

Interpretation: Saddle-point

Claim : Result of II \supseteq Result of I

Given an arbitrary pair $(\tilde{w}, \tilde{z}) \in D$

$$\min_{w \in W} f(w, \tilde{z}) \leq f(\tilde{w}, \tilde{z}) \leq \max_{z \in Z} f(\tilde{w}, z) \quad \forall \tilde{w}, \tilde{z} \in D$$

$$\min_{w \in W} f(w, \tilde{z}) \leq \max_{z \in Z} f(\tilde{w}, z)$$

$$\text{Thus } \max_{z \in Z} \min_{w \in W} f(w, z) \leq \min_{w \in W} \max_{z \in Z} f(w, z)$$

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Interpretation: Saddle-point

$$\text{Example : } f(w, z) \quad w = \begin{matrix} 1 \\ 2 \\ 3 \end{matrix} \begin{bmatrix} 1 & 1 & 1 \\ 2 & 2 & 2 \\ 3 & 3 & 3 \end{bmatrix} \\ z = 1, 2, 3$$

$$\min_{w \in W} f(w, 1) = 1$$

$$\min_{w \in W} f(w, 2) = 1$$

$$\min_{w \in W} f(w, 3) = 1$$

$$\max_{z \in Z} \min_{w \in W} f(w, z) = 1$$

$$\max_{z \in Z} f(1, z) = 1$$

$$\max_{z \in Z} f(2, z) = 2$$

$$\max_{z \in Z} f(3, z) = 3$$

$$\min_{w \in W} \max_{z \in Z} f(w, z) = 1$$

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Interpretation: Saddle-point

Example : $f(w, z)$ $w = \begin{matrix} 1 & 3 & 2 \\ 2 & 1 & 3 \\ 3 & 2 & 1 \end{matrix}$
 $z = 1, 2, 3$

$$\min_{w \in W} f(w, 1) = 1$$

$$\min_{w \in W} f(w, 2) = 1$$

$$\min_{w \in W} f(w, 3) = 1$$

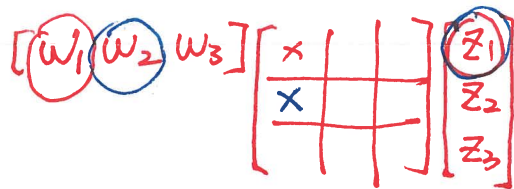
$$\max_{z \in Z} \min_{w \in W} f(w, z) = 1$$

$$\max_{z \in Z} f(1, z) = 3$$

$$\max_{z \in Z} f(2, z) = 3$$

$$\max_{z \in Z} f(3, z) = 3$$

$$\min_{w \in W} \max_{z \in Z} f(w, z) = 3$$



$$\left[\frac{1}{3} \ \frac{1}{3} \ \frac{1}{3} \right] \begin{bmatrix} 1 & 3 & 2 \\ 2 & 1 & 3 \\ 3 & 2 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} 2 & 2 & 2 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} = 2$$

Interpretation: Saddle-point

Example : $f(w, z)$ $w = \begin{matrix} 1 & 4 & 6 & 3 \\ 2 & 2 & 3 & 1 \\ 3 & 3 & 5 & 2 \end{matrix}$
 $z = 1, 2, 3$

$$w=2 \quad f(w, z) = 2$$

$$\min_{w=2} \max_{z=2} f(2, z) = 3$$

$$\max_{z=2} \min_{w=2} f(w, z)$$

$$\min_{w=2} f(w, z) = 3 \quad \begin{matrix} w \neq A \text{ midterm} \\ w \neq B \end{matrix}$$

Proof: Necessity:

Assume that

$$\min_w \max_z f(w, z) = \max_z \min_w f(w, z)$$

$$\text{Let } \tilde{w} = \arg \min_w \max_z f(w, z)$$

$$\tilde{z} = \arg \max_z \min_w f(w, z)$$

We have

$$f(\tilde{w}, \tilde{z}) \underset{\textcircled{1}}{\leq} \max_z f(\tilde{w}, z) = \min_w \underset{\textcircled{3}}{f(w, \tilde{z})} \leq \underset{\textcircled{2}}{f(\tilde{w}, \tilde{z})}$$

By definition

(\tilde{w}, \tilde{z}) is a saddle point.

