

CSE203B Convex Optimization

Lecture 2 Convex Set

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Chapter 2 Convex Set

Example: Support Vector Machines

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2. Basic Convex Sets
3. Hyperplanes and Convex Sets
 1. Separating Hyperplanes,
 2. Supporting Hyperplanes
4. Dual Cones

Convex Optimization Problem: Example

Problem Statement

$$\min_x f_0(x), x \in R^n$$

Subject to

$$f_i(x) \leq 0, i = 1, \dots, m$$

$$h_i(x) = 0, i = 1, \dots, p$$

The conditions of convex optimization problem

1. Objective Function, $f_0(x)$, is a convex function

2. Constraint set $S_1 \cap S_2$ is a convex set

$S_1 = \{x | f_i(x) \leq 0, i = 1, \dots, m\}$ is a convex set

$S_2 = \{x | h_i(x) = 0, i = 1, \dots, p\}$ is a convex set

Example: Support Vector Machines

Input: $(x_i, y_i), i = 1, \dots, m; x_i \in R^n, y_i \in \{1, -1\}$

Hard Margin: Find $(a, b), a \in R^n, b \in R,$

Min $|a|^2$

Subject to

$$y_i(a^T x_i - b) \geq 1, \text{ for all } 1 \leq i \leq m$$

Example: Support Vector Machines

Input: $(x_i, y_i), i = 1, \dots, m; x_i \in R^n, y_i \in \{1, -1\}$

Soft margin 1: Find $(w, b), w \in R^n, b \in R,$

$$\text{Min } \lambda |a|^2 + \frac{1}{m} \sum \max(0, 1 - y_i(a^T x_i - b))$$

Soft 2: Find $(a, b), a \in R^n, b \in R, c_i \in R_+, i = 1, \dots, m$

$$\text{Min } \lambda |a|^2 + \frac{1}{m} \sum c_i$$

such that

$$y_i(a^T x_i - b) \geq 1 - c_i, \text{ for all } 1 \leq i \leq m$$

Remark:

1. Support Vector Machine: Find separation hyperplane.
2. Could we derive any better formulations?
3. One key concept is hyperplane, $a^T x = b$.

1. Set Convexity and Specification

1. Convexity Definition
2. Set Specification: Qualification vs. Enumeration Oriented Description

1.1 Set Convexity and Specification: Definition

Definition: A set S is convex if we have

$$\alpha x + \beta y \in S, \forall \alpha + \beta = 1, \alpha, \beta \geq 0, \forall x, y \in S$$

Examples:

1.1 Set Convexity and Specification: Convexity

Definition: A set S is convex if we have

$$\alpha x + \beta y \in S, \forall \alpha + \beta = 1, \alpha, \beta \geq 0, \forall x, y \in S$$

Examples:

Hyperplane $H_1 \{x | a^T x - b = 0\}$, $a, x \in R^n, b \in R$

Half space $H_2 \{x | a^T x - b > 0\}$, $H_3 \{x | a^T x - b < 0\}$

1.1 Set Convexity and Specification: Convexity

Definition: A set S is convex if we have

$$\alpha x + \beta y \in S, \forall \alpha + \beta = 1, \alpha, \beta \geq 0, \forall x, y \in S$$

Remark:

1. Most used sets in the class
 1. Scalar set: $S \subset R$
 2. Vector set: $S \subset R^n$
 3. Matrix set: $S \subset R^{n \times m}$
2. Set S is convex if every two points in S has the connected straight segment in the set.
3. For convex sets S_1 and S_2 : $S_1 \cap S_2$ is also convex

Example:

1.1 Set Convexity and Specification: Convexity

Definition: A set S is convex if we have

$$\alpha x + \beta y \in S, \forall \alpha + \beta = 1, \alpha, \beta \geq 0, \forall x, y \in S$$

Remark:

3. For convex sets S_1 and S_2 : $S_1 \cap S_2$ is also convex

Proof:

Let $U = S_1 \cap S_2$.

By definition $\forall x, y \in U$

$\alpha x + \beta y \in S_1$ because S_1 is convex

$\alpha x + \beta y \in S_2$ because S_2 is convex

Therefore, we have

$\alpha x + \beta y \in U$.

