

Interpretation: Saddle-point

Example : $f(w, z)$ $w = \begin{matrix} 1 & [1 & 3 & 2] \\ 2 & [2 & 1 & 3] \\ 3 & [3 & 2 & 1] \end{matrix}$
 $z = 1, 2, 3$

$$\min_{w \in W} f(w, 1) = 1$$

$$\max_{z \in Z} f(1, z) = 3$$

$$\min_{w \in W} f(w, 2) = 1$$

$$\max_{z \in Z} f(2, z) = 3$$

$$\min_{w \in W} f(w, 3) = 1$$

$$\max_{z \in Z} f(3, z) = 3$$

$$\max_{z \in Z} \min_{w \in W} f(w, z) = 1$$

$$\min_{w \in W} \max_{z \in Z} f(w, z) = 3$$

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Interpretation: Saddle-point

Example : $f(w, z)$ $w = \begin{matrix} 1 & [4 & 6 & 3] \\ 2 & [2 & \textcircled{3} & 1] \\ 3 & [3 & 5 & 2] \end{matrix}$
 $z = 1, 2, 3$

$$w=2 \quad f(w, z) :$$

$$\min_{w=2} \max_{z=2} f(2, z) = 3$$

$$\max_{z=2} \min_{w=2} f(w, z)$$

$$\min_{w=2} f(w, z) = 3 \quad w \triangleright A \text{ midterm}$$

w \triangleright B

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Interpretation: Saddle-point

Claim : Result of II \geq Result of I

Given an arbitrary pair $(\tilde{w}, \tilde{z}) \in D$

$$\min_{w \in W} f(w, \tilde{z}) \leq f(\tilde{w}, \tilde{z}) \leq \max_{z \in Z} f(\tilde{w}, z) \quad \forall \tilde{w}, \tilde{z} \in D$$

$$\min_{w \in W} f(w, \tilde{z}) \leq \max_{z \in Z} f(\tilde{w}, z)$$

$$\text{Thus } \max_{z \in Z} \min_{w \in W} f(w, z) \leq \min_{w \in W} \max_{z \in Z} f(w, z)$$

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Interpretation: Saddle-point

Example : $f(w, z)$ $w = \begin{matrix} 1 & 1 & 1 \\ 2 & 2 & 2 \\ 3 & 3 & 3 \end{matrix}$
 $z = 1, 2, 3$

