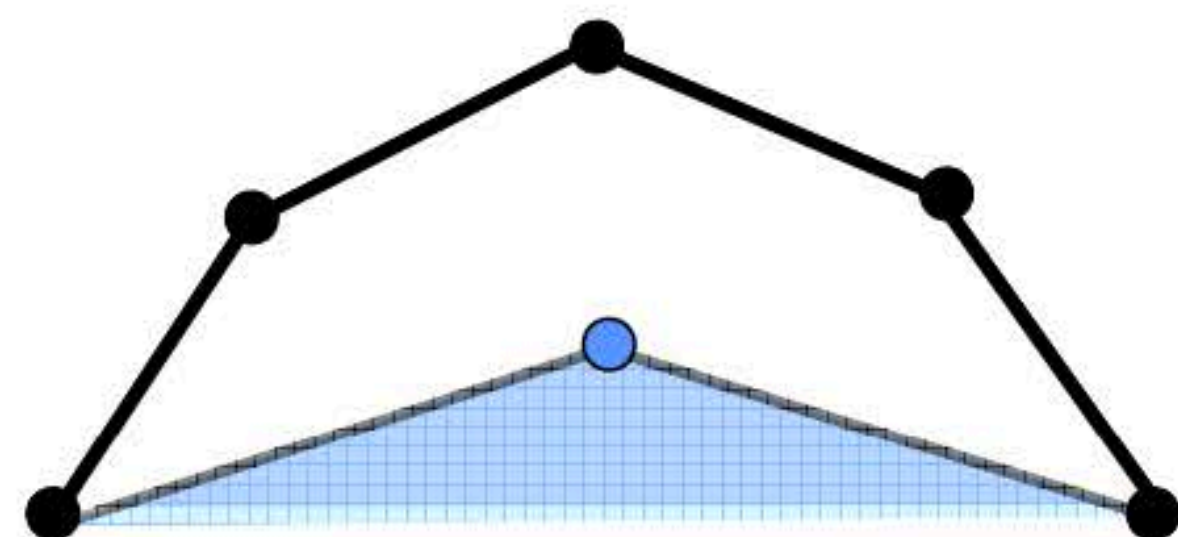


- Minimize the distance-squared to neighboring triangle planes

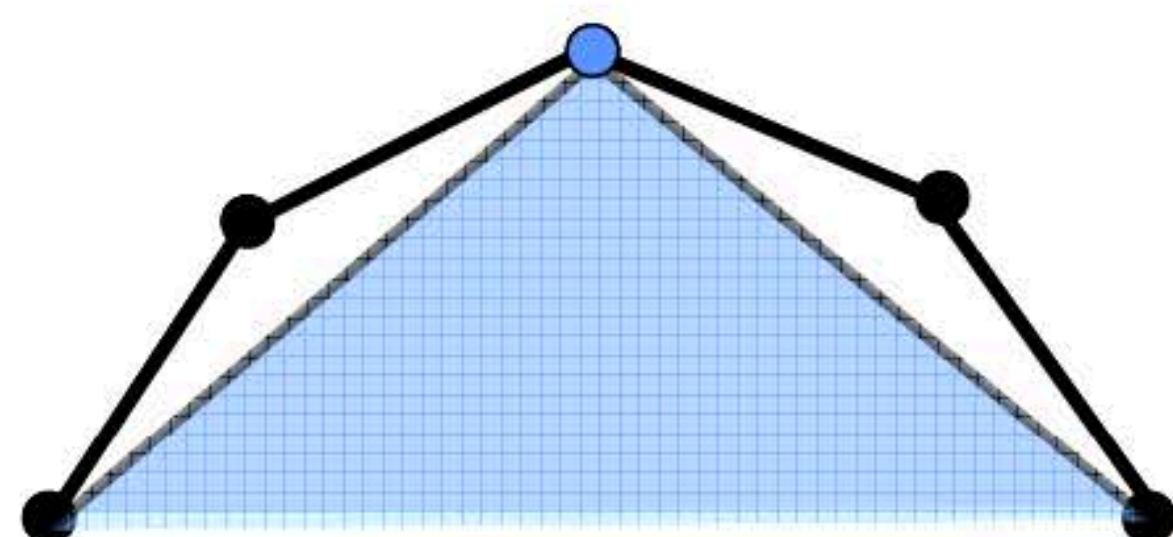




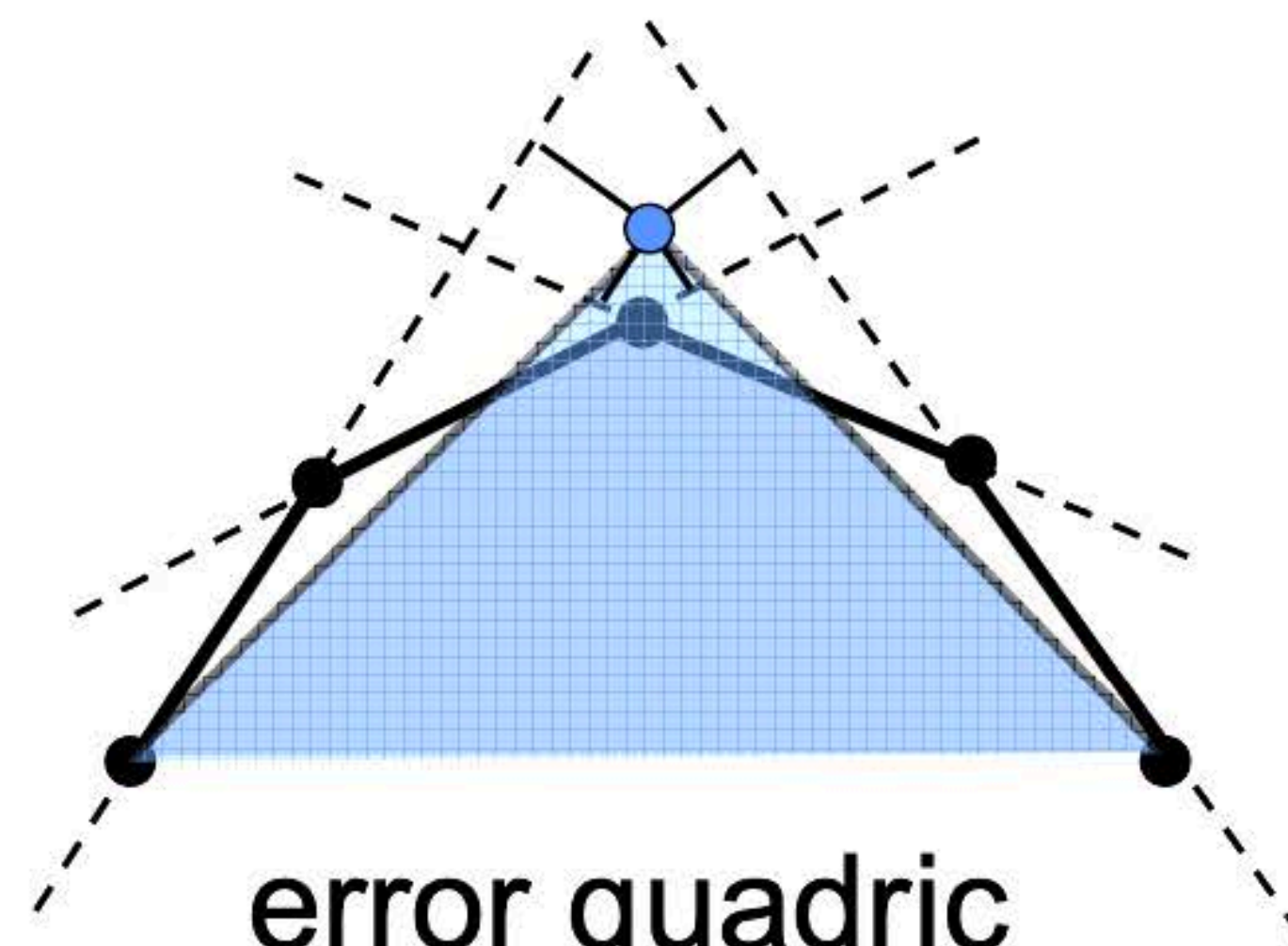
# Mesh simplification



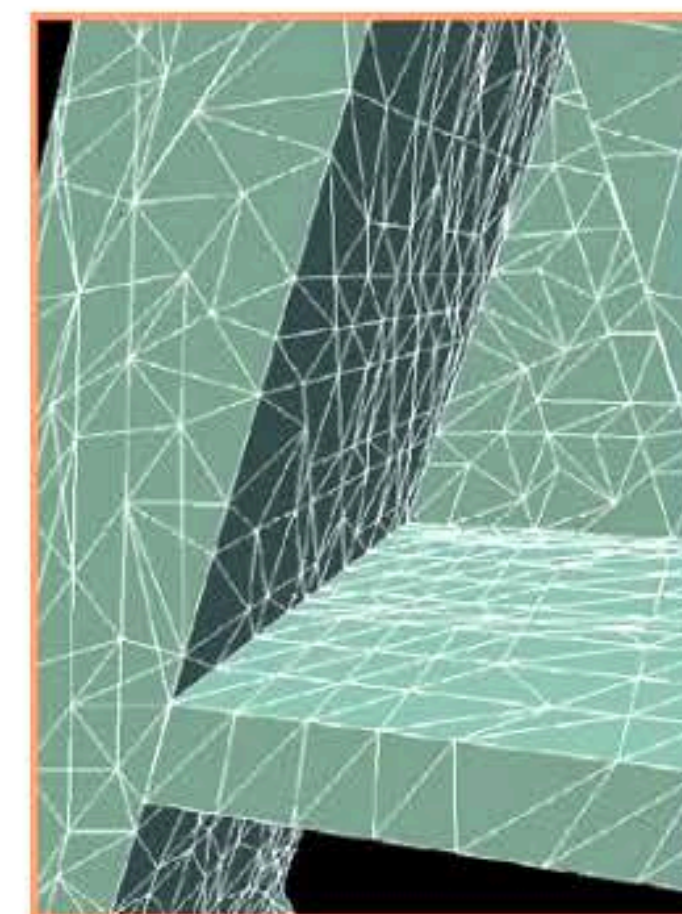
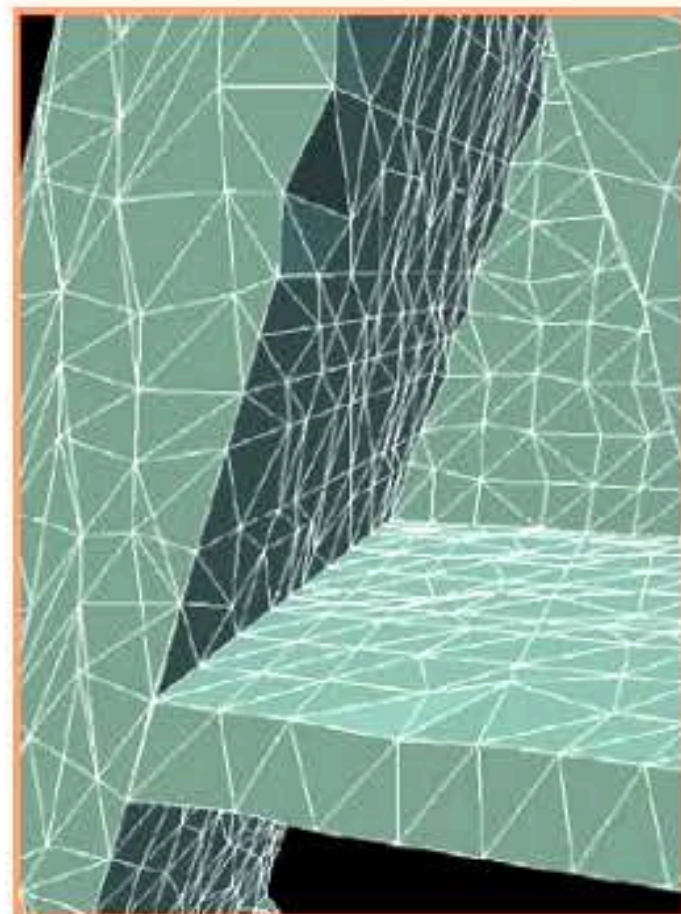
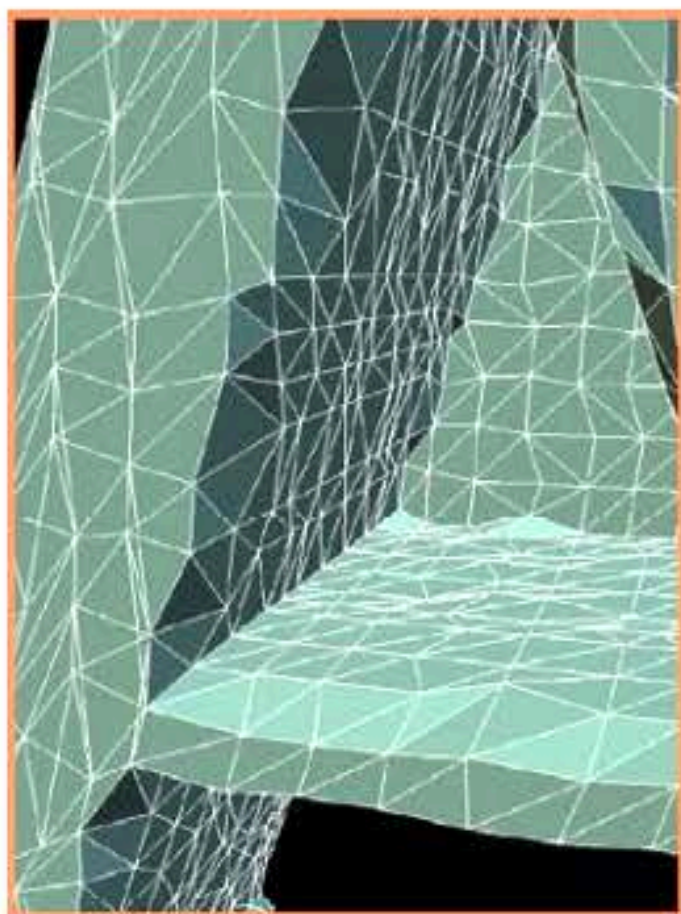
average



median



error quadric

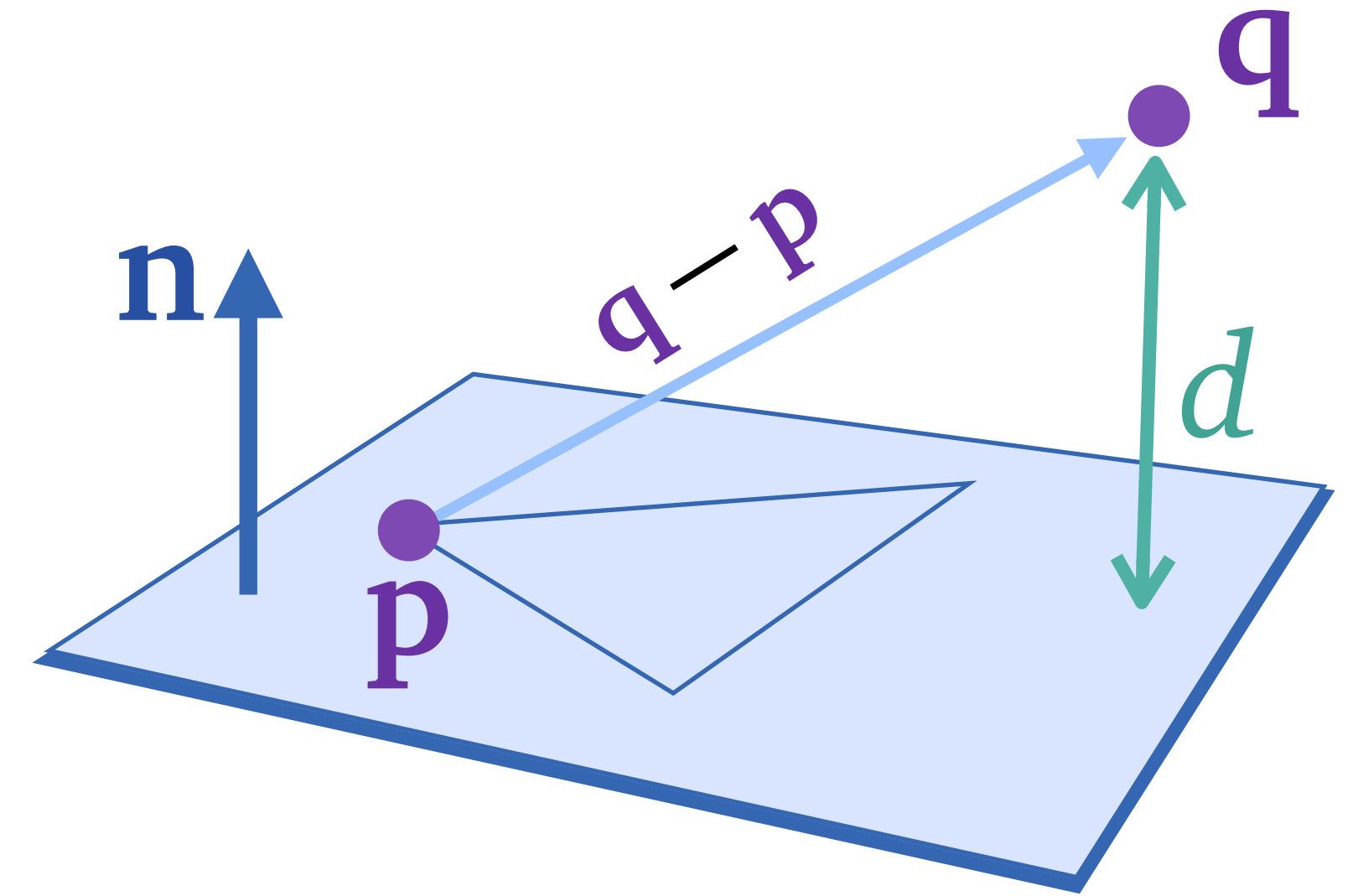




# Mesh simplification

- What is the distance  $d$  between a point  $\mathbf{q} \in \mathbb{R}^3$  and a plane?
- Suppose the plane passes through  $\mathbf{p} \in \mathbb{R}^3$  with unit normal  $\mathbf{n} \in \mathbb{S}^2$

$$\begin{aligned}d &= \mathbf{n} \cdot (\mathbf{q} - \mathbf{p}) \\ &= \begin{bmatrix} - & \mathbf{n}^T & - & (-\mathbf{n} \cdot \mathbf{p}) \end{bmatrix} \begin{bmatrix} | \\ \mathbf{q} \\ | \\ 1 \end{bmatrix} \\ &= \mathbf{a}_{4D}^T \mathbf{q}_{4D}\end{aligned}$$