Thus, U is a convex set.

1.2 Set Specification via Qualification or Enumeration

Qualification Oriented Expression $S_Q = \{x | Ax \leq b, x \in R^n\}$

Enumeration Oriented Expression $S_E = \{Ax | x \in R_+^n\}$

Qualification Oriented

Expression:

Constraints

Min $f_o(x)$

Subject to

$$Ax \leq b, x \in R^n$$

Enumeration Oriented

Expression:

Obj function

Min $f_o(Ax), x \in R_+^n$

1.2 Qualification vs Enumeration Oriented Description

Qualification Oriented Expression

Example: $\{x | Ax \leq b\}$

Remark: Simultaneous linear constraints imply **AND**, **Intersection** of the constraints

$$\begin{array}{rclcl} x_1 & +2x_2 & +3x_3 & \leq & 4 \\ 2x_1 & -x_2 & & \leq & 3 \\ & x_2 & +x_3 & \leq & 5 \\ & & x_3 & \leq & 10 \end{array}$$

$$A = \begin{bmatrix} 1 & 2 & 3 \\ 2 & -1 & 0 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \end{bmatrix}, x = \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix}, b = \begin{bmatrix} 4 \\ 3 \\ 5 \\ 10 \end{bmatrix}$$

1.2 Qualification vs. Enumeration Oriented Description

$S_1 = \{x | Ax \leq b, x \in R^n\}$ is a convex set

Proof: Given two vectors $u, v \in S_1$, *i. e.* $Au \leq b, Av \leq b$

For $w = \theta_1 u + \theta_2 v, \forall \theta_1 + \theta_2 = 1, \theta_1, \theta_2 \geq 0$

we have $Aw \leq b$.

($Aw = \theta_1 Au + \theta_2 Av \leq \theta_1 b + \theta_2 b = b$)

The inequality implies $w \in S_1$

By definition, set S_1 is convex.

Remark:

1. Simultaneous linear constraints imply **AND**, **Intersection** of the constraints
2. Linear constraints constitute a convex set.

1.2 Qualification vs. Enumeration Oriented Description

Example:

$$S_2 = \{x \mid Ax \geq b, x \in R^n\}$$

$$S_3 = \{x \mid Ax = b, x \in R^n\}$$

1.2 Qualification Oriented Expression

Example: $S = \{x \in R^m \mid |p_x(t)| \leq 1 \text{ for } |t| \leq \frac{\pi}{3}\}$

where $p_x(t) = x_1 \cos t + x_2 \cos 2t + \cdots + x_m \cos mt$

1.2 Enumeration Oriented Expression

Expression Conversion

Example: $\{x | Ax \leq b, x \in R^n\}$ vs $\{U\theta | 1^T \theta = 1, \theta \in R_+^m\}$

$$\begin{bmatrix} 1 & 1 \\ 1 & 0 \\ -1 & 0 \\ 0 & -1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} \leq \begin{bmatrix} 2 \\ 1 \\ 1 \\ 1 \end{bmatrix} \quad \begin{bmatrix} 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 3 \end{bmatrix} \begin{bmatrix} \theta_1 \\ \theta_2 \\ \theta_3 \\ \theta_4 \end{bmatrix}$$

1.2 Qualification vs. Enumeration Oriented Description

Remark:

Qualification Oriented Expression: Constraints of the problem

Enumeration Oriented Enumeration: The objective function

The interchange may not be trivial.

$$\begin{aligned} \min f_0(x) \\ \text{s. t. } Ax \leq b \\ x \in R^n \end{aligned}$$

$$\begin{aligned} \min f_0(U\theta) \\ \text{s. t. } I^T \theta \leq 1 \\ U \in R^{nm}, \theta \in R_+^m \end{aligned}$$

Every vector u_i in matrix U is a solution of
 n equations in constraint $Ax \leq b$

p equations
 n variables



$\text{comb}(p, n)$ possible
vertex points.