$$\min_{w \in W} f(w, 1) = 1$$

$$\max_{z \in Z} f(1, z) = 1$$

$$\min_{w \in W} f(w, 2) = 1$$

$$\max_{z \in Z} f(2, z) = 2$$

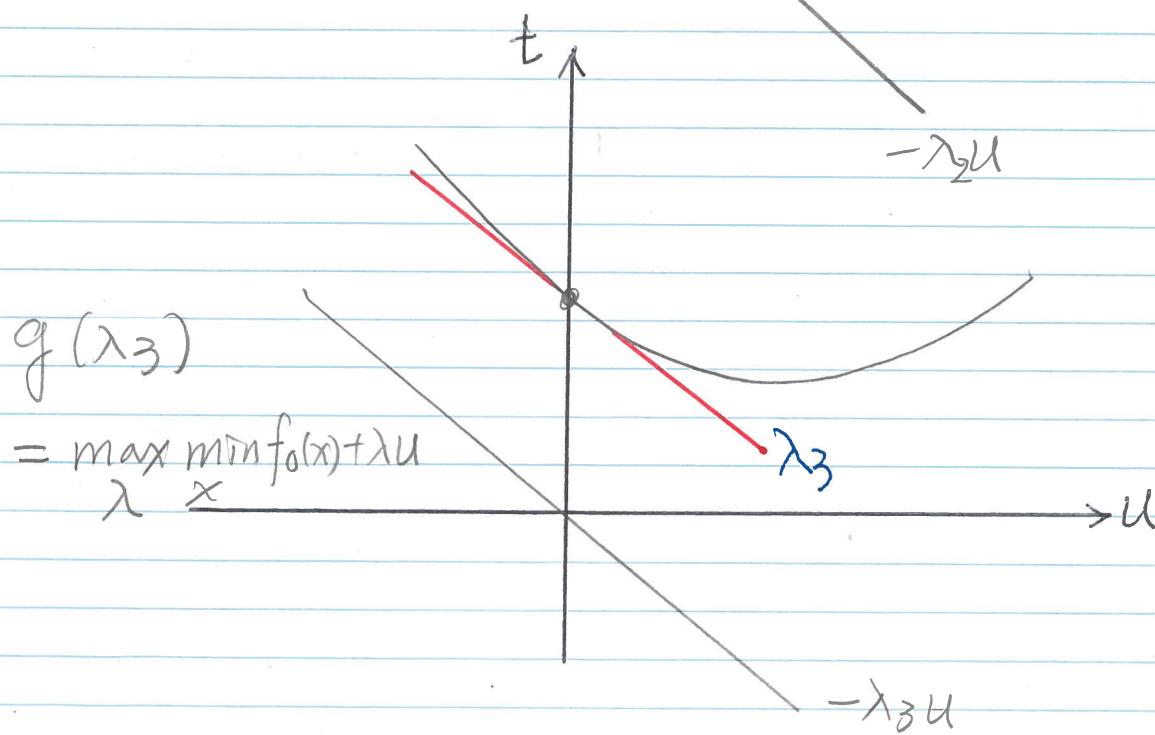
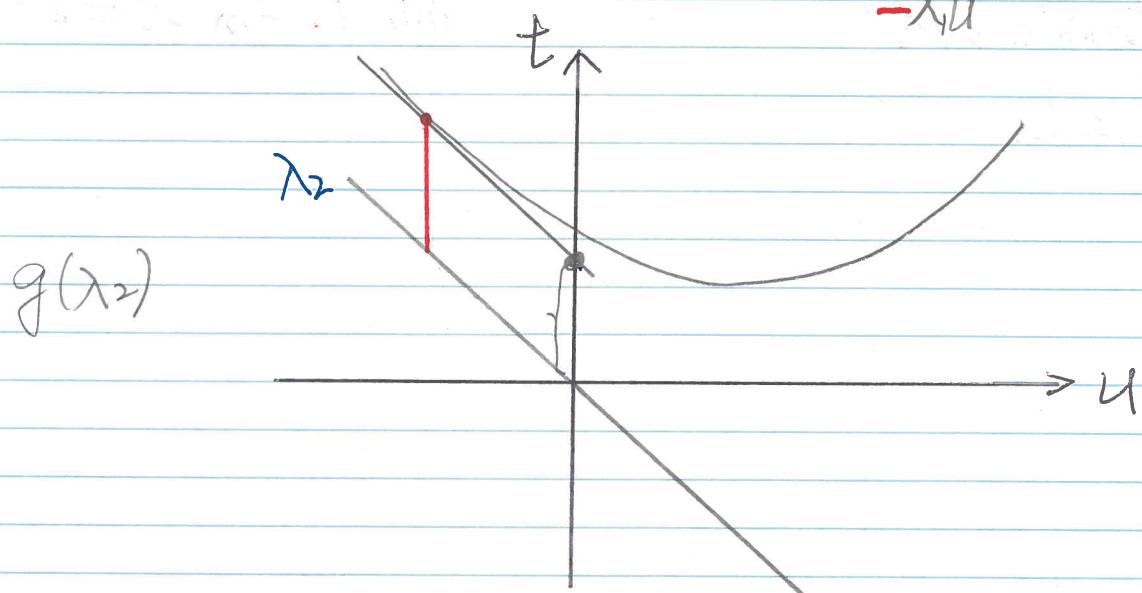
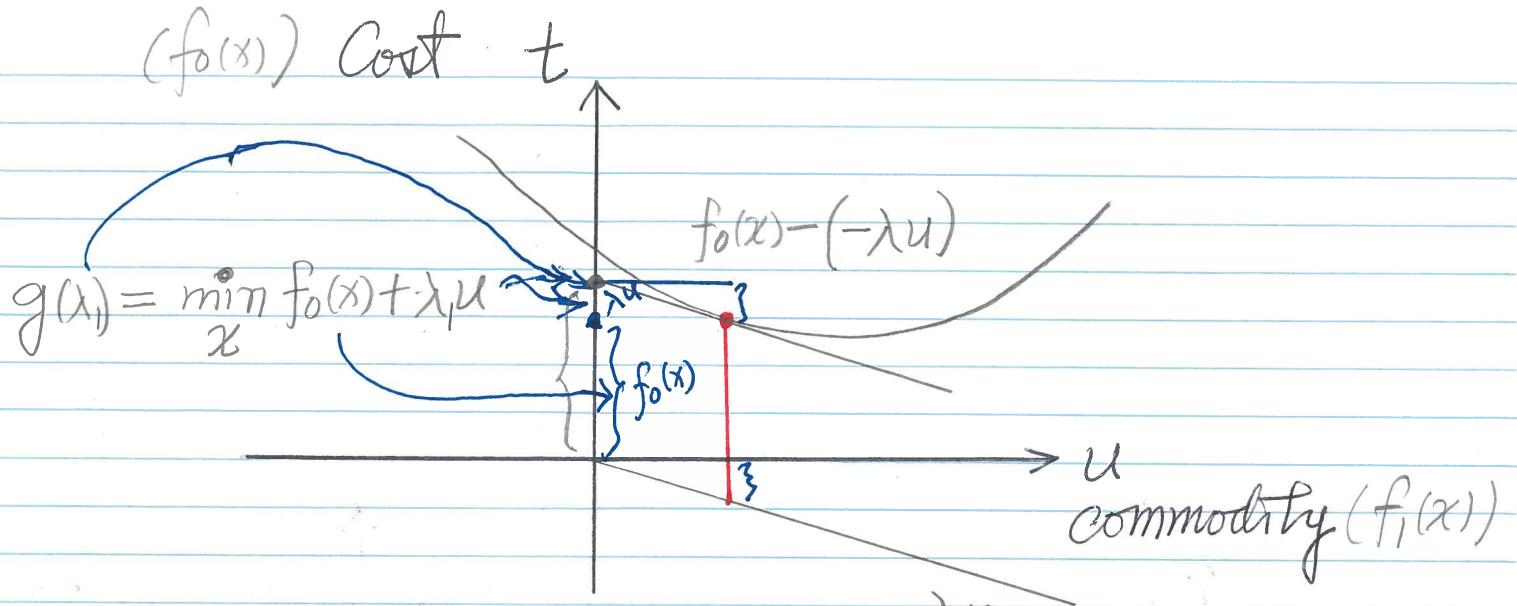
$$\min_{w \in W} f(w, 3) = 1$$