- Every plane is now a 4D row vector  $\mathbf{a}_{4D}^T$

$$d^2 = \mathbf{q}_{4D}^T (\mathbf{a}_{4D} \mathbf{a}_{4D}^T) \mathbf{q}_{4D}$$



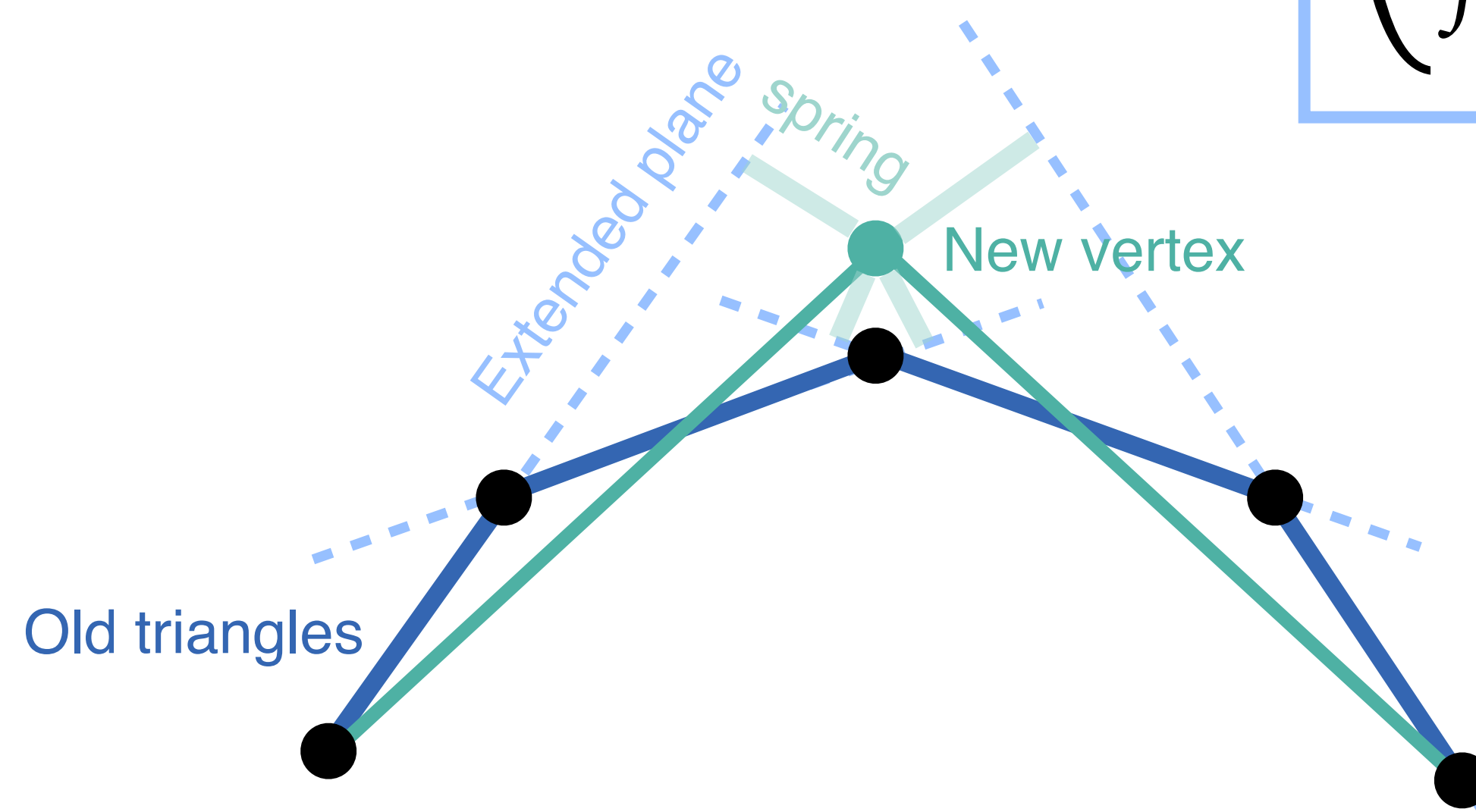
# Mesh simplification

- What is the distance  $d$  between a point  $\mathbf{q} \in \mathbb{R}^3$  and a plane?

$$d^2 = \mathbf{q}_{4D}^T (\mathbf{a}_{4D} \mathbf{a}_{4D}^T) \mathbf{q}_{4D}$$

- The total spring energy  $U(\mathbf{q}) = \frac{1}{2} \mathbf{q}_{4D}^T \left( \sum_{j \in \text{neighboring triangles}} \mathbf{a}_j \mathbf{a}_j^T \right) \mathbf{q}_{4D}$

$$\mathbf{K} \in \mathbb{R}^{4 \times 4}$$



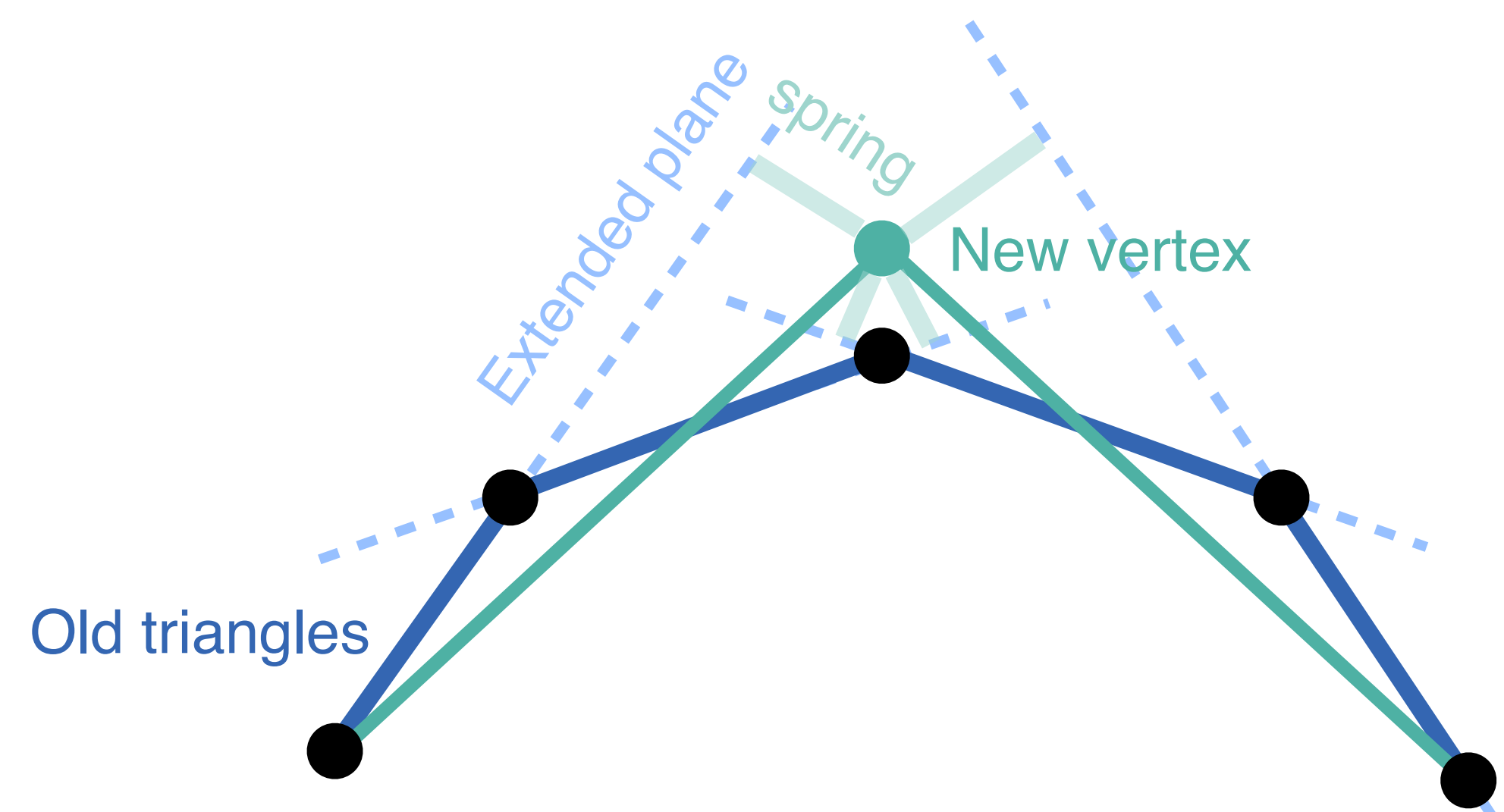


# Mesh simplification

- The position  $\mathbf{q} \in \mathbb{R}^3$  that minimizes the spring energy

$$U(\mathbf{q}) = \frac{1}{2} \begin{bmatrix} \mathbf{q}^\top & 1 \end{bmatrix} \underbrace{\begin{bmatrix} \mathbf{K}_{4 \times 4} \end{bmatrix}}_{\begin{bmatrix} \mathbf{A}_{3 \times 3} & \mathbf{b}_{3 \times 1} \\ \mathbf{b}^\top & c \end{bmatrix}} \begin{bmatrix} \mathbf{q} \\ 1 \end{bmatrix}$$