1.2 Qualification vs. Enumeration Oriented Description

Mixed Description

$$S_4 = \left\{ \frac{Ax + b}{c^T x + d} \mid (c^T x + d) > 0, x \in C_4 \right\} \text{ (Projective Function)}$$

$$S_5 = \left\{ \frac{z}{t} \mid z \in R^n, t > 0, (z, t) \in C_5 \right\} \text{ (Perspective Function)}$$

S_4 is convex if C_4 is convex

S_5 is convex if C_5 is convex

1.2 Qualification vs. Enumeration Oriented Description

Statement: S_5 is convex if C_5 is convex.

Proof: Given $\begin{pmatrix} z_1 \\ t_1 \end{pmatrix} \in S_5$, $\begin{pmatrix} z_2 \\ t_2 \end{pmatrix} \in S_5$, let us set

$$z_3 = \alpha z_1 + \beta z_2, t_3 = \alpha t_1 + \beta t_2, \forall \alpha + \beta = 1, \alpha, \beta \geq 0$$

We have
$$\frac{z_3}{t_3} = \frac{\alpha z_1 + \beta z_2}{\alpha t_1 + \beta t_2} = \frac{\alpha t_1}{\alpha t_1 + \beta t_2} \frac{z_1}{t_1} + \frac{\beta t_2}{\alpha t_1 + \beta t_2} \frac{z_2}{t_2}$$

Let $\alpha' = \frac{\alpha t_1}{\alpha t_1 + \beta t_2}$, $\beta' = \frac{\beta t_2}{\alpha t_1 + \beta t_2}$

(Note that $\alpha' + \beta' = 1$, $\alpha', \beta' \geq 0$),

we have
$$\frac{z_3}{t_3} = \alpha' \frac{z_1}{t_1} + \beta' \frac{z_2}{t_2} \in S_5$$

Therefore, by definition S_5 is convex.