$$\max_{z \in Z} f(3, z) = 3$$

$$\max_{z \in Z} \min_{w \in W} f(w, z) = 1$$

$$\min_{w \in W} \max_{z \in Z} f(w, z) = 1$$

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$$= \max_{\lambda} \min_x f_0(x) + \lambda u$$

A.21.1

Primal

$$\min f_0(x)$$

s.t.

$$f_i(x) \leq 0$$

$$h_i(x) = 0$$

$$\min_x \max_{\lambda, \nu} L(x, \lambda, \nu)$$

The constraints are
enforced.

Dual

$$\max g(\lambda)$$

$$\min_x L(x, \lambda, \nu)$$

$$\max_{\lambda, \nu} \min_x L(x, \lambda, \nu)$$

$\min L$ for $x \in D$

$$L(x, \lambda, \nu) = f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{j=1}^p \nu_j h_j(x)$$

Let

$$f(a, b) = \min_w \max_z f(w, z), \quad f(c, d) = \max_z \min_w f(w, z)$$

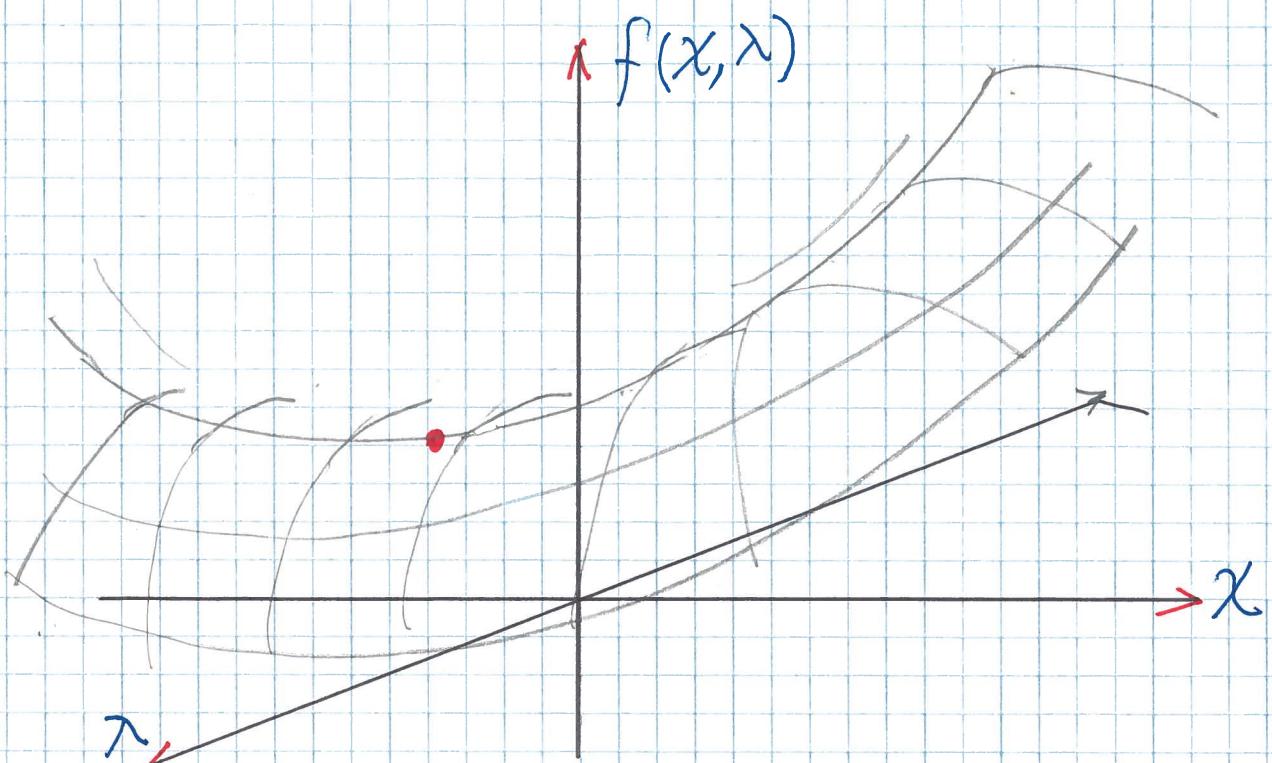
Then

$$f(a, b) \geq f(a, d) \geq f(c, d)$$

$f(w, z)$ is convex w.r.t. w

concave w.r.t. z

A.21.0



Definition

I. Saddle Point

Given function $f(w, z)$

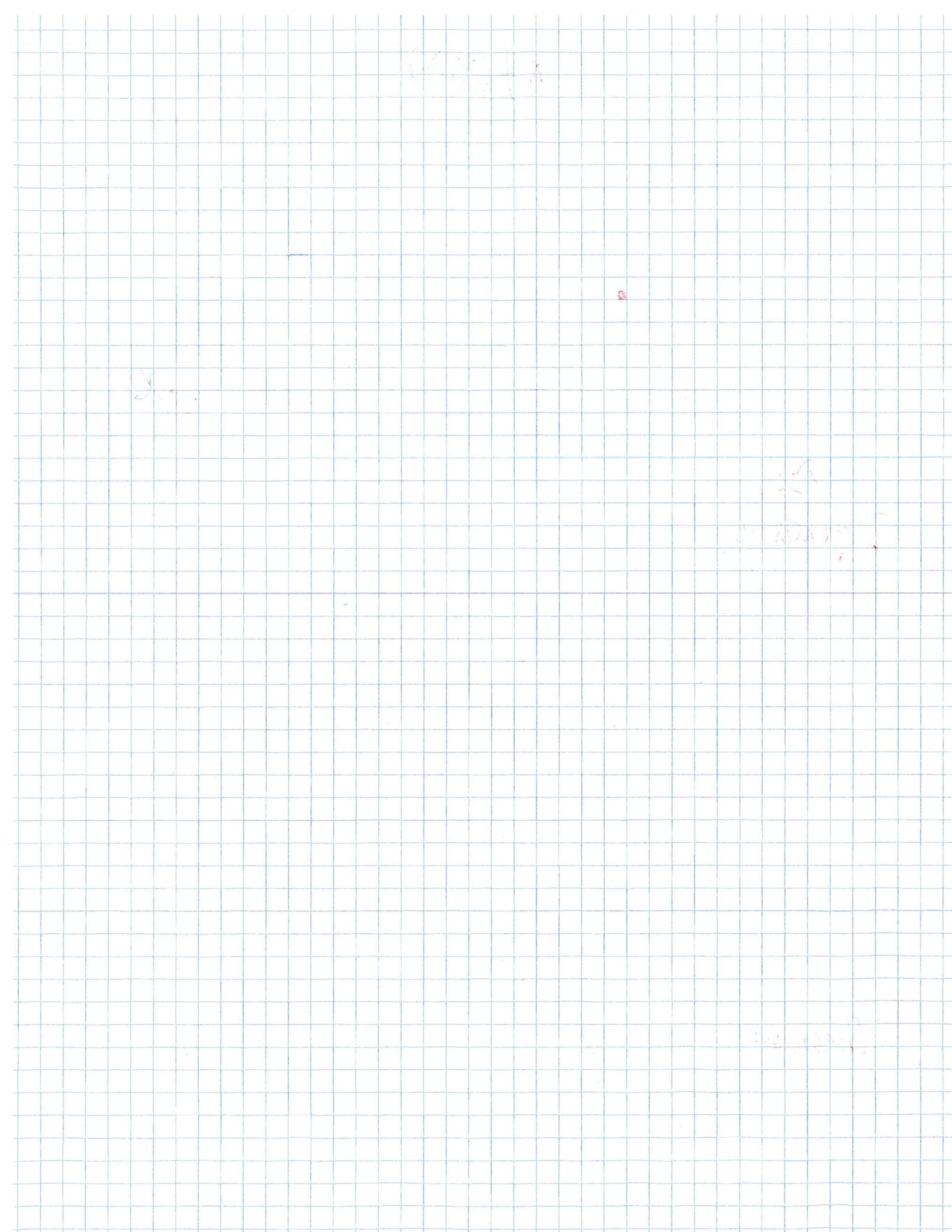
(\tilde{w}, \tilde{z}) is a saddle point of f