is the solution to  $\mathbf{A}\mathbf{q} = -\mathbf{b}$





# Mesh simplification

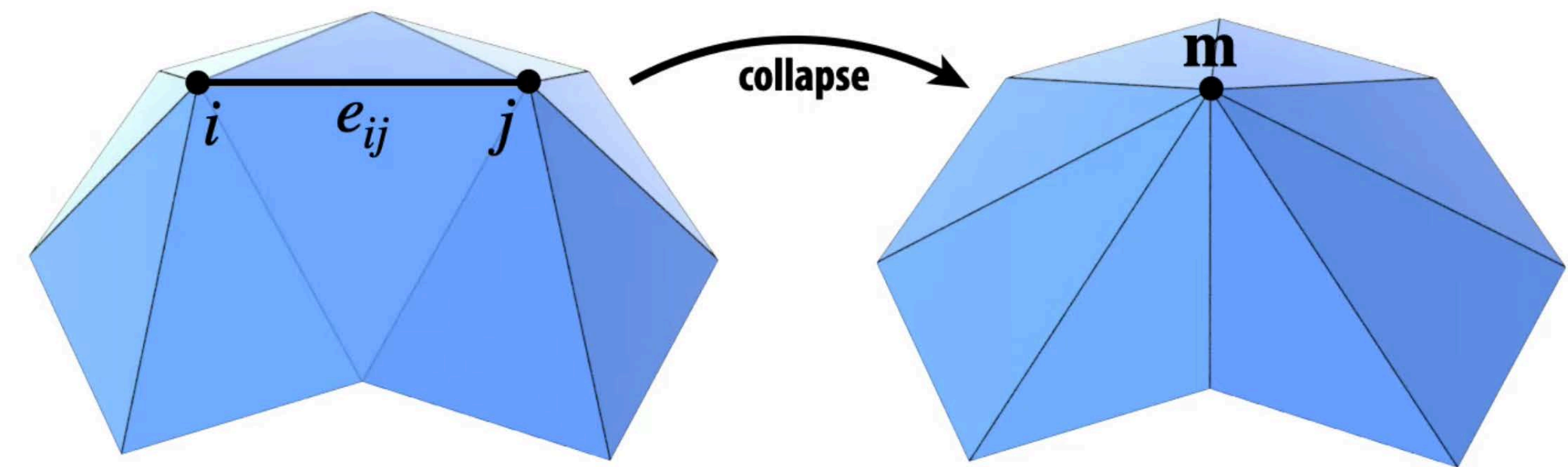
## Algorithm (Assign energy per edge)

- For each edge
  - ▶ Compute the spring matrix  $\mathbf{K}$

$$U(\mathbf{q}) = \frac{1}{2} \mathbf{q}_{4D}^T \left( \sum_{j \in \text{neighboring triangles}} \mathbf{a}_j \mathbf{a}_j^T \right) \mathbf{q}_{4D}$$

$\mathbf{K} \in \mathbb{R}^{4 \times 4}$

$$\mathbf{K} = \begin{bmatrix} \mathbf{A}_{3 \times 3} & \mathbf{b}_{3 \times 1} \\ \mathbf{b}^T & c \end{bmatrix}$$



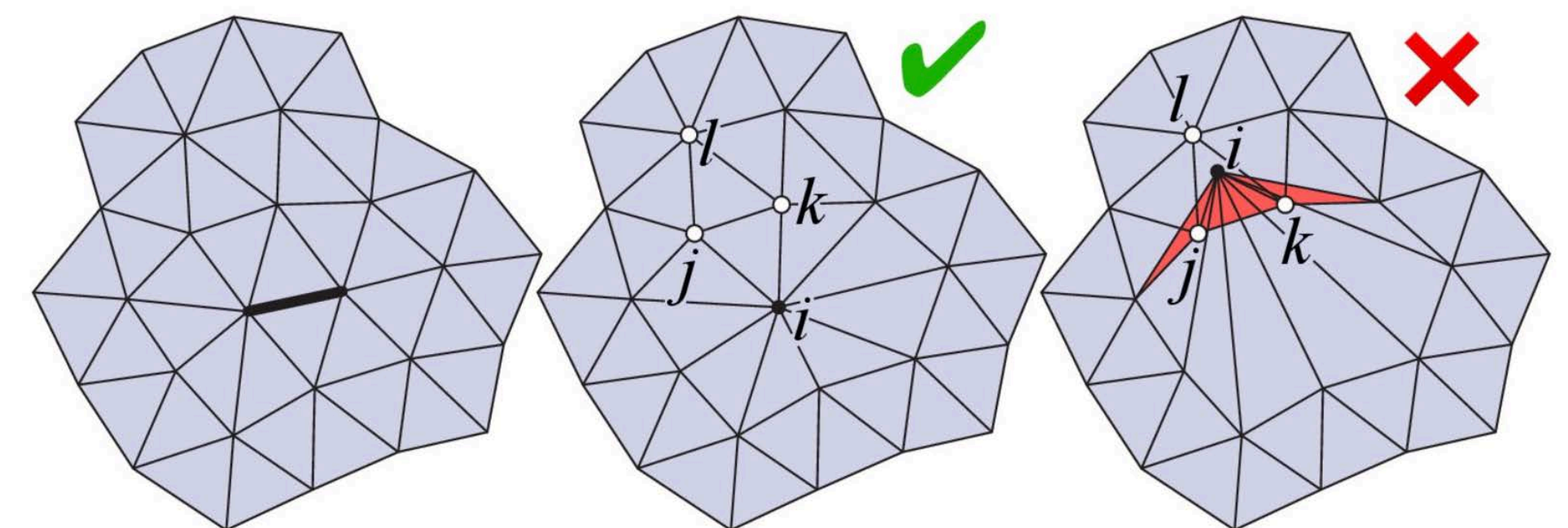
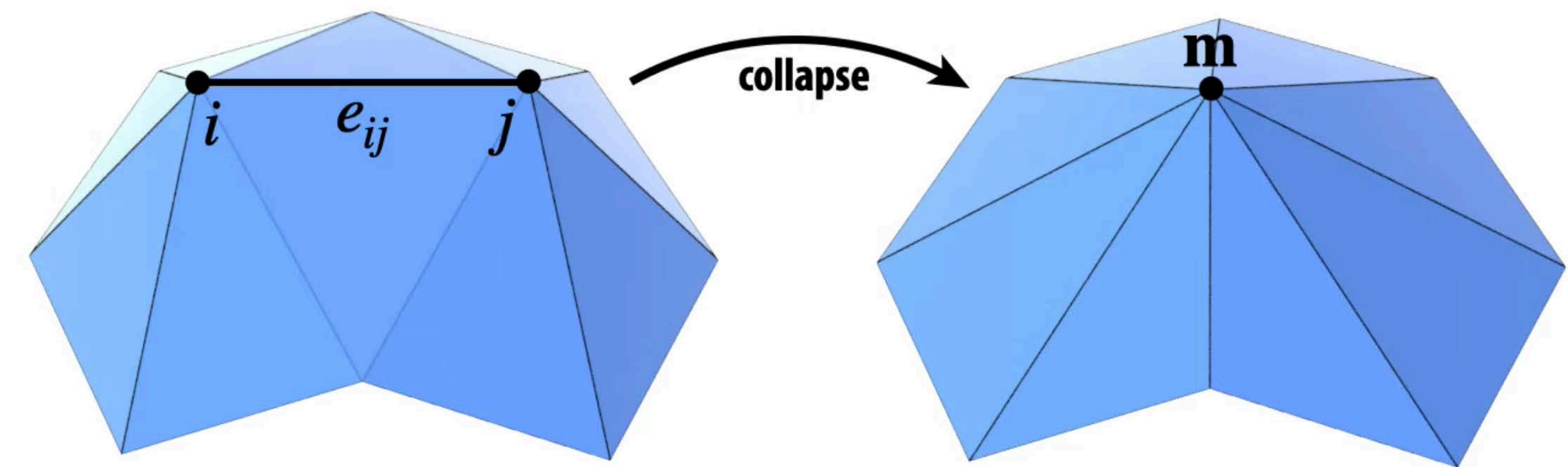
- ▶ Compute the optimal new vertex position  $\mathbf{q} = -\mathbf{A}^{-1} \mathbf{b}$
  - ▶ Assign the edge energy as  $U(\mathbf{q})$
- EndFor



# Mesh simplification

## Algorithm (Edge Collapse)

- ▶ Assign each edge a cost
- ▶ Collapse edge with least cost
  - ▶ Exclude the cases where the triangles fold over
- ▶ Repeat until target number of elements is reached