2. Basic Convex Sets: Terms and Definitions

Definitions: I. Affine Set, II. Cone, and III. Convex Hull

Given $u_1, u_2, \dots, u_k \in R^n$,

function $f(u, \theta) = \theta_1 u_1 + \theta_2 u_2 + \dots + \theta_k u_k$

and two conditions

1. $\theta_1 + \theta_2 + \dots + \theta_k = 1$
2. $\theta_i \geq 0 \forall i$

I. $\{f(u, \theta) \mid \text{condition 1}\}$: **Affine set**

II. $\{f(u, \theta) \mid \text{condition 2}\}$: **Cone**

III. $\{f(u, \theta) \mid \text{conditions 1 and 2}\}$: **Convex hull**

$$\text{Ex1: } \theta_1 u_1 + \theta_2 u_2 = u_1 + \theta_2 (u_2 - u_1)$$

$$\text{Ex2: } \theta_1 u_1 + \theta_2 u_2 + \theta_3 u_3$$

2. Basic Convex Sets: Terms and Definitions

Definitions: I. Affine Set, II. Cone, and III. Convex Hull

Given $u_1, u_2, \dots, u_k \in R^n$,

function $f(u, \theta) = \theta_1 u_1 + \theta_2 u_2 + \dots + \theta_k u_k$

and two conditions

1. $\theta_1 + \theta_2 + \dots + \theta_k = 1$

I. $\{f(u, \theta) \mid \text{condition 1}\}$: **Affine set**

2. Basic Convex Sets: Terms and Definitions

Definitions: I. Affine Set, II. Cone, and III. Convex Hull

Given $u_1, u_2, \dots, u_k \in R^n$,

function $f(u, \theta) = \theta_1 u_1 + \theta_2 u_2 + \dots + \theta_k u_k$

and two conditions 2. $\theta_i \geq 0 \forall i$

II. $\{f(u, \theta) \mid \text{condition 2}\}$: **Cone**

2. Basic Convex Sets: Terms and Definitions

Definitions: I. Affine Set, II. Cone, and III. Convex Hull

Given $u_1, u_2, \dots, u_k \in R^n$,

function $f(u, \theta) = \theta_1 u_1 + \theta_2 u_2 + \dots + \theta_k u_k$

and two conditions

1. $\theta_1 + \theta_2 + \dots + \theta_k = 1$

2. $\theta_i \geq 0 \forall i$

III. $\{f(u, \theta) \mid \text{conditions 1 and 2}\}$: **Convex hull**

2. Basic Convex Sets: Terms and Definitions

Definitions: I. Affine Set, II. Cone, and III. Convex Hull

Given $u_1, u_2, \dots, u_k \in R^n$,

function $f(u, \theta) = \theta_1 u_1 + \theta_2 u_2 + \dots + \theta_k u_k$

and two conditions

1. $\theta_1 + \theta_2 + \dots + \theta_k = 1$
2. $\theta_i \geq 0 \forall i$

I. $\{f(u, \theta) \mid \text{condition 1}\}$: Affine set

II. $\{f(u, \theta) \mid \text{condition 2}\}$: Cone

III. $\{f(u, \theta) \mid \text{conditions 1 and 2}\}$: Convex hull

Affine Hull: The affine hull of a set $C \subset R^n$ is the smallest affine set containing set C .

Cone: A convex set $C \subset R^n$ is a cone (with apex at the origin) if $\forall x \in C$ and $\forall a \in R_+$, $ax \in C$.

Convex Hull: In geometry, the convex hull of a shape is the smallest convex set that contains it. (Wikipedia)

2. Basic Sets and Definitions: VI. Hyperplanes and Half Spaces

Hyperplane $\{x \mid a^T x = b\}, a \in R^n, b \in R$

or $\{x \mid a^T (x - x_0) = 0\}, \text{ for any } x_0 \in R^n, a \in R^n, b \in R$

Half Space $\{x \mid a^T x \leq b\} \quad a \in R^n, b \in R$

or $\{x \mid a^T (x - x_0) \leq 0\}$

Ex: $\{x \mid x_1 + x_2 = 1\}, \text{ or } \{x \mid [1,1] \left(\begin{bmatrix} x_1 \\ x_2 \end{bmatrix} - \begin{bmatrix} 0.5 \\ 0.5 \end{bmatrix} \right) = 0\},$

or $\{x \mid a^T (x - x_0) = 0\}, \quad a^T = [1,1], b = 1, x_0 = [2, -1]$

For many applications, we standardize the expression:

normalized expression: $\frac{a^T}{\|a\|_2} x = \frac{b}{\|a\|_2}$

2. Basic Sets and Definitions: VI. Hyperplanes

Ex : 3 variables

$$\{x | a^T x = b\}, \quad a^T = (1,1,1), \quad b = 6$$

Ex : 4 variables

$$\{x | a^T x = b\}, \quad a^T = (1,1,1,1), \quad b = 6$$

(1) degrees of freedom on variables: $n - 1$ (R^n).

(2) Vector $(x - y)$ is orthogonal to direction a ,

i.e. $a^T (x - y) = 0, \quad \forall x, y$ in the hyperplane

Proof:

Exercise: Conceptually (visually) construct hyperplane.

2. Basic Sets and Definitions: VI. Hyperplanes

Hyperplane : as an Equal potential of cost function

$$\min f_0(x) = c^T x$$

Vector c is the sensitivity or cost of vector $\begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$

$$e.g. f_0(x) = [1, 2] \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$$

$$\frac{\partial f_0(x)}{\partial x_1} = 1$$

$$\frac{\partial f_0(x)}{\partial x_2} = 2$$

Vector $c = \begin{bmatrix} 1 \\ 2 \end{bmatrix}$ is the sensitivity or cost of vector $\begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$

2. Basic Sets and Definitions

VI. Hyperplane : as a linearized constraint

$$a^T x \leq b, x \in R^n$$

$$e.g. [1, 2] \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} \leq 10$$

Remark :

- Hyperplane is one key building block of convex optimization. (theory, algorithms, applications for machine learning, deep learning, ...)
- Each hyperplane separates the space into two half spaces.
- If $n \geq p$, p hyperplanes can separate the space into 2^p disjoint regions.

