

## Infeasible Start Newton's Method

The search of the feasible start point,

$$\begin{bmatrix} \nabla^2 f(x) & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} \Delta x \\ v \end{bmatrix} = - \begin{bmatrix} \nabla f(x) \\ Ax - b \end{bmatrix}$$

We can write in incremental derivation,

$$\begin{bmatrix} \nabla^2 f(x) & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} \Delta x \\ \Delta v \end{bmatrix} = - \begin{bmatrix} \nabla f(x) + A^T v \\ Ax - b \end{bmatrix}$$

$$\begin{bmatrix} \nabla^2 f(x) & 0 \\ 0 & -A \nabla^2 f(x)^{-1} A^T \end{bmatrix} \begin{bmatrix} \Delta x \\ \Delta v \end{bmatrix}$$

$$\begin{aligned} \min f(x) \\ \underbrace{Ax = b}_{\Delta x} \rightarrow r = Ax - b \\ A \Delta x = -Ax + b \\ \begin{array}{l} r_{dual} \\ r_{pri} \end{array} \downarrow \\ A(x + \Delta x) = b \end{aligned}$$

KLU, SuperLU  
direct matrix solvers.

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## Newton Method: Infeasible Start

*Algorithm.*

Given  $x \in D$ ,  $v$ , tolerance  $\epsilon > 0$ ,  $\alpha \in (0, \frac{1}{2})$ ,  $\beta \in (0, 1)$ .

Repeat

1. Compute primal and dual Newton steps  $\Delta x_{nt}, \Delta v_{nt}$
2. Line search on  $\|r(x, v)\|_2 = \|(r_{dual}(x, v), r_{pri}(x, v))\|_2$

$$t := 1$$

while  $\|r(x + t\Delta x_{nt}, v + t\Delta v_{nt})\|_2 > (1-\alpha t)\|r(x, v)\|_2$   
 $t := \beta t$ .

3. Update  $x := x + t\Delta x_{nt}, v := v + t\Delta v_{nt}$

Until  $Ax = b$  and  $\|r(x, v)\|_2 \leq \epsilon$

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# Summary

KKT Linear Equations:  
Quadratic objective function + linear equality constraints

Newton's Method:  
**Twice differentiable obj function** + linear equality constraints

Interior Point Method:  
Twice differentiable obj function + linear equality + **inequality constraints**

# CSE203B Convex Optimization

## Chapter 11 Interior Point Methods

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## Chapter 11: Interior-Point Methods

- Formulation
  - Inequality constrained optimization
- Barrier Method
- Generalized Inequalities Problems
- Primal Dual Interior Point Methods
- Summary

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# Formulation: The problem

Problem:  $\min f_0(x)$

Subject to  $f_i \leq 0, i = 1, \dots, m$

$$Ax = b$$

Function  $f_i$ s are convex, twice continuously differentiable  
We assume that  $\text{rank } A = p, A \in R^{p \times n}$ .

Issues:

1.  $m$  can be exponential.
2. When to put  $f_i = 0$  (active)? There are  $2^m$  combination.

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# Formulation: logarithmic barrier

Problem:

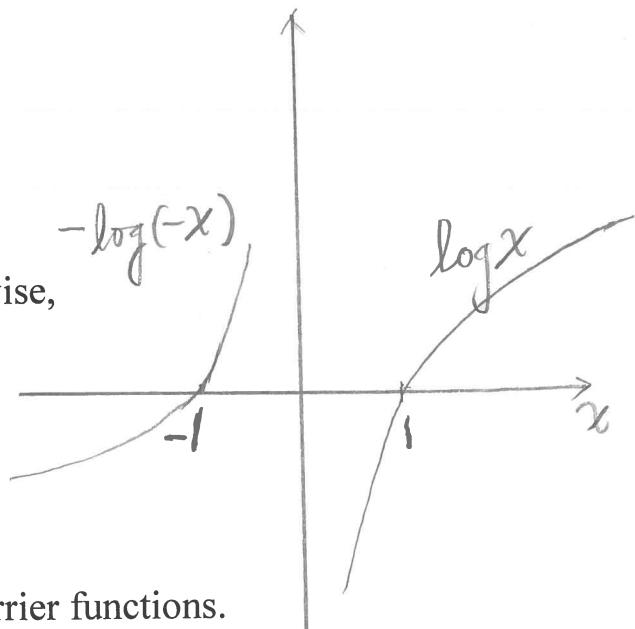
$$\min f_0(x) + \sum_{i=1}^m I_{f_i(x)}$$