Sufficiency

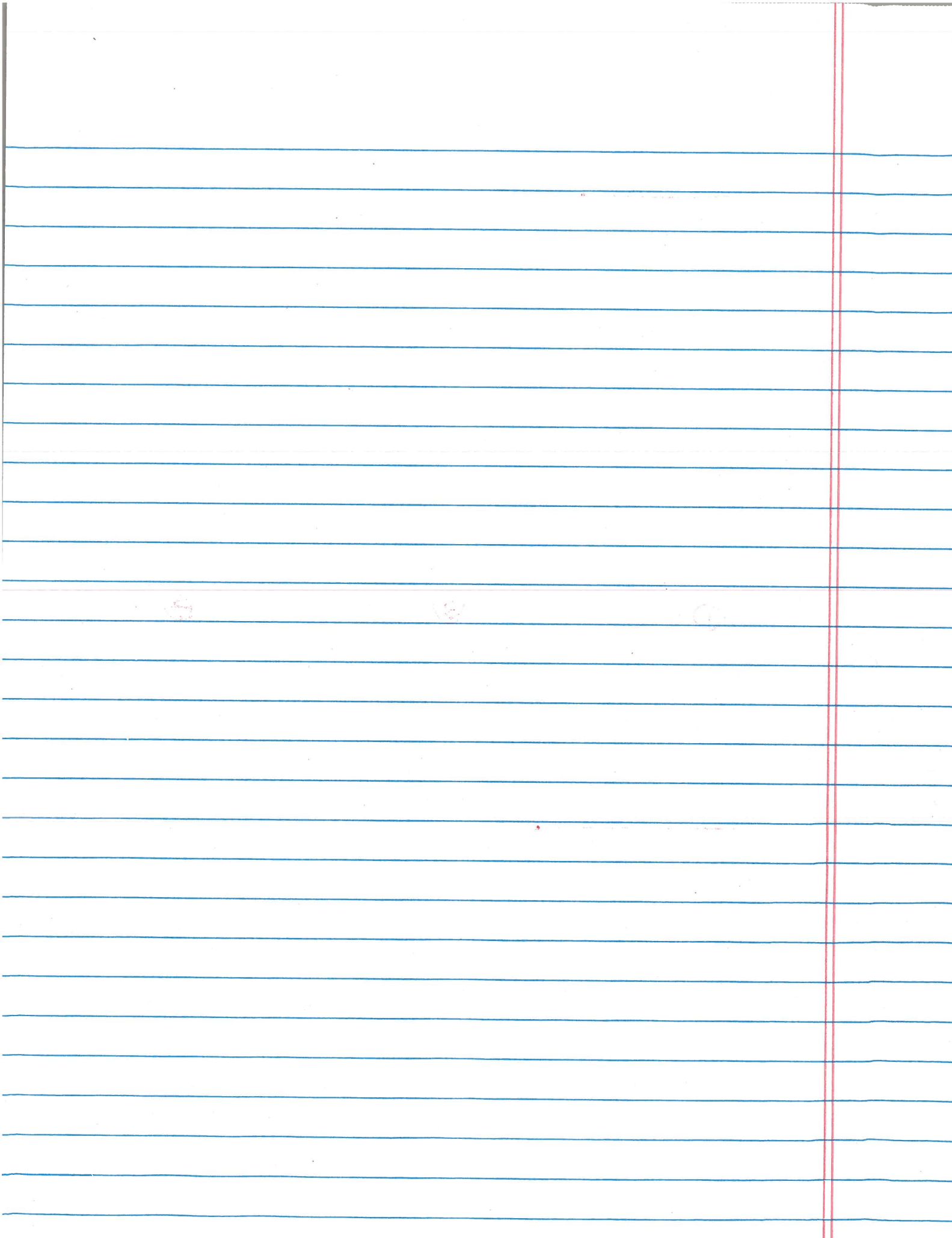
Assume that (\tilde{w}, \tilde{z}) is a saddle point