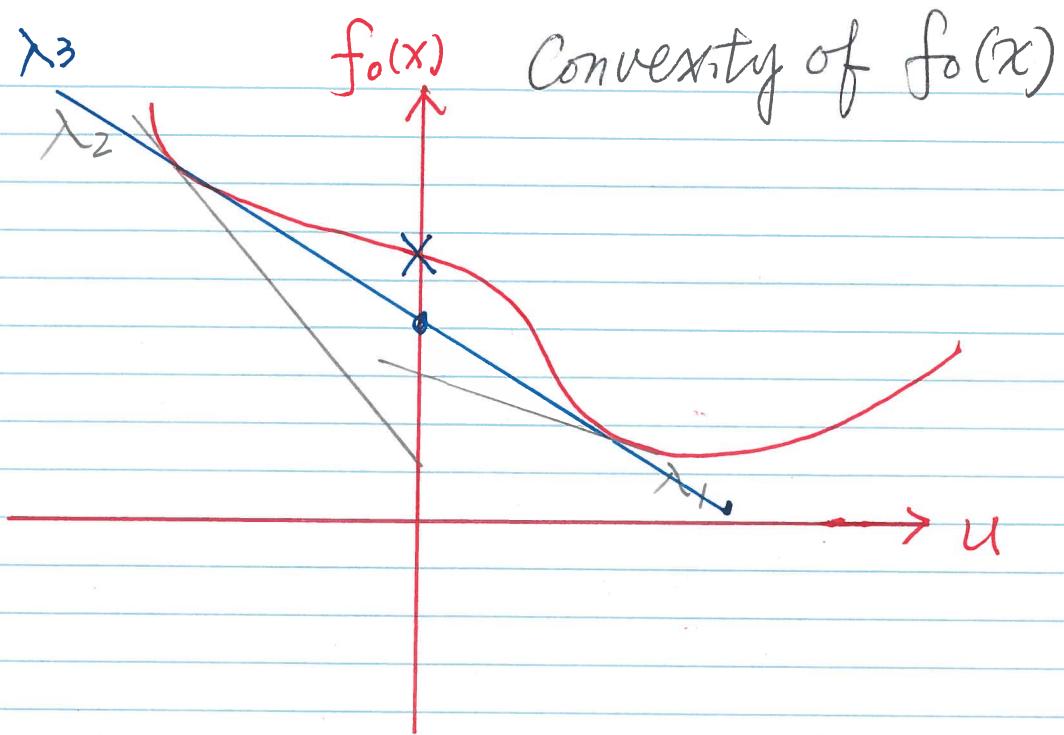
if $\max_z f(\tilde{w}, z) = f(\tilde{w}, \tilde{z})$

$$\min_w f(w, \tilde{z}) = f(\