$$\text{s.t. } Ax = b$$

When  $I_u = 0$  if  $u \leq 0, I_u = \infty$ . Otherwise,

$$\min f_0(x) + \frac{-1}{t} \sum_{i=1}^m \log(-f_i(x))$$

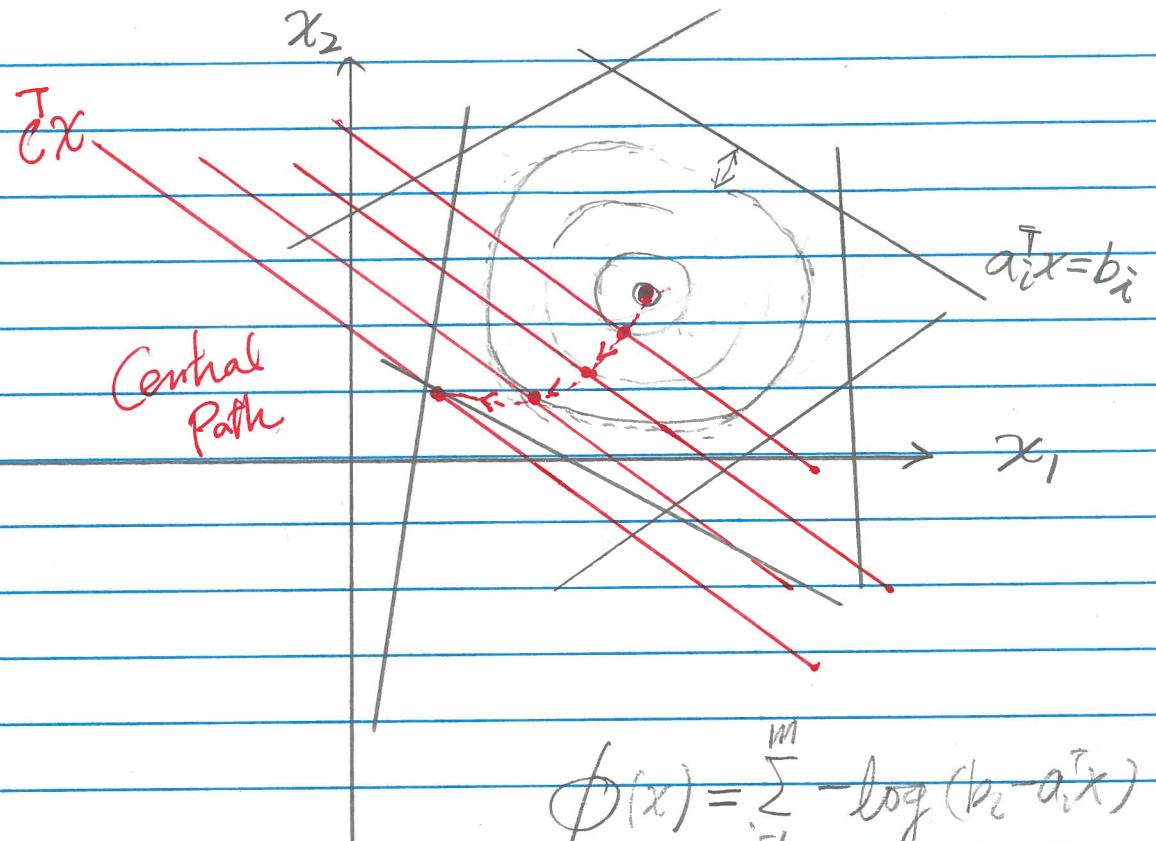
$$\text{s.t. } Ax = b$$



Remark:

