

# CSE203B Convex Optimization:

## Chapter 4: Problem Statement

CK Cheng  
Dept. of Computer Science and Engineering  
University of California, San Diego

1

## Convex Optimization Formulation

1. Introduction
  - I. Eliminating equality constants
  - II. Slack variables
  - III. Absolute values, softmax
2. Optimality Conditions
  - I. Local vs. global optimum
  - II. Optimality criterion for differentiable  $f_0$ 
    - i. Optimization without constraints
    - ii. Opt. with inequality constraints
    - iii. Opt. with equality constraints
  - III. Quasi-convex optimization
3. Linear Optimization
4. Quadratic Optimization
5. Geometric Programming
6. Generalized Inequality Constraints

*Variations* → Dual Problems  
*Same original problems*

2

# 1. Introduction

Formulation: One of the most critical processes to conduct a project.

$$\min f_0(x)$$

$$s.t. f_i(x) \leq 0 \quad i = 1, \dots, m$$

$$h_i(x) = 0 \quad i = 1, \dots, p \quad (Ax = b \text{ Affine set})$$

$$x \in R^n$$

$$D_{f_0} f_0: R^n \rightarrow R$$

$$D_{f_i} f_i: R^n \rightarrow R$$

$$D_{h_i} h_i: R^n \rightarrow R$$

$f_0, f_1, \dots, f_m$  are convex

$D = \cap_{i=0,m} D_f \cap_{i=0,p} D_{h_i}$  Domain of functions, but not the feasible set.

**Feasible Set:** The set which satisfies the constraints (is convex for convex problems).

3

## 1.1 Introduction: Eliminating Equality Constraints

$$\begin{aligned} & \min f_0(x) && \{x | Ax \leq b\} \Rightarrow \{x | Ax = b, x \geq 0\} \\ & s.t. f_i(x) \leq 0 \quad i = 1, \dots, m \\ & Ax = b \end{aligned} \quad (\text{note that this is equality relation})$$

- a. Convert  $\{x | Ax = b\}$  to  $\{Fz + x_0 | z \in R^k\}$
- b. We have an equivalent problem

$$\begin{aligned} & \min f_0(Fz + x_0) \\ & s.t. f_i(Fz + x_0) \leq 0 \end{aligned}$$

Remark: Matrix  $F$  contains columns of null space basis

$$\begin{array}{c} P \begin{bmatrix} n \\ A \\ I \\ n \end{bmatrix} P^{-1} = \begin{bmatrix} B \\ F \\ C \end{bmatrix} \xrightarrow{\text{elimination}} \begin{array}{l} A \\ F \\ C \end{array} \\ AP = B \quad AC = B \\ IP = C \Rightarrow P = C \end{array}$$

4

## 1.2 Introduction: Slack Variables

$$\begin{aligned} & \min f_0(x) \\ & s.t. f_i(x) \leq 0, i = 1, \dots, m \\ & Ax = b \end{aligned}$$

Add slack variables to convert to an equivalent problem

- a. Convert the objective function with variable  $t$

$$\begin{aligned} & \min t \\ & s.t. f_0(x) - t \leq 0 \\ & f_i(x) \leq 0, i = 1, \dots, m \\ & A^T x = b \end{aligned}$$

- b. Convert the inequality with variables  $s_i$

$$\begin{aligned} & \min f_0(x) \\ & s.t. f_i(x) + s_i = 0 \\ & A^T x = b \\ & s_i \in R_+, i = 1, \dots, m \end{aligned}$$

5

## 1.3 Introduction: Absolute values and Softmax

- a. Absolute values

$$\begin{aligned} & |f_i(x)| \leq b \\ & \Rightarrow f_i(x) \leq b \text{ and} \\ & -f_i(x) \leq b \end{aligned}$$

- b. Maximum values

$$\max\{f_1, f_2, \dots, f_m\}$$

$$\text{Softmax: } \frac{1}{\alpha} \log(e^{\alpha f_1} + e^{\alpha f_2} + \dots + e^{\alpha f_m})$$

Example:  $\max\{1, 5, 10, 2, 3\} \Rightarrow \text{Softmax}$

$$\frac{1}{\alpha} \log(e^\alpha + e^{5\alpha} + e^{10\alpha} + e^{2\alpha} + e^{3\alpha}) \approx 10$$

6

## 2.1 Optimality Conditions: Local vs. Global Optima

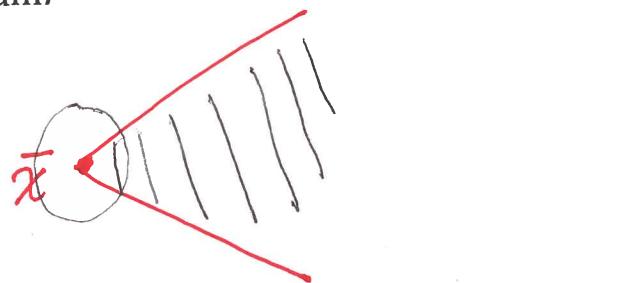
Definition: Local Optima

Given a convex optimization problem and a point  $\bar{x} \in R^n$

If there exists a  $r > 0$

s.t.  $f_0(z) \geq f_0(\bar{x})$  for all  $z \in$  Feasible Set, and  $\|z - \bar{x}\|_2 \leq r$

Then  $\bar{x}$  is a local optimum.



7

## 2.2 Optimality Conditions

Theorem: Given a convex opt. problem

If  $\bar{x}$  is a local optimum, then  $\bar{x}$  is a global optimum

Proof: By contradiction

Suppose that  $\exists y \in$  Feasible Set

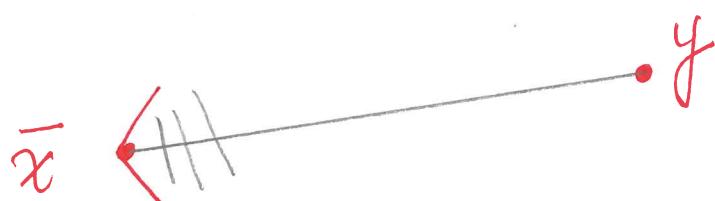
s.t.  $f_0(\bar{x}) > f_0(y)$

We have  $f_0(\bar{x}) > (1 - \theta)f_0(\bar{x}) + \theta f_0(\bar{y})$  (by assumption) I

$> f_0((1 - \theta)\bar{x} + \theta\bar{y})$  ( $f_0$  is convex) II

And  $(1 - \theta)\bar{x} + \theta\bar{y}$  is feasible (Feasible set is convex)

The inequality contradicts to the assumption of local optima.



8

# Null space

$$A = \begin{bmatrix} 1 & 2 \\ 2 & 4 \end{bmatrix}$$

$$\begin{bmatrix} A \\ I \end{bmatrix} P = \begin{bmatrix} B \\ C \end{bmatrix}$$

$$\begin{bmatrix} 1 & 2 \\ 2 & 4 \\ 1 & 0 \\ 0 & 1 \end{bmatrix} P = \begin{bmatrix} 1 & 0 \\ 2 & 0 \\ 1 & -2 \\ 0 & 1 \end{bmatrix}$$

$$AC = B$$

$$A \begin{bmatrix} 1 & 2 \\ 0 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 2 & 0 \end{bmatrix}$$

$C_1 \quad C_2$

$$C_2 - 2 \cdot C_1$$

$$\Rightarrow \text{Null space } \left\{ \begin{bmatrix} -2 \\ 1 \end{bmatrix} + \begin{bmatrix} 0 \\ 0 \end{bmatrix} \mid \theta \in \mathbb{R} \right\}$$

A4