We have

$$\max_z \min_w f(w, z) \geq \min_w f(w, \tilde{z}) = f(\tilde{w}, \tilde{z})$$

$$\min_w \max_z f(w, z) \leq \max_z f(\tilde{w}, z) = f(\tilde{w}, \tilde{z})$$

$$\text{Thus, } \max_z \min_w f(w, z) = \min_w \max_z f(w, z)$$



Convexity \Rightarrow Saddle Point

Formulation: The row & column selection is formulated as a bilinear optimization problem.

$$\min_{\omega} \max_{z} f(\omega, z) = \sum_i \sum_j a_{ij} \omega_i z_j \quad \Bigg| \quad \max_{z} \min_{\omega} f(\omega, z)$$

I. row & column selection constraints

where $\omega_i, z_j \in \{0, 1\}$ $\sum \omega_i = 1$ & $\sum z_j = 1$

II. relaxed constraints

$$\sum \omega_i = 1 \text{ \& \ } \sum z_j = 1 \quad \omega_i \geq 0, z_j \geq 0. \forall i, j.$$

A. The optimization problem with relaxed constraints can be solved with algorithms Dantzig

$$\min_{\omega} \max_{z} f(\omega, z) = \max_{z} \min_{\omega} f(\omega, z)$$

B. Since $f(\omega, z)$ is convex w.r.t ω
concave w.r.t z .

the solution can reduce to constraint I (row & column selection)

$$f(\tilde{\omega}, \tilde{z})$$

C. From B, $(\tilde{\omega}, \tilde{z})$ is a saddle point.

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Geometric Interpretation

$$\min f_0(x) \quad (t)$$

$$\text{s.t. } f_1(x) \leq 0 \quad (u \leq 0)$$

$$g(\lambda) = \min_{(u,t) \in G} t + \lambda u \quad G = \{(f_1(x), f_0(x)) | x \in D\}$$

$$g(\lambda) = \lambda u + t$$

supporting hyperplane to G

that intersects t axis at $t = g(\lambda)$

u

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Duality via Separating Hyperplane

$$\text{Set } G = \{(f_1(x), \dots, f_m(x), h_1(x), \dots, h_p(x), f_0(x)) | x \in D\},$$

$$G \in R^m \times R^p \times R, p^* = \inf\{t | (u, w, t) \in g, u \leq 0, w = 0\}$$

$$\text{Lagrangian } L = (\lambda, v, 1)^T (u, w, t) = \sum_{i=1}^m \lambda_i u_i + \sum_{i=1}^p v_i w_i + t$$

$$\text{Dual Problem } g(\lambda, v) = \inf_{(u, w, t) \in G} (\lambda, v, 1)^T (u, w, t)$$

u, w, t

Separating hyperplane: Example

$$\{(u, t) | f_0(x) \leq t, f_1(x) \leq u, \exists x \in D\}$$

$$(\tilde{\lambda}, \tilde{v}, \tilde{\mu})^T (u, w, t) \geq \alpha, \quad \forall (u, w, t) \in A$$

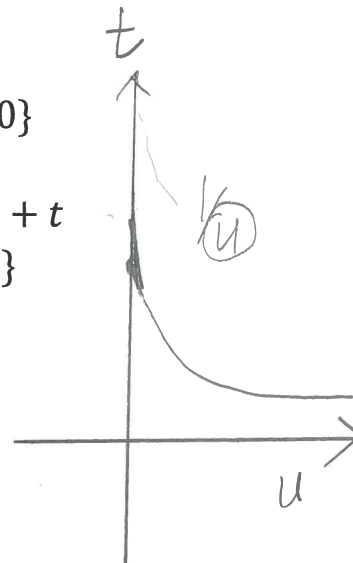
$$(\tilde{\lambda}, \tilde{v}, \tilde{\mu})^T (u, w, t) \leq \alpha, \quad \forall (u, w, t) \in B$$

Since $\tilde{\mu} \neq 0$, we can have $(\lambda, v, 1) = \left(\frac{\tilde{\lambda}}{\tilde{\mu}}, \frac{\tilde{v}}{\tilde{\mu}}, 1\right)$

$$A = \{(u, w, t) | \exists x \in D, f_i(x) \leq u_i, i = 1, \dots, m,$$

$$h_i(x) = w_i, i = 1, \dots, p, f_0(x) \leq t\}$$

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Lagrange dual problem

$$\begin{aligned} \max g(\lambda, v) \\ \text{s.t. } \lambda \geq 0 \end{aligned}$$

$$\min 1/u$$

Properties

This is a convex problem.