2. Basic Sets and Definitions

V. Polyhedra (plural) : Poly (many) Hedron (face)

$$P = \{x | Ax \leq b, Cx = d\}$$

$$A = \begin{bmatrix} a_1^T \\ a_2^T \\ \dots \\ a_m^T \end{bmatrix} \quad C = \begin{bmatrix} c_1^T \\ c_2^T \\ \dots \\ c_p^T \end{bmatrix}$$

2. Basic Sets and Definitions

VI. Matrix Space : **Positive Semidefinite Cone**

$$\textcircled{1} S^n = \{X \in R^{n \times n} | X = X^T\} \text{ Symmetric Matrix}$$

$$\textcircled{2} S_+^n = \{X \in S^n | X \succeq 0\} \quad i. e. \quad y^T X y \geq 0, \forall y$$

$$S_{++}^n = \{X \in S^n | X \succ 0\} \quad i. e. \quad y^T X y > 0, \forall y \neq 0$$

$$\text{Ex: } X = \begin{bmatrix} x & y \\ y & z \end{bmatrix} \in S_+^2 \Leftrightarrow x \geq 0, z \geq 0, xz \geq y^2$$

$$[a \ b] X \begin{bmatrix} a \\ b \end{bmatrix} = a^2 x + b^2 z + 2aby \geq 0, \forall a, b \in \mathbb{R}$$

2. Basic Sets and Definitions

$$\text{Ex: } X = \begin{bmatrix} x & y \\ y & z \end{bmatrix} \in S_+^2 \Leftrightarrow x \geq 0, z \geq 0, xz \geq y^2$$

$$[a \ b]X \begin{bmatrix} a \\ b \end{bmatrix} = a^2x + b^2z + 2aby \geq 0, \forall a, b \in \mathbb{R}$$

Proof: Let $R = \begin{bmatrix} 1 & -x^{-1}y \\ 0 & 1 \end{bmatrix}$

$$\begin{aligned} \text{We have } [a \ b]X \begin{bmatrix} a \\ b \end{bmatrix} &= [a \ b]R^{-T}R^T X R R^{-1} \begin{bmatrix} a \\ b \end{bmatrix} \\ &= [a \ b]R^{-T} \begin{bmatrix} x & 0 \\ 0 & z - x^{-1}y^2 \end{bmatrix} R^{-1} \begin{bmatrix} a \\ b \end{bmatrix} \end{aligned}$$

$$\begin{bmatrix} 1 & 0 \\ -x^{-1}y & 1 \end{bmatrix} \begin{bmatrix} x & y \\ y & z \end{bmatrix} \begin{bmatrix} 1 & -x^{-1}y \\ 0 & 1 \end{bmatrix} = \begin{bmatrix} x & 0 \\ 0 & z - x^{-1}y^2 \end{bmatrix}$$

3. Hyperplanes and Convex Sets

1. Separating Hyperplanes
2. Supporting Hyperplanes

3.1 Separating Hyperplane

$\{x | a^T x = b\}$ (Classification, Optimization, Duality)

Theorem : Given two convex sets $C \cap D = \emptyset$ in R^n

$$\exists a \in R^n, b \in R, \text{ s.t. } a^T x \leq b, \forall x \in C$$

$$a^T x \geq b, \forall x \in D$$

$$\text{Actually, } a = d - c, b = \frac{\|d\|_2^2 - \|c\|_2^2}{2}$$

$$\text{i.e. } f(x) = a^T x - b = (d - c)^T \left(x - \frac{d+c}{2}\right)$$

$$\text{For } \text{dist}(C, D) = \inf\{\|u - v\|_2 \mid u \in C, v \in D\}$$

3.1 Separating Hyperplane

Proof: $\forall v \in D, a^T v \geq a^T d$ should be true

By contradiction, suppose that $a^T v < a^T d$

Then we can show that $d + t(v - d)$ is close to c for $t > 0$

Because $\frac{d}{dt} \|d + t(v - d) - c\|_2^2 = 2(d - c)^T (v - d) < 0$

We have $\|d + t(v - d) - c\|_2 < \|d - c\|_2$ for tiny $t > 0$

3.2 Supporting Hyperplane

Given set $C \in R^n$, and a point x_0 on the boundary of C , the hyperplane $\{x | a^T x = a^T x_0\}$ is called supporting hyperplane of C if $a^T x \leq a^T x_0, \forall x \in C$.