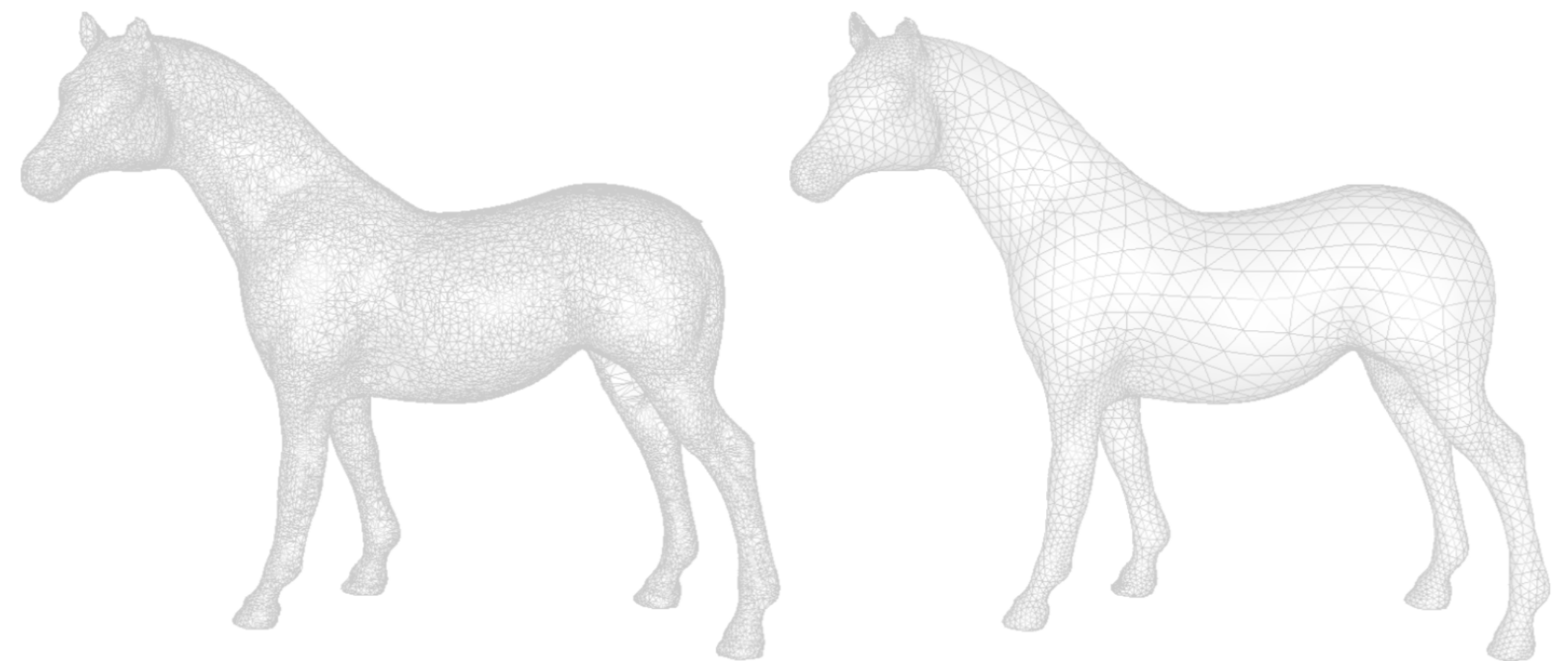


# Remeshing

- Mesh simplification



- Improve mesh quality



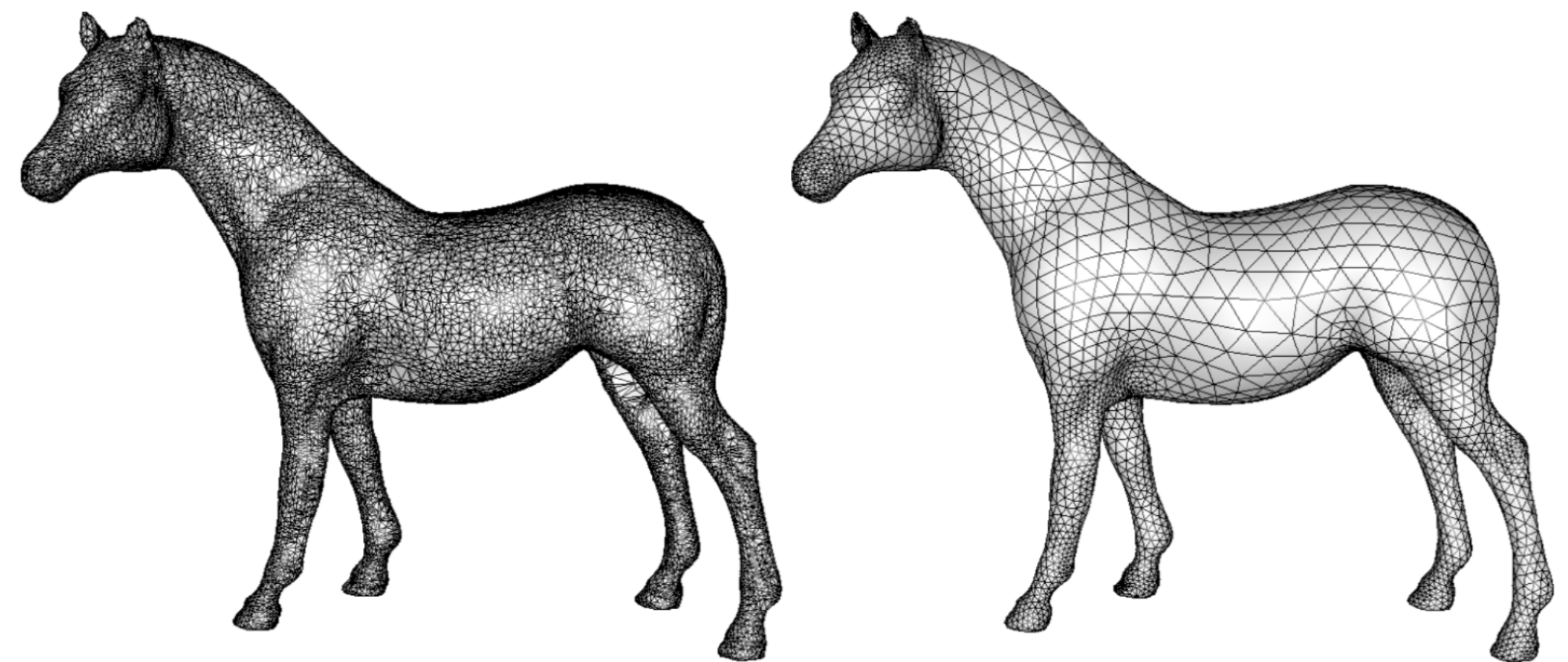


# Remeshing

- Mesh simplification



- Improve mesh quality

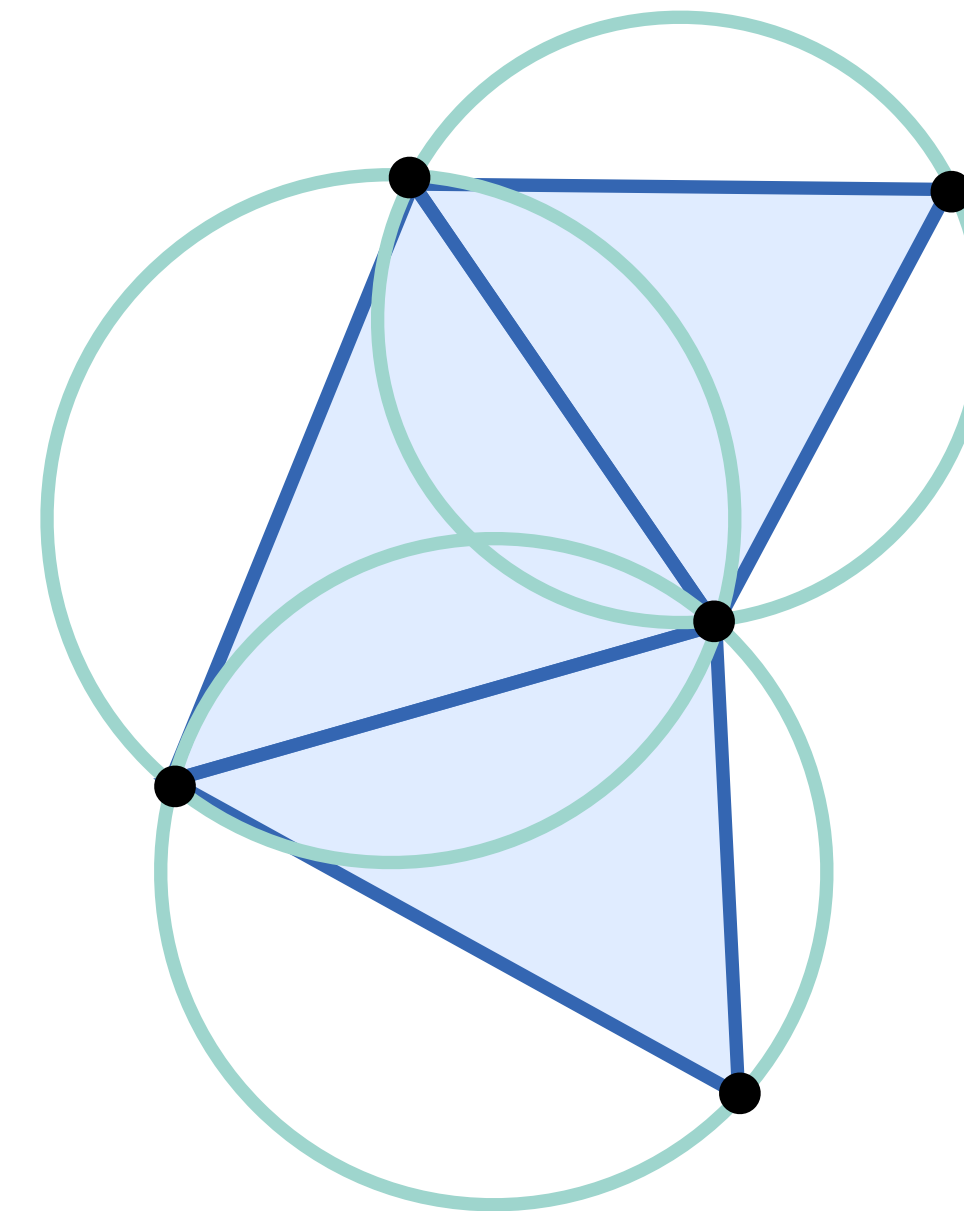




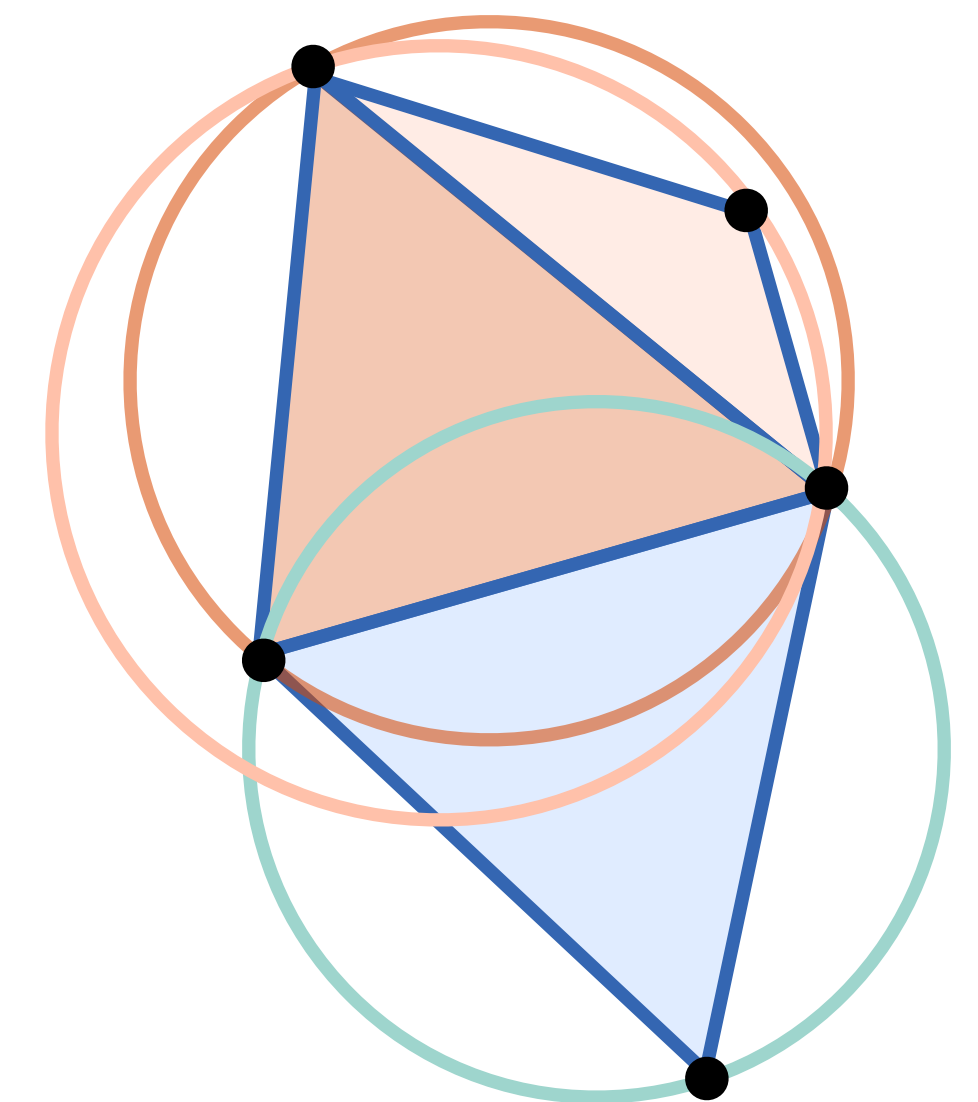