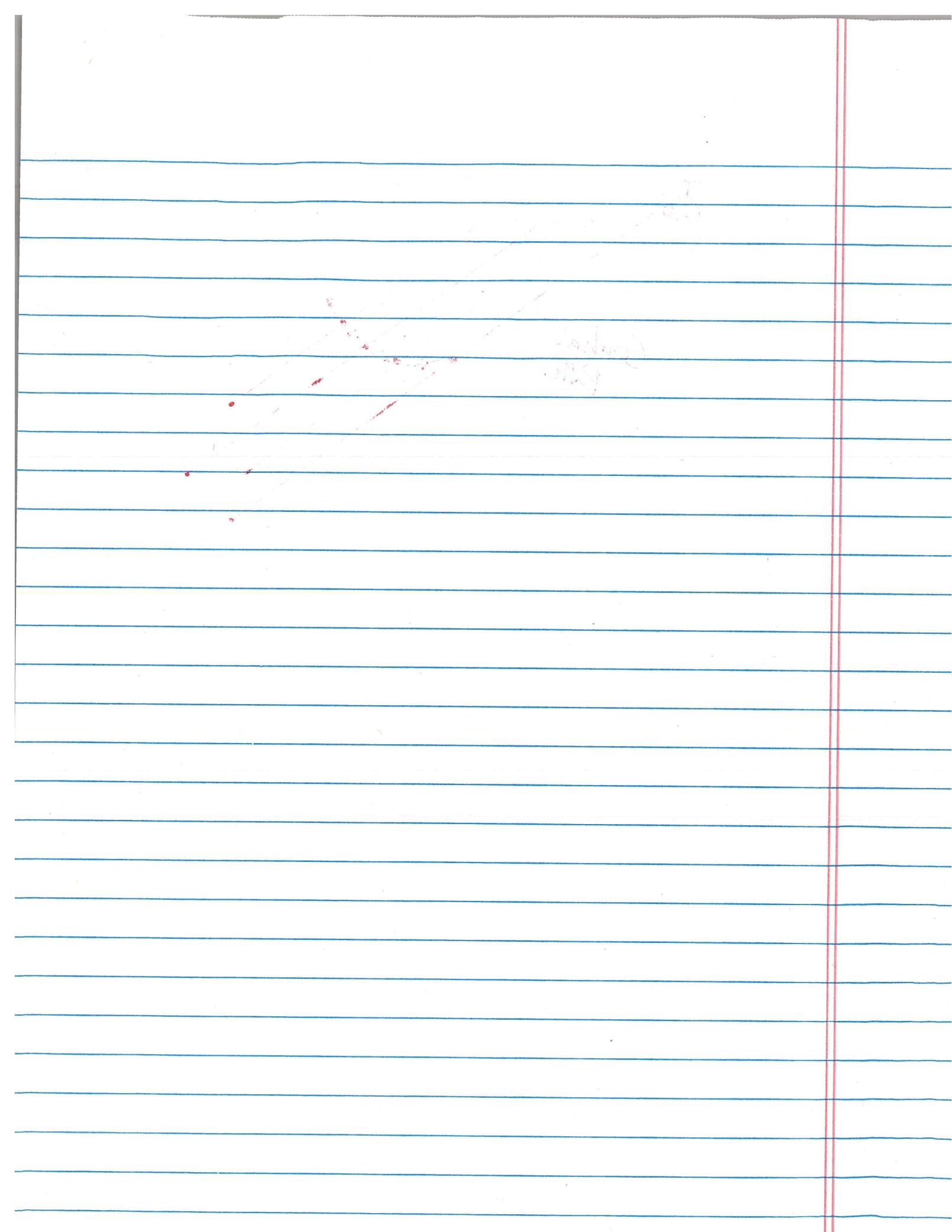
1. Convert inequality constraints to barrier functions.
2. Incorporate barrier functions in objective function.
3. Increase  $t$  to improve accuracy.

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$$\phi(x) = \sum_{i=1}^m -\log(b_i - a_i^T x)$$

$$\nabla \phi(x) = \sum_{i=1}^m -\frac{a_i}{b_i - a_i^T x}$$



## Formulation: logarithmic barrier

Let us set

$$\phi(x) = -\sum_{i=1}^m \log(-f_i(x)), \quad \text{dom } \phi = \{x | f_i(x) < 0\}$$

$\phi(x)$  is convex and twice differentiable

$$\nabla \phi(x) = \sum_{i=1}^m -\frac{1}{f_i(x)} \nabla f_i(x)$$

$$\nabla^2 \phi(x) = \sum_{i=1}^m \frac{1}{f_i(x)^2} \nabla f_i(x) \nabla f_i(x)^T - \frac{1}{f_i(x)} \nabla^2 f_i(x)$$

Central Path is  $\{x^*(t) | t > 0\}$

$$\min t f_0(x) + \phi(x)$$

$$\text{s.t. } Ax = b$$

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## Formulation: logarithmic barrier

Ex:

Problem:  $\min c^T x$   
 s.t.  $a_i^T x \leq b_i, \quad i = 1, \dots, m$

Log barrier formulation:

$$\min t c^T x - \sum_{i=1}^m \log(b_i - a_i^T x)$$

Hyperplane  $c^T x = c^T x^*(t)$  is tangent to real curve  $\varphi$  through  $x^*(t)$ .