tilde{w}, \tilde{z})$$

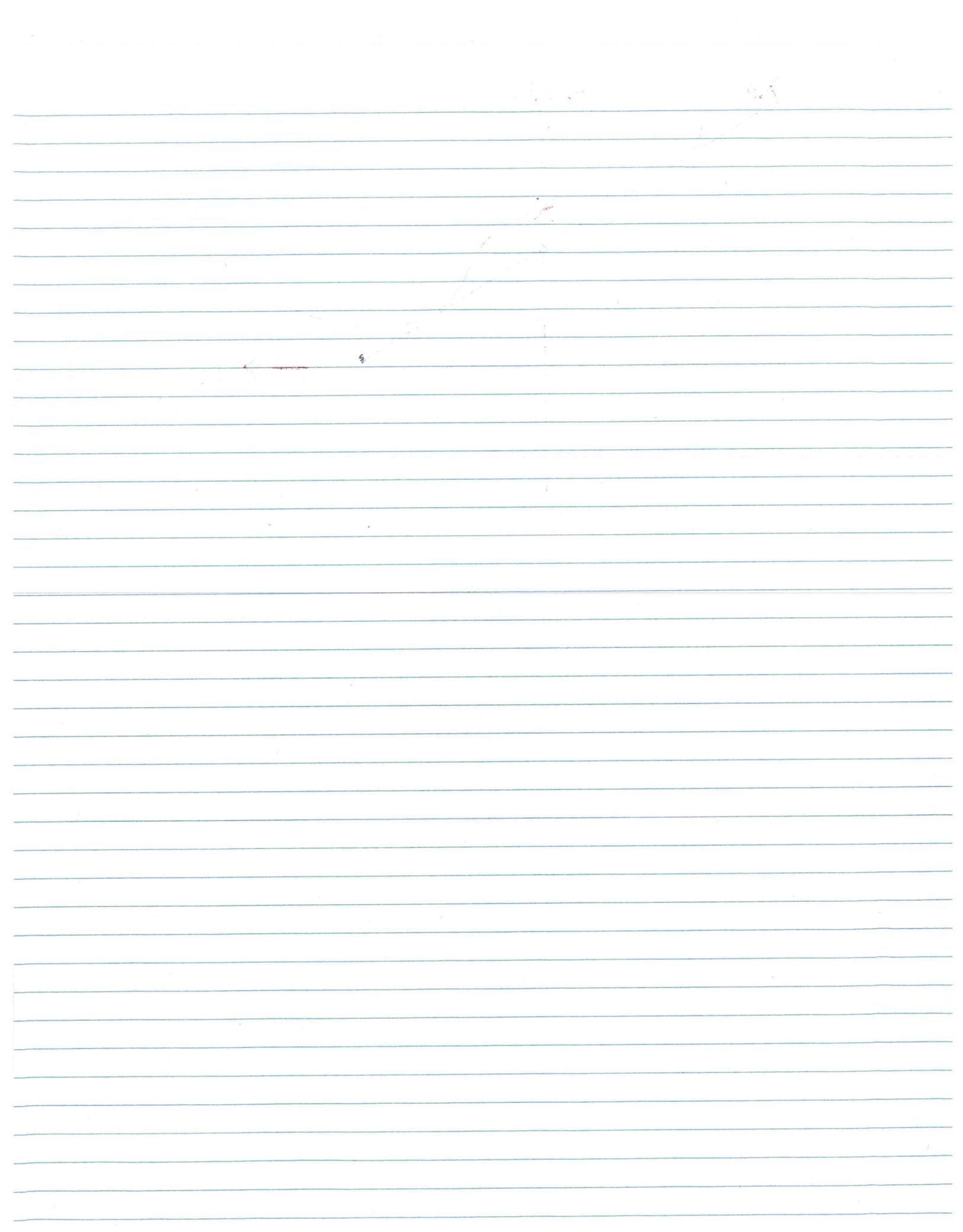
II. Theorem: $\max_z \min_w f(w, z) = \min_w \max_z f(w, z)$