# Make it more Delaunay

- A triangle mesh is **Delaunay** if the circumcircle of each triangle does not contain any vertex of any adjacent triangle

- How do we make a mesh “more Delaunay”?



Delaunay

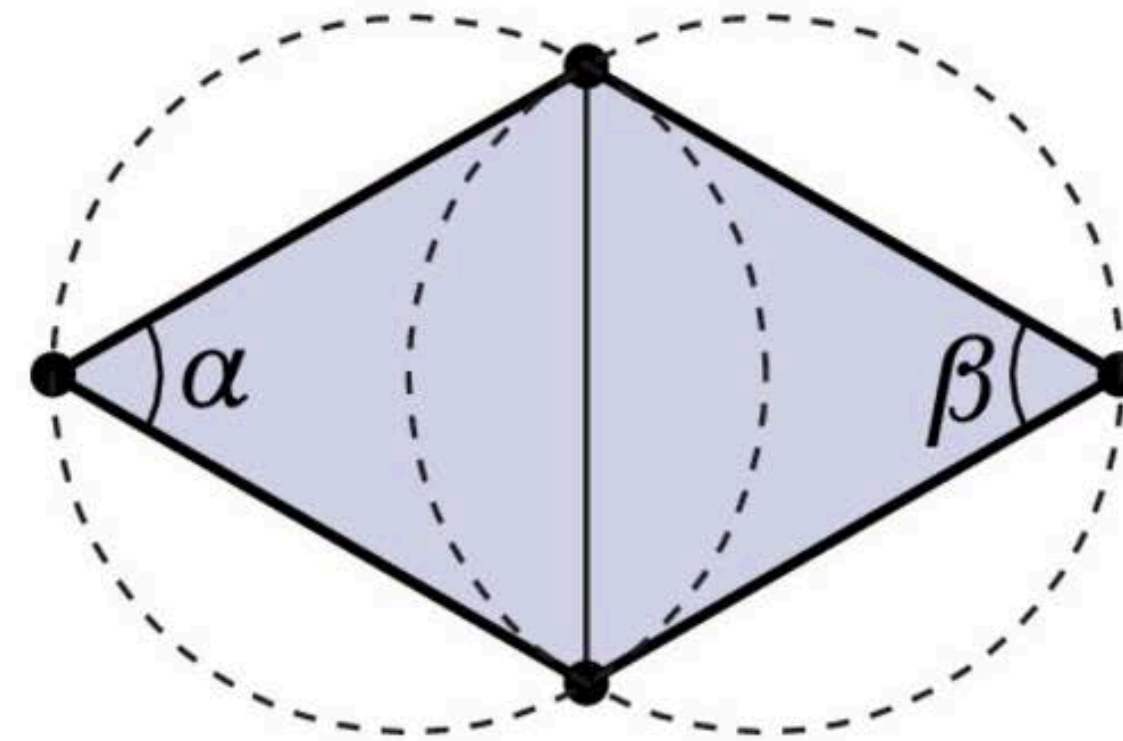
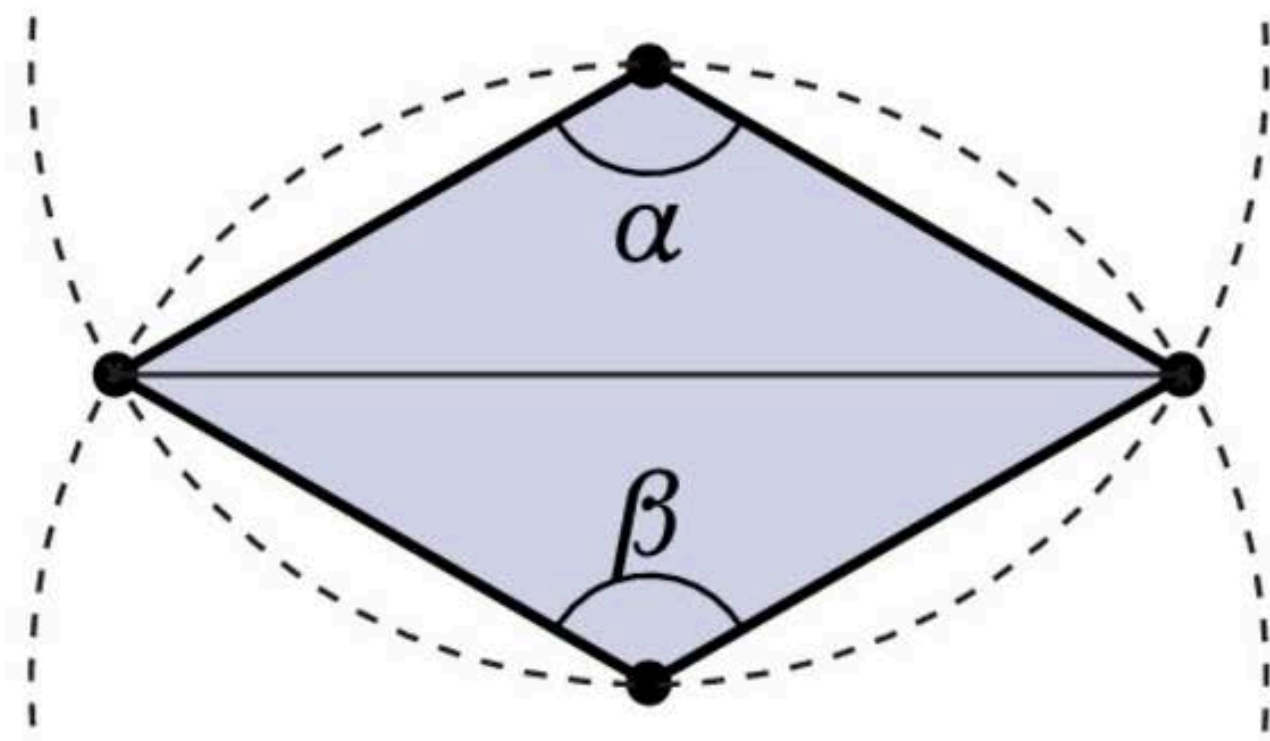


Non-Delaunay



# Make it more Delaunay

- If  $\alpha + \beta > 180^\circ$ , flip the edge.