Solution  $x^*(t)$  balance the force between  $-t \nabla f_0(x)$  and  $\sum_{i=1}^m -\frac{1}{f_i(x)} \nabla f_i(x)$ .

$$t c + \sum_{i=1}^m -\frac{1}{(b_i - a_i^T x)} a_i = 0$$

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# Formulation: logarithmic barrier

Ex:

$$\text{Problem: } \min c^T x$$

$$\text{s.t. } a_i^T x \leq b_i \quad i = 1, \dots, m$$

$$-t \nabla f_0(x) = -tc$$

$$\left\{ \sum_{i=1}^m -\frac{1}{f_i(x)} \nabla f_i(x) = \sum_{i=1}^m -\frac{1}{b_i - a_i^T x} a_i \right.$$

$$\text{Note that } \min \left\| \frac{1}{b_i - a_i^T x} a_i \right\|_2 = \frac{1}{\text{dist}(x_i H_i)}, \quad H_i = \{x | a_i^T x = b_i\}$$

$$a_i^T x = b_i$$

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## Barrier Method: Algorithm

Given strictly feasible  $x, t = t^0 > 0, \mu > 1, \epsilon > 0$

Repeat (10-20)

1. **Centering step** to find solution  $x^*(t)$

Problem:  $\min t f_0(x) + \phi(x) \quad (\text{Newton's method})$

s.t.  $Ax = b$

2. Update  $x = x^*(t)$

3. Stopping criterion: exit if  $\frac{m}{t} < \epsilon$

4. Increase  $t = \mu t$

Complexity: # Repeats (Outer iterations) =  $\frac{\log(\frac{m}{\epsilon t(0)})}{\log \mu}$

Plus the initial centering step  $x^*(t^{(0)})$

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## Barrier Method: Newton's Step for Modified KKT

$$\begin{bmatrix} t\nabla^2 f_o(x) + \nabla^2 \phi(x) & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} \Delta x \\ v \end{bmatrix} = - \begin{bmatrix} t\nabla f_o(x) + \nabla \phi(x) \\ 0 \end{bmatrix}$$

$$\begin{aligned} \nabla \sum_{i=1}^m (-\log(-f_i(x))) &= \sum_{i=1}^m -\frac{1}{f_i(x)} \nabla f_i(x) \\ \nabla^2 \sum_{i=1}^m (-\log(-f_i(x))) &= \sum_{i=1}^m \left[ -\frac{1}{f_i(x)} \nabla^2 f_i(x) + \frac{1}{f_i(x)^2} \nabla f_i(x) \nabla f_i(x)^T \right] \end{aligned}$$

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## Barrier Method: Central Path

$$\begin{aligned} \text{Min } f_0(x) + \frac{-1}{t} \sum_{i=1}^m \log(-f_i(x)) \\ \text{s.t. } Ax = b \end{aligned}$$

$$\text{Lagrangian: } L(x, v) = f_0(x) + \frac{-1}{t} \sum_{i=1}^m \log(-f_i(x)) + v^T(Ax - b)$$

For an optimal solution, we have  $(x^*(t), \bar{v}(t))$

$$\nabla f_0(x^*) + \sum -1/(tf_i(x^*)) \nabla f_i(x^*) + A^T \bar{v} = 0$$

We can view the dual points from central path:

$$\lambda_i^*(t) = -1/(tf_i(x^*)), i = 1, \dots, m$$

Original Lagrangian:

$$L(x, \lambda, v) = f_0(x) + \sum \lambda_i f_i(x) + v^T(Ax - b)$$

Replace with  $(x^*(t), \lambda^*(t), \bar{v}(t))$ :

$$L(x^*, \lambda^*, \bar{v}) = f_0(x^*) + \sum \lambda_i^* f_i(x^*) + \bar{v}^T(Ax^* - b) = f_0(x^*) - \frac{m}{t}$$

Thus, we have  $f_0(x^*(t)) - p^* \leq m/t$

$$P^* = \max_{\lambda, v} g(\lambda, v) \geq g(\lambda^*, \bar{v}) = \min_x L(x, \lambda^*, \bar{v}) = f_0(x^*) - \frac{m}{t}$$

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# Barrier Method: Feasible Solution Search

Search 1:

$$\min s$$

$$s.t. f_i(x) \leq s, i = 1, \dots, m$$

$$Ax = b, s \in R$$

Search 2:

$$\min 1^T s, \quad s \in R_+^m$$

$$s.t. f_i(x) \leq s_i, i = 1, \dots, m$$

$$Ax = b$$

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## Barrier Method: complexity analysis

#Repeats (outer iterations)

$$= \text{Ceiling}(\log(m/(\epsilon t^0))/\log \mu)$$

#Newton steps per outer iteration (self-concordance)

$$= \frac{m(\mu - 1 - \log \mu)}{\gamma} + \log_2 \log_2 1/\epsilon_{nt},$$

$$\text{where } \gamma = \alpha \beta (1 - 2\alpha)^2 / (20 - 8\alpha)$$

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