iff a saddle point of f exists.





A.21.2



Proof: Necessity

Assume that

$$\min_w \max_z f(w, z) = \max_z \min_w f(w, z)$$

Let $\tilde{w} = \arg \min_w \max_z f(w, z)$

$$\tilde{z} = \arg \max_z \min_w f(w, z)$$

We have

$$f(\tilde{w}, \tilde{z}) \leq \max_z f(\tilde{w}, z) = \min_w f(w, \tilde{z}) \leq f(\tilde{w}, \tilde{z})$$

By definition

(\tilde{w}, \tilde{z}) is a saddle point.

Sufficiency

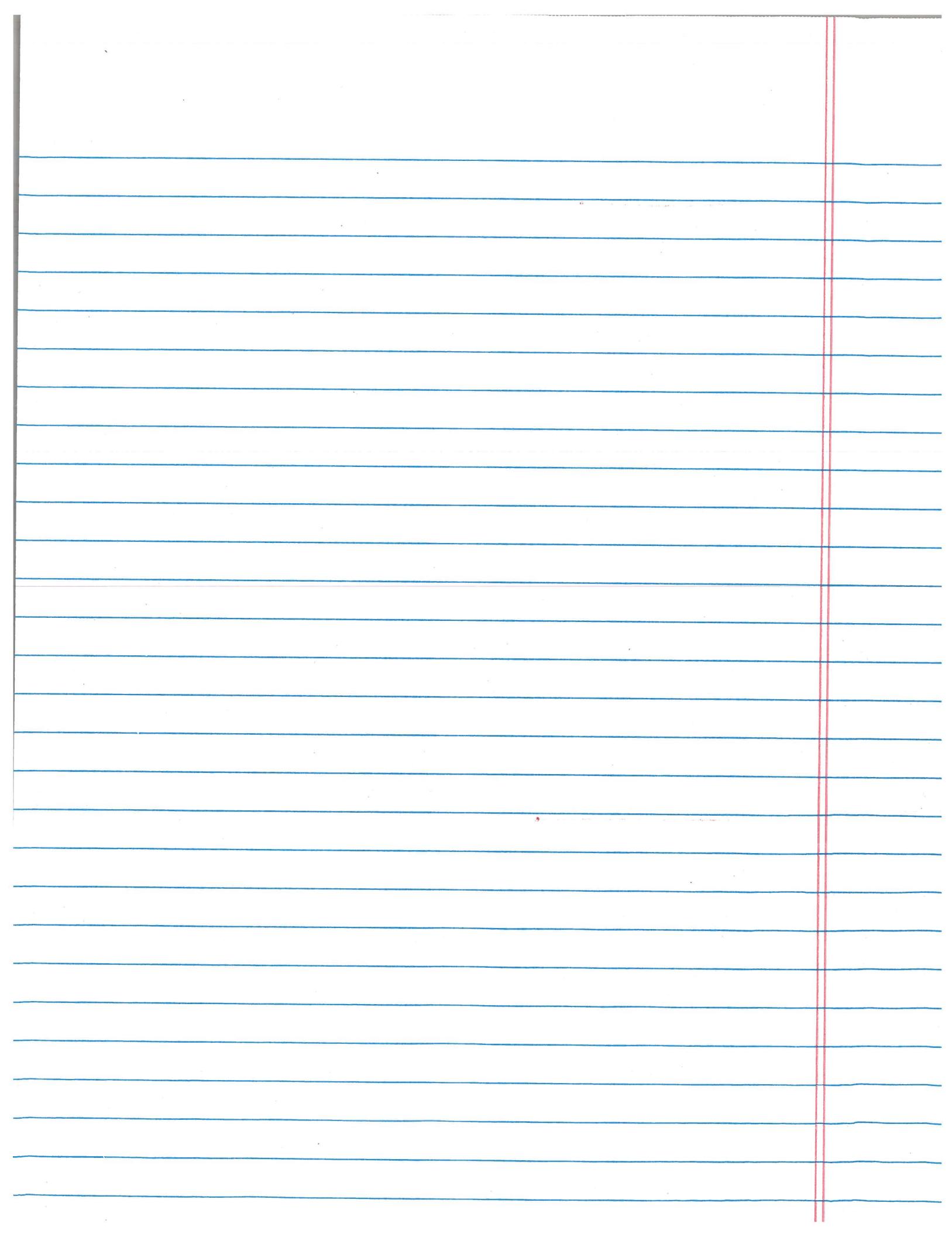
Assume that (\tilde{w}, \tilde{z}) is a saddle point

We have

$$\max_z \min_w f(w, z) \geq \min_w f(w, \tilde{z}) = f(\tilde{w}, \tilde{z})$$

$$\min_w \max_z f(w, z) \leq \max_z f(\tilde{w}, z) = f(\tilde{w}, \tilde{z})$$

Thus, $\max_z \min_w f(w, z) = \min_w \max_z f(w, z)$



Formulation: The row & column selection is formulated as a bilinear optimization problem.

$$\min_{\omega} \max_{z} f(\omega, z) = \sum_i \sum_j a_{ij} \omega_i z_j \quad \left| \quad \max_{z} \min_{\omega} f(\omega, z)$$

I. row & column selection constraints

where

$$\omega_i, z_j \in \{0, 1\} \quad \sum \omega_i = 1 \quad \sum z_j = 1$$

II. relaxed constraints

$$\sum \omega_i = 1 \quad \sum z_j = 1 \quad \omega_i \geq 0, z_j \geq 0, \forall i, j.$$

A. The optimization problem with relaxed constraints can be solved with algorithms Dantzig

$$\min_{\omega} \max_{z} f(\omega, z) = \max_{z} \min_{\omega} f(\omega, z)$$

B. Since $f(\omega, z)$ is convex w.r.t ω concave w.r.t z .

The solution can reduce to constraint I (row & column selection)

$$f(\tilde{\omega}, \tilde{z})$$

C. From B, $(\tilde{\omega}, \tilde{z})$ is a saddle point.

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