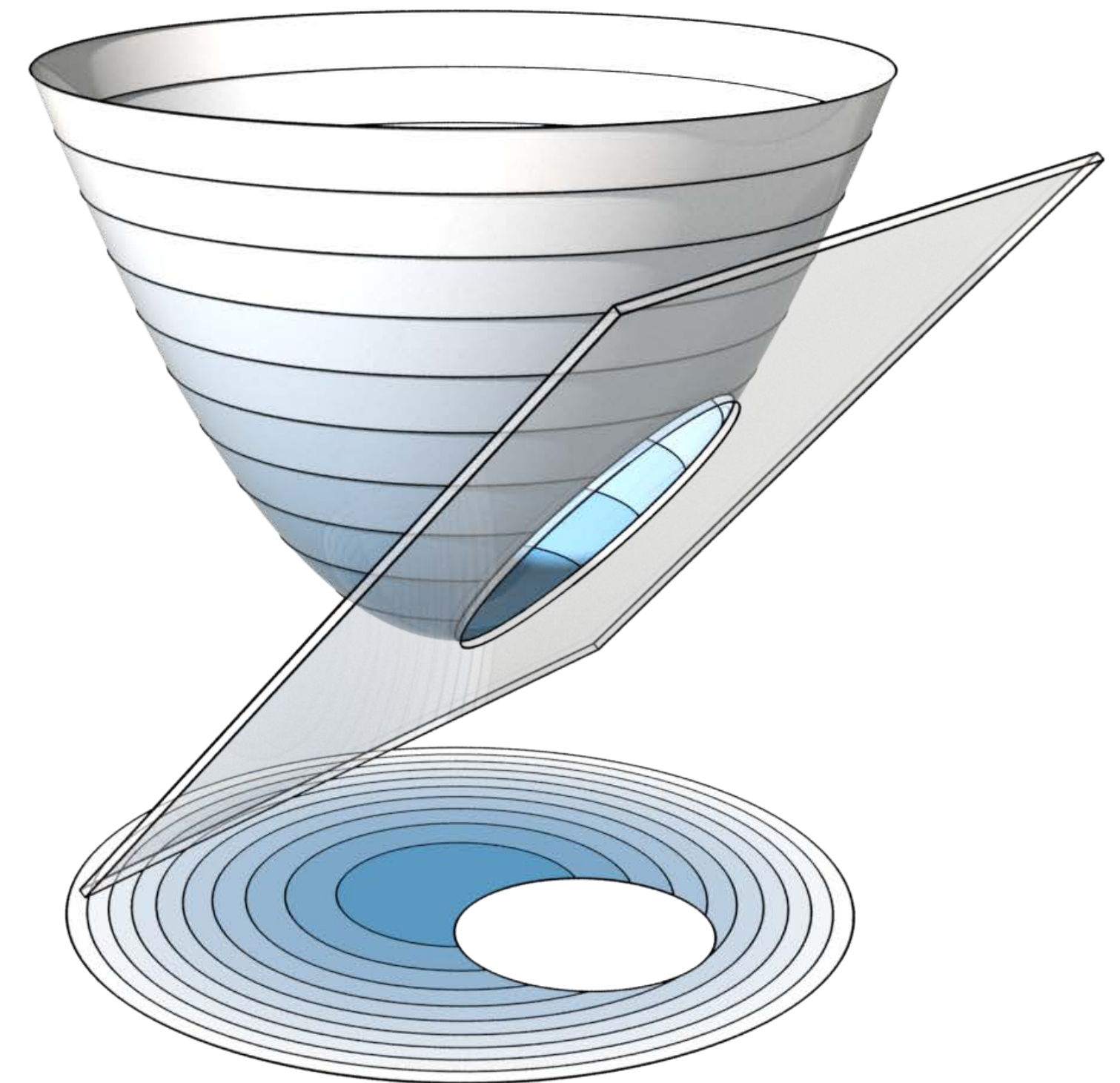
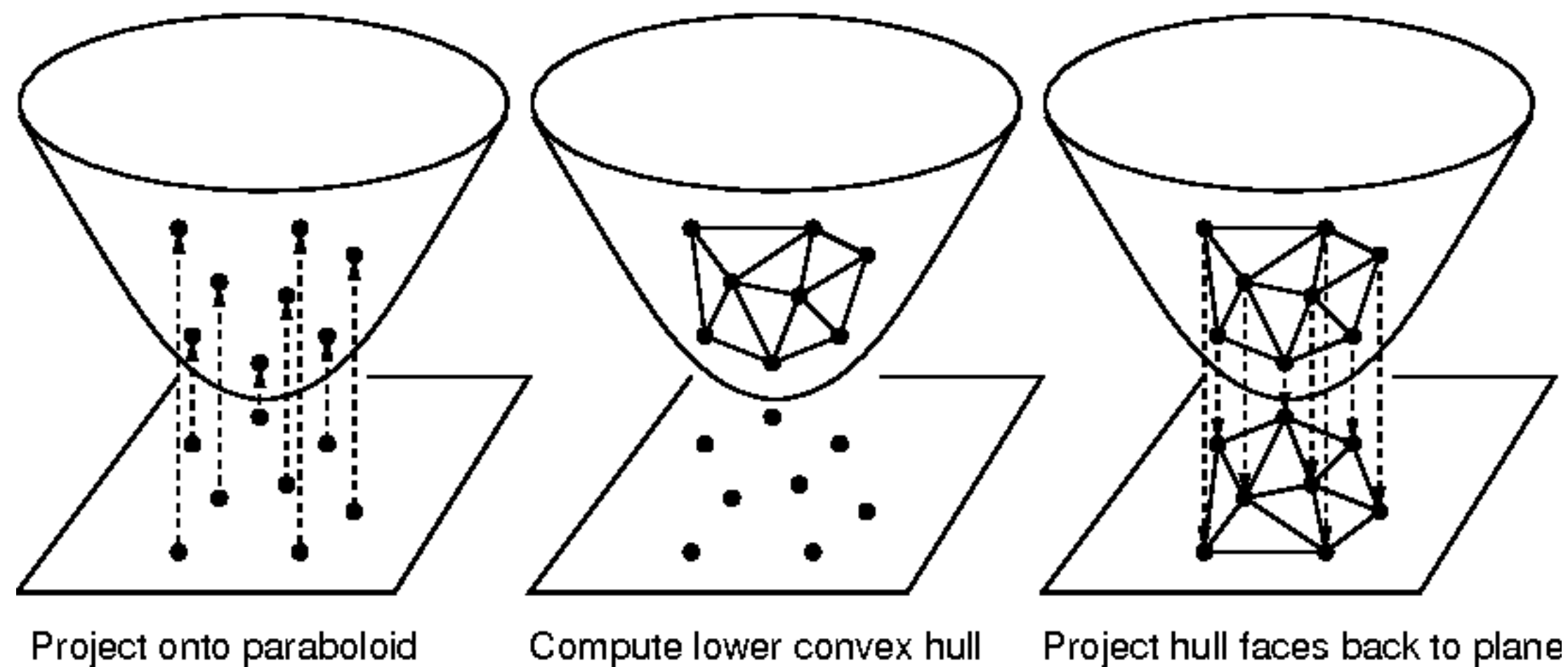
(Delaunay condition is equivalent to  $\alpha + \beta \leq 180^\circ$  for all edges)

- For a planar (2D) mesh, this eventually yields Delaunay mesh
- For surfaces in 3D, this is still good heuristics in practice