The opt. solution is denoted as d^*

$$p^* - d^* = \text{gap} \geq 0$$

If $\text{gap} > 0$, it is a weak duality.

If $\text{gap} = 0$, it is a strong duality.

$$u \leq 0$$

$$u \in \mathbb{R}_+$$

Slater's condition

relint: relative interior of set D

Given that the primal problem is convex,

$$\text{If } f_i(x) < 0, i = 1, \dots, m, \exists x \in \text{relint } D$$

Then strong duality holds.

$$\text{relint } C = \{x \in C \mid B(x, r) \cap \text{aff } C \subseteq C \text{ for some } r > 0\}$$

$$B(x, r) = \{y \mid \|y - x\| \leq r\}$$

any norm

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Shadow Price Interpretation: Food vs. Vitamin

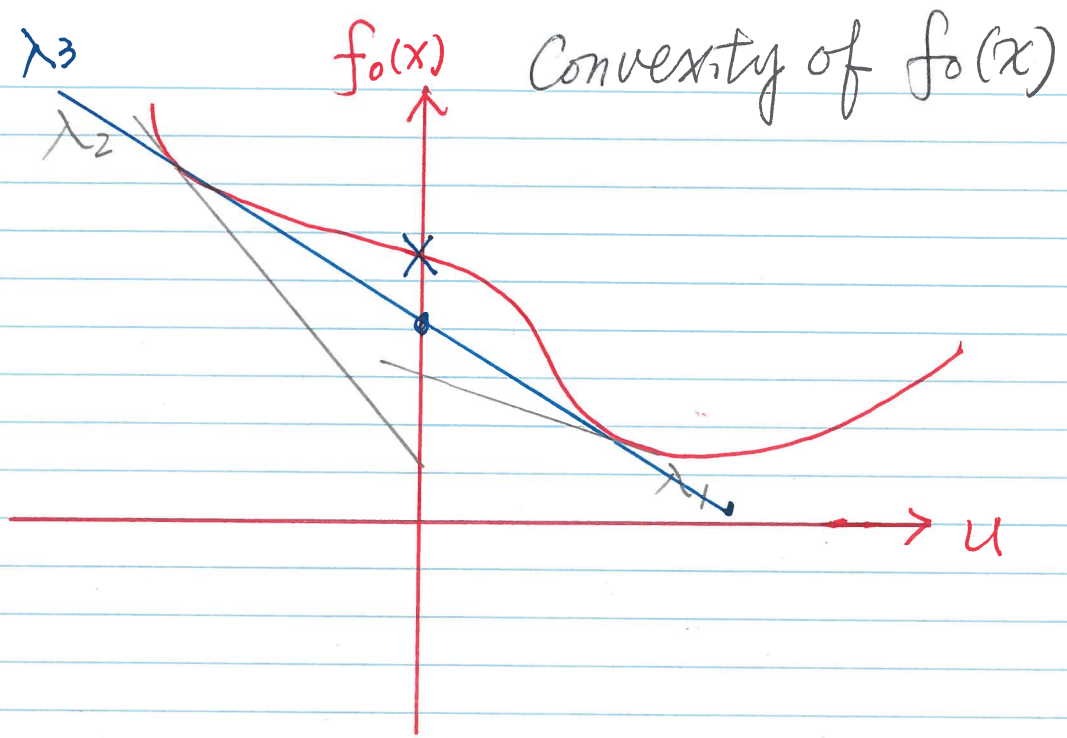
			Flour	protein powder
Primal	$\min c^T x$	$\min c^T x$	Veg.	vitamins A,B,D,E,K
	$\text{s.t. } Ax \geq b$	$\text{s.t. } -Ax + b \leq 0$	Fruits	minerals
	$x \geq 0$	$-x \leq 0$		