Supporting Hyperplane Theorem: For any nonempty convex set C , and a point x_0 on the boundary of C , There exists a support hyperplane to C at x_0 .

Proof: A separating hyperplane that separates interior C and $\{x_0\}$ is a supporting hyperplane.

4. Dual Cones

Definition: Given Cone K (i.e. $K = \{\sum_{i=1}^k \theta_i u_i \mid \theta_i > 0, u_i \in R^n, \forall i\}$)

$$K^* = \{y \mid x^T y \geq 0, \forall x \in K\}$$

Ex: 1. $K = R_+^n : K^* = R_+^n$

2. $K = S_+^n : K^* = S_+^n$

3. $K = \{(x, t) \mid \|x\|_2 \leq t\} : K^* = \{(x, t) \mid \|x\|_2 \leq t\}$

4. $K = \{(x, t) \mid \|x\|_1 \leq t\} : K^* = \{(x, t) \mid \|x\|_\infty \leq t\}$

4. Dual Cones

Show that cone $K = \{(x, t) \mid \|x\|_1 \leq t\}$ has its dual

$$K^* = \{(x, t) \mid \|x\|_\infty \leq t\}$$

Proof :

Statement $x^T u + tv \geq 0, \forall \|x\|_1 \leq t \leftrightarrow \|u\|_\infty \leq v$

L= \Rightarrow R By contradiction, suppose that $\|u\|_\infty > v$

We can find $\exists x$ s.t $\|x\|_1 \leq 1, x^T u > v$

Setting $t=1$, then we have $u^T(-x) + v < 0$.

R= \Rightarrow L Given $\|x\|_1 \leq t, \|u\|_\infty \leq v$

$$u^T \|-x/t\|_1 \leq \|u\|_\infty \leq v$$

Thus, $u^T(-x) \leq vt$

4. Dual Cones

Definition: $x \leq_K y$ if $y - x \in K$

Theorem: $x \leq_K y$ iff $\lambda^T x \leq \lambda^T y, \forall \lambda \in K^*$

Examples

4. Dual Cones

The polyhedral cone $V = \{x | Ax \geq 0\}$ has its dual cone

$$V^* = \{A^T v | v \geq 0\}$$

Proof : By definition

$$V^* = \{y | x^T y \geq 0, \forall x \in V\}$$

$$\text{Thus } V^* = \{y | x^T y \geq 0, \forall Ax \geq 0\}$$

$$\text{Let } y = A^T v, \text{ we have } x^T y = x^T A^T v > 0 \text{ if } v \geq 0$$

$$\text{Ex: } A = \begin{bmatrix} 1 & 2 \\ 1 & -1 \end{bmatrix} \text{ i.e. } x_1 + 2x_2 \geq 0, x_1 - x_2 \geq 0$$

$$A^T = \begin{bmatrix} 1 & 1 \\ 2 & -1 \end{bmatrix} \text{ i.e. } \{\theta_1 \begin{bmatrix} 1 \\ 2 \end{bmatrix} + \theta_2 \begin{bmatrix} 1 \\ -1 \end{bmatrix} | \theta_1, \theta_2 \geq 0\}$$

4. Dual Cones

Remark: $\{x_0 + \Delta x | \Delta x \in K\}$

(1) K cone can be translated to x_0

(2) If $a \in K^*$, then $a^T x \geq 0, \forall x \in K$, i.e. $-ax$ is a supporting hyperplane of cone K

(3) Suppose that the current feasible search region is at corner x_0
and $\{x_0 + \Delta x | \Delta x \in K, \|\Delta x\| < r\}$ is a local feasible region of a convex set

If $\nabla f_0(x_0) \in K^*$, i.e. $\nabla f_0(x_0)^T \Delta x \geq 0, \forall \Delta x \in K$,

Then x_0 is an optimal solution

Summary

- Set specification
 - Hyperplane and formulation
 - Implicit vs. explicit specification
- Convexity
- Duality