# Another algorithm for planar Delaunay

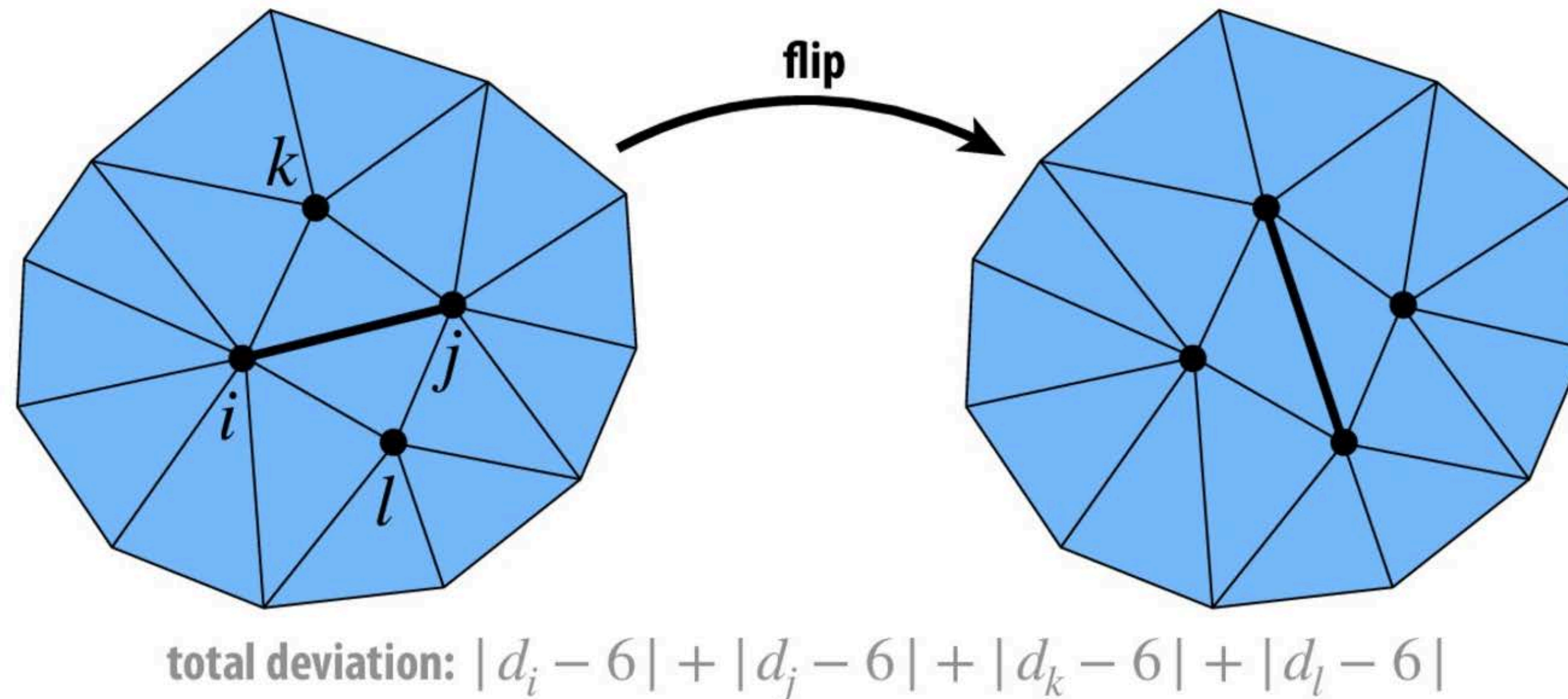
- Notice that every planar section of the paraboloid  $z = x^2 + y^2$  is always circle when viewed from top
- Lift the vertices to the paraboloid



- Delaunay condition in the plane is equivalent to convexity of the lifted triangular mesh

# Alternatively: Improve vertex degree

- How do we improve vertex degree?
- Same tool: Edge flip.
- If total deviation from 6 gets smaller, flip it!

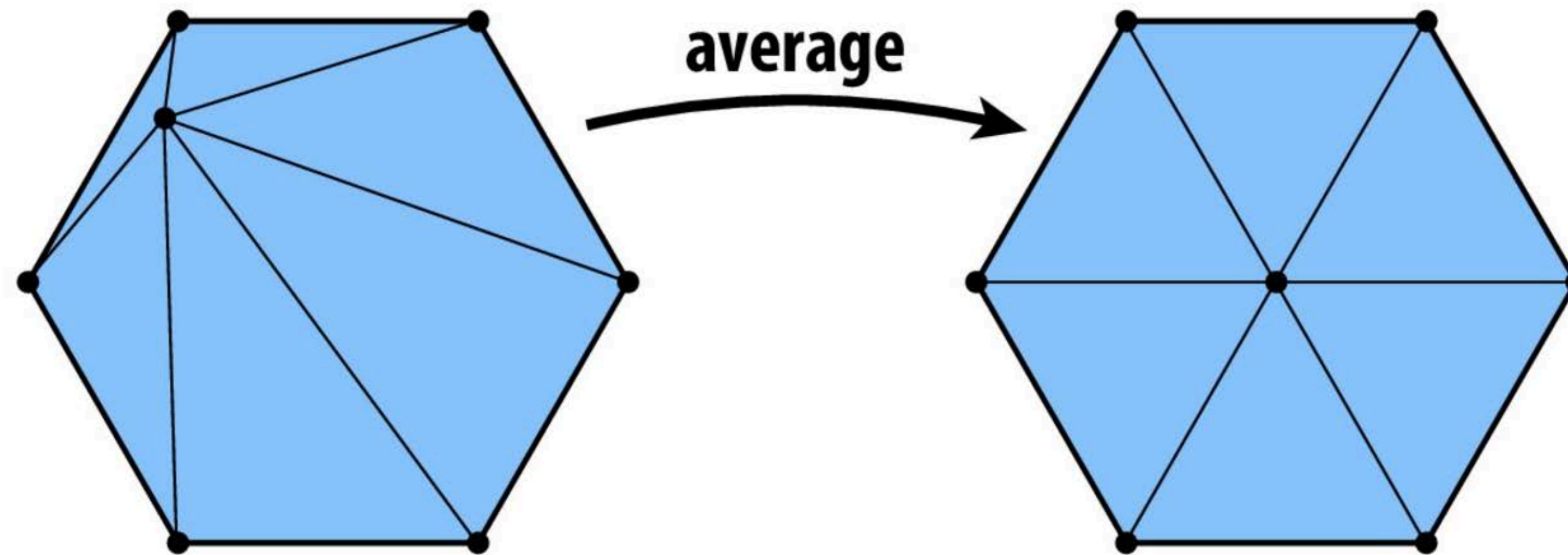


- Works well in practice. No known guarantees.



# Smoothing

- Delaunay doesn't guarantee triangle angles are close to  $60^\circ$
- Improve triangle shapes by centering vertices

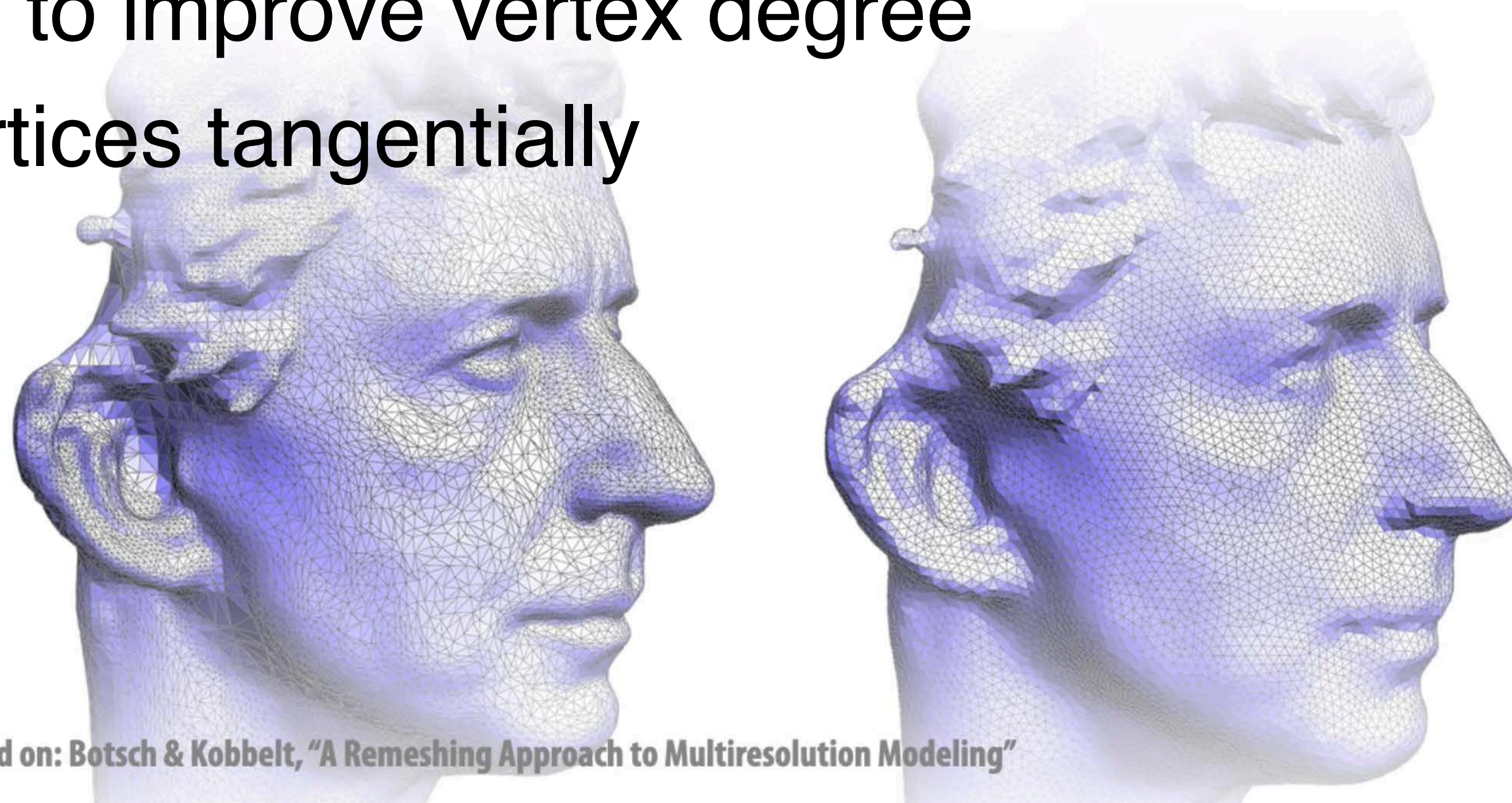


- On surface: move only in the *tangent* direction
  - ▶ How? Remove the normal component of the update vector



# Isotropic Remeshing Algorithm

- Put all tricks together: make triangles uniform in shape & size
- Repeat the 4 steps
  - ▶ Split any edge over  $\frac{4}{3}$  mean edge length
  - ▶ Collapse edge less than  $\frac{4}{5}$  mean edge length
  - ▶ Flip edges to improve vertex degree
  - ▶ Center vertices tangentially





# What we've learned today

- Surfaces are geometric signals
- Basic remeshing