$$\min c_1 x_1 + c_2 x_2 + c_3 x_3 \quad \begin{bmatrix} v_{11} & v_{12} & v_{13} \\ v_{21} & v_{22} & v_{23} \\ v_{31} & v_{32} & v_{33} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} \geq \begin{bmatrix} b_1 \\ b_2 \\ b_3 \end{bmatrix}, x_i \geq 0, \forall i$$

$$\begin{aligned} \text{Dual} \quad \max \lambda^T b & \quad \max \lambda_1 b_1 + \lambda_2 b_2 + \lambda_3 b_3 \\ \text{s.t. } A^T \lambda \leq c & \quad \text{s.t. } \begin{bmatrix} v_{11} & v_{21} & v_{31} \\ v_{12} & v_{22} & v_{32} \\ v_{13} & v_{23} & v_{33} \end{bmatrix} \begin{bmatrix} \lambda_1 \\ \lambda_2 \\ \lambda_3 \end{bmatrix} \leq \begin{bmatrix} c_1 \\ c_2 \\ c_3 \end{bmatrix} \\ \lambda \geq 0 & \end{aligned}$$

$$\begin{aligned} \text{Lagrangian} \quad L(x, \lambda) &= c^T x + \lambda_1^T (-Ax + b) + \lambda_2^T (-x) \\ &= [c^T + \lambda_1^T (-A) - \lambda_2^T] x + \lambda_1^T b \\ c^T &= \lambda_1^T (A) + \lambda_2^T, \text{ or } A^T \lambda_1 \leq c \end{aligned}$$

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Primal

$$\min f_0(x)$$

s.t.

$$f_i(x) \leq 0$$

$$h_i(x) = 0$$

$$\min_x \max_{\lambda, \nu} L(x, \lambda, \nu)$$

↑
The constraints are enforced.

Dual

$$\max g(\lambda)$$

$$\min_x L(x, \lambda, \nu)$$

$$\max_{\lambda, \nu} \min_x L(x, \lambda, \nu)$$

↑
min L for $x \in D$

$$L(x, \lambda, \nu) = f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x)$$

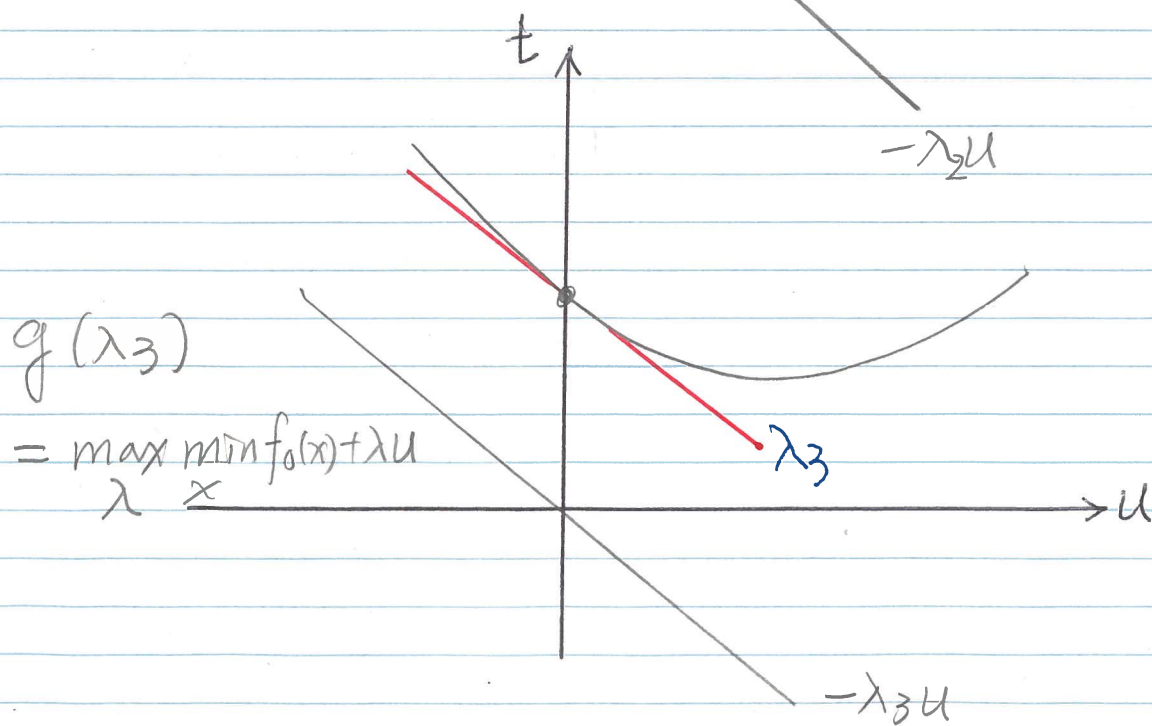
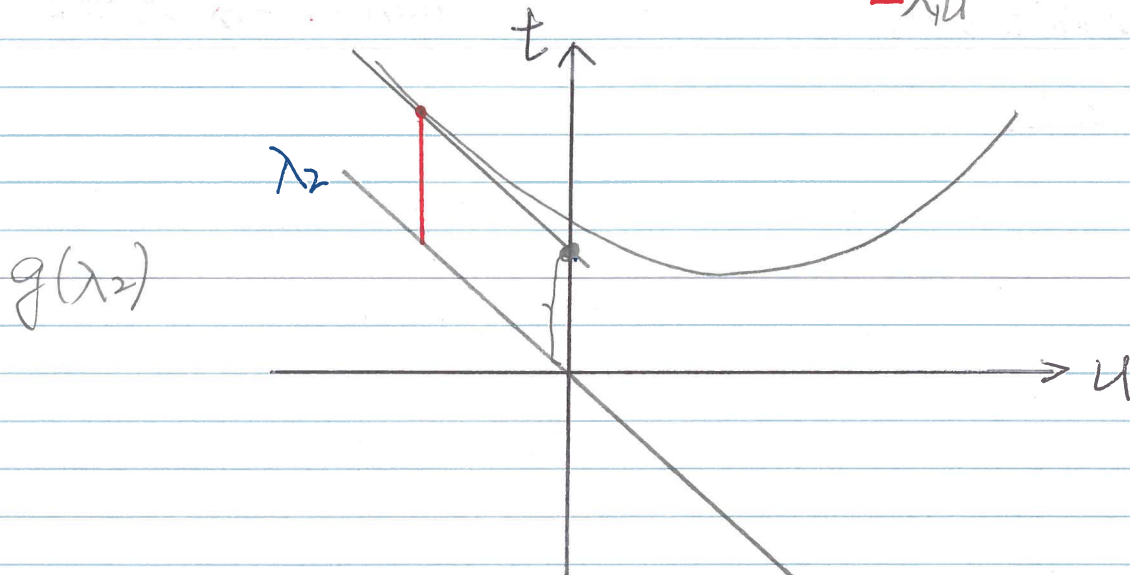
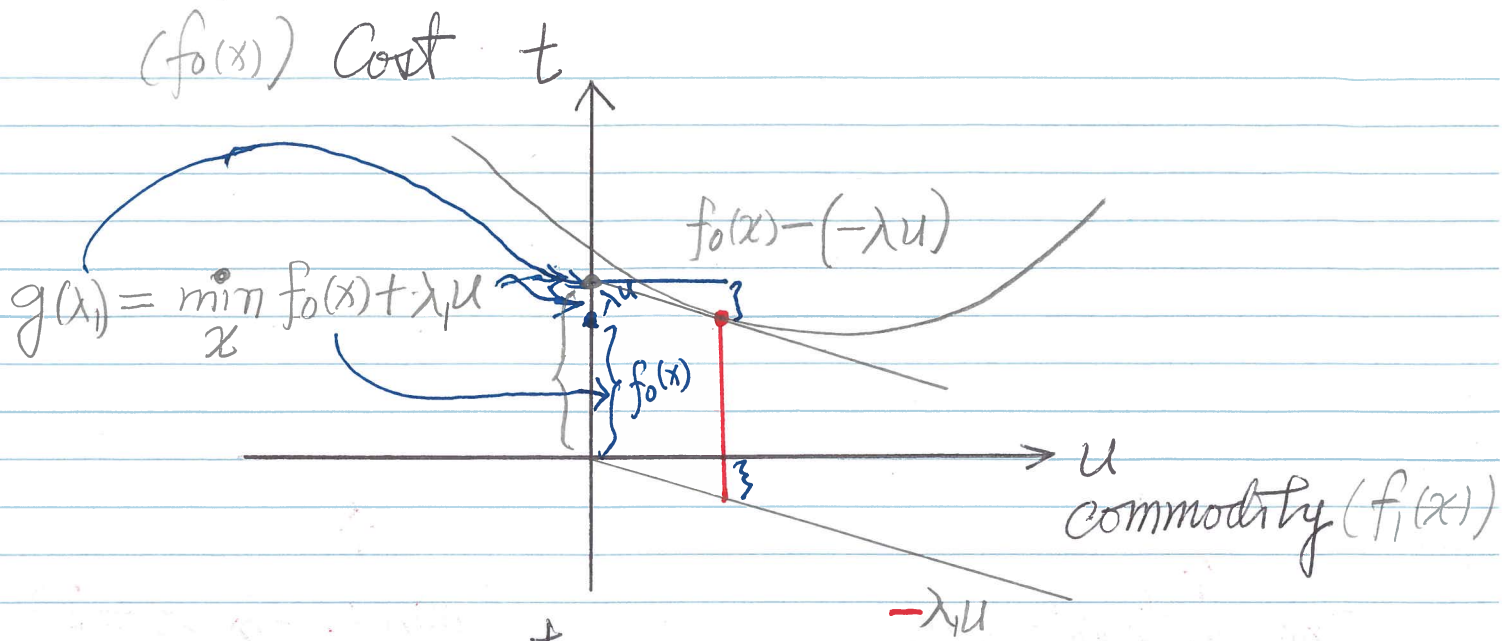
Let

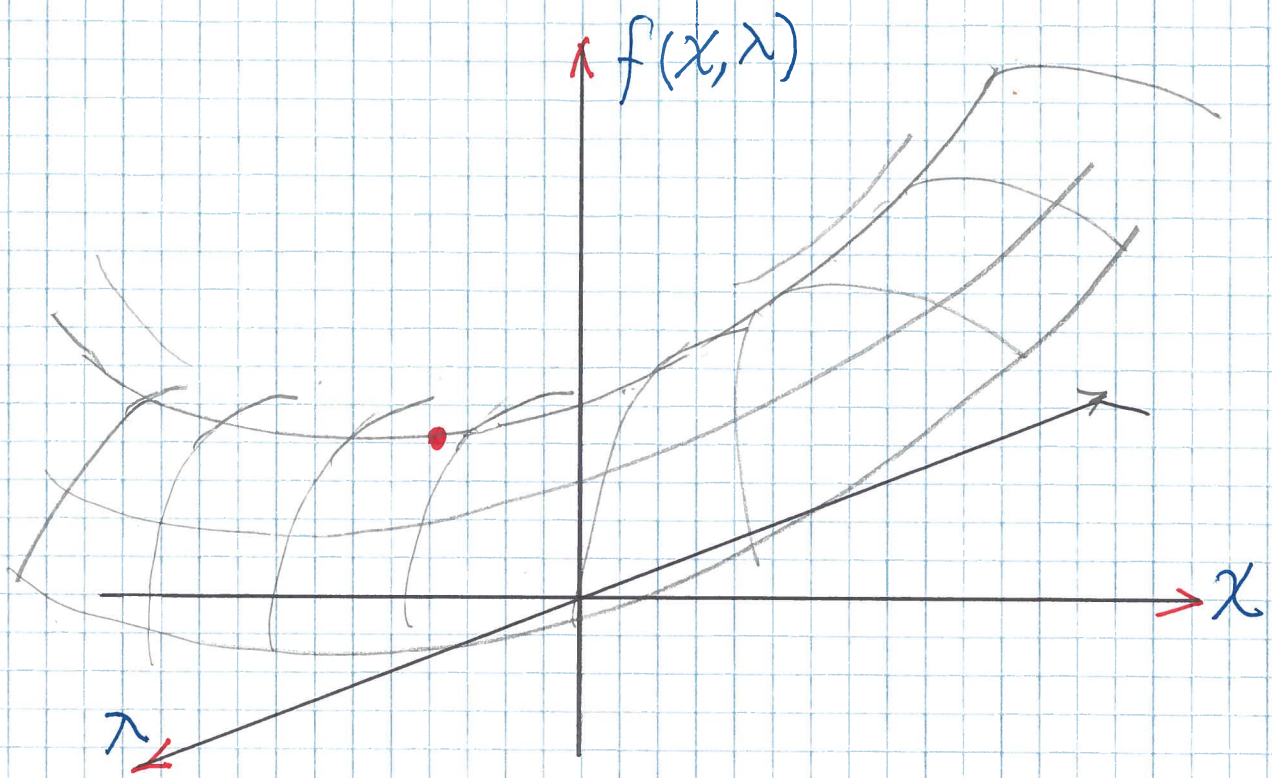
$$f(a, b) = \min_w \max_z f(w, z), \quad f(c, d) = \max_z \min_w f(w, z)$$

Then

$$f(a, b) \geq f(a, d) \geq f(c, d)$$

$f(w, z)$ is convex w.r.t. w
concave w.r.t. z





Definition

I. Saddle Point

Given function $f(\omega, z)$

$(\tilde{\omega}, \tilde{z})$ is a saddle point of f

if
$$\max_z f(\tilde{\omega}, z) = f(\tilde{\omega}, \tilde{z})$$

$$\min_{\omega} f(\omega, \tilde{z}) = f(\tilde{\omega}, \tilde{z})$$

II. Theorem:

$$\max_z \min_{\omega} f(\omega, z) = \min_{\omega} \max_z f(\omega, z)$$

iff a saddle point of f exists

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Primal Dual Interpretation

Flour, Vegetable, Fruits
 x_1, x_2, x_3

Protein, Vitamin A-K, Mineral
 $\lambda_1, \lambda_2, \lambda_3, \lambda_4, \lambda_5$

$$\text{Obj: } \min C_1 x_1 + C_2 x_2 + C_3 x_3$$

$$\begin{array}{l} \text{Protein } \lambda_1 \\ \text{Vit A } \lambda_2 \\ \text{Vit B } \lambda_3 \\ \text{Vit E } \lambda_4 \\ \text{Mineral } \lambda_5 \end{array} \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \\ a_{41} & a_{42} & a_{43} \\ a_{51} & a_{52} & a_{53} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} \geq \begin{bmatrix} b_1 \\ b_2 \\ b_3 \\ b_4 \\ b_5 \end{bmatrix}$$

$$x_1, x_2, x_3 \geq 0$$

$$\text{Obj: } \max \lambda_1 b_1 + \lambda_2 b_2 + \lambda_3 b_3 + \lambda_4 b_4 + \lambda_5 b_5$$

$$\begin{bmatrix} a_{11} & a_{21} & a_{31} & a_{41} & a_{51} \\ a_{12} & a_{22} & a_{32} & a_{42} & a_{52} \\ a_{13} & a_{23} & a_{33} & a_{43} & a_{53} \end{bmatrix} \begin{bmatrix} \lambda_1 \\ \lambda_2 \\ \lambda_3 \\ \lambda_4 \\ \lambda_5 \end{bmatrix} \leq \begin{bmatrix} C_1 \\ C_2 \\ C_3 \end{bmatrix}$$

$$\lambda_i \geq 0 \quad \forall i$$

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