CSE 105
THEORY OF COMPUTATION

"Winter" 2018

http://cseweb.ucsd.edu/classes/wi18/cse105-ab/

Today's learning goals

Sipser Ch 4.2

- Trace high-level descriptions of algorithms for computational problems.
- Use counting arguments to prove the existence of unrecognizable (undecidable) languages.
- Use diagonalization in a proof of undecidability.

Computational problems

A computational problem is **decidable** iff the language encoding the problem instances is decidable.

In Siper 4.1: The computational problems below

A_{DFA}, A_{NFA}, A_{REX}, A_{CFG}

E_{DFA}, E_{NFA}, E_{REX}, E_{CFG}

EQ_{DFA}, EQ_{NFA}, EQ_{REX}, EQ_{CFG}

are all decidable

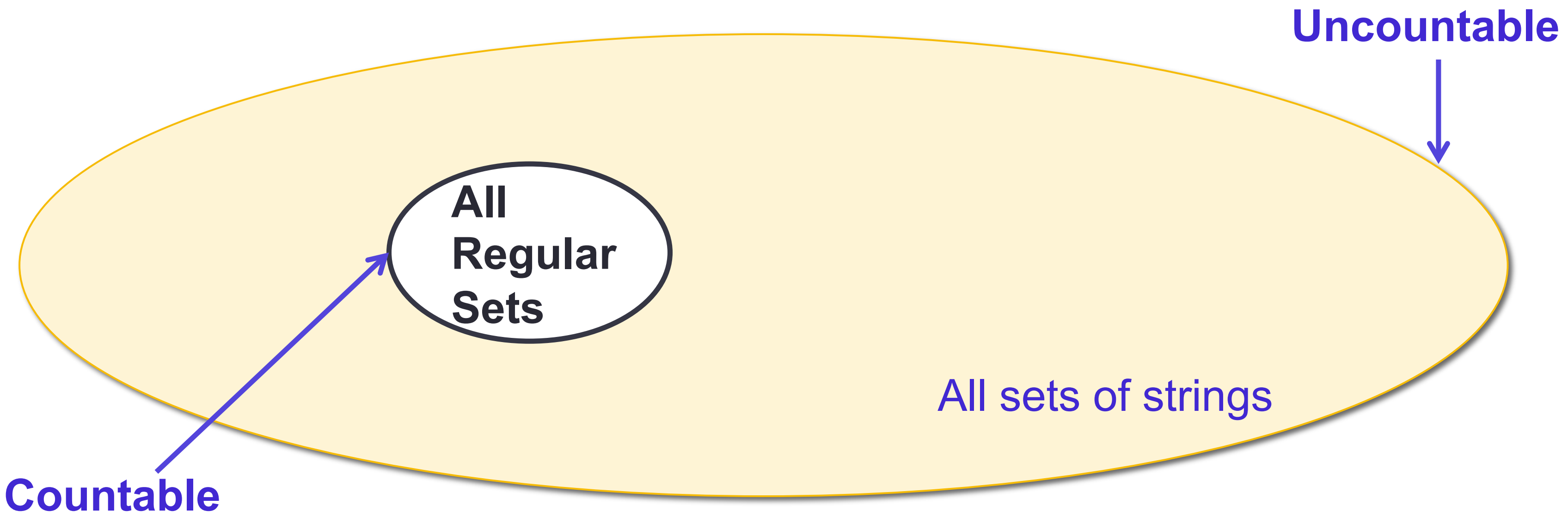
Undecidable?

- There are many ways to prove that a problem ***is*** decidable.
- How do we find (and prove) that a problem ***is not*** decidable?

Counting arguments

Before we proved the Pumping Lemma …

We proved there was a set that was not regular because



Reminder: countable/ uncountable

Sets A and B have the **same size** |A| = |B| means there is a function between them that is both one-to-one and onto.

Countable (finite or same size as N)	**Uncountable**
N, Z, Q	R
Σ^*	{ infinite sequences over Σ}
The set of all TMs	$P(\Sigma^*)$

Sipser p. 202-204

Reminder: countable/ uncountable

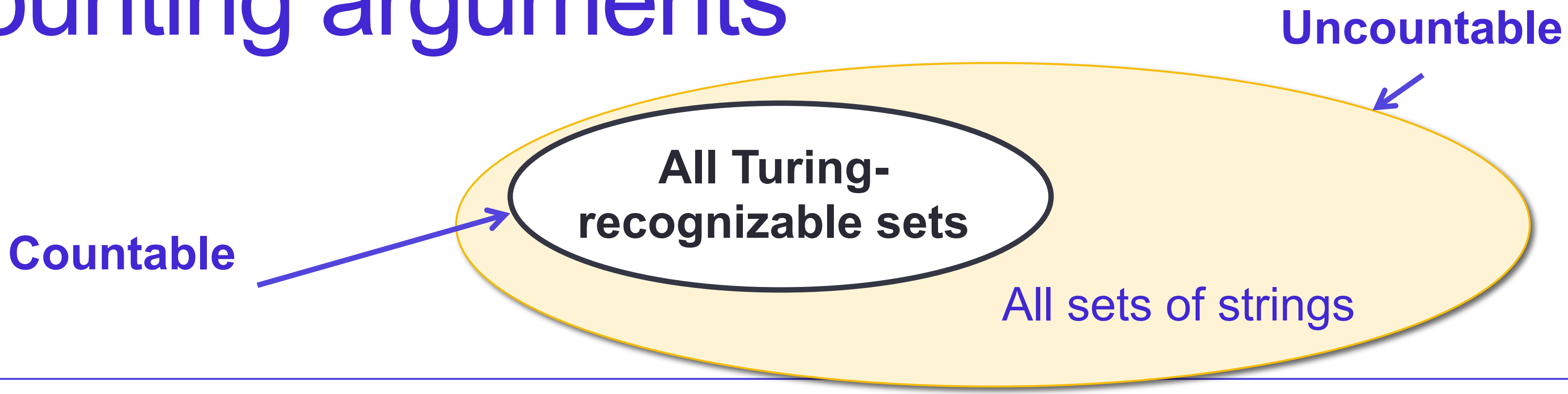
Sets A and B have the **same size** $|A| = |B|$ means there is a function between them that is both one-to-one and onto.

What type of elements are in the set $P(\Sigma^*)$?

A. Strings
B. Regular expressions
C. Languages
D. Sets of regular expressions
E. I don't know

Sipser p. 202-204

Counting arguments



Why is the set of Turing-recognizable languages **countable**?

A. It's equal to the set of all TMs, which we showed is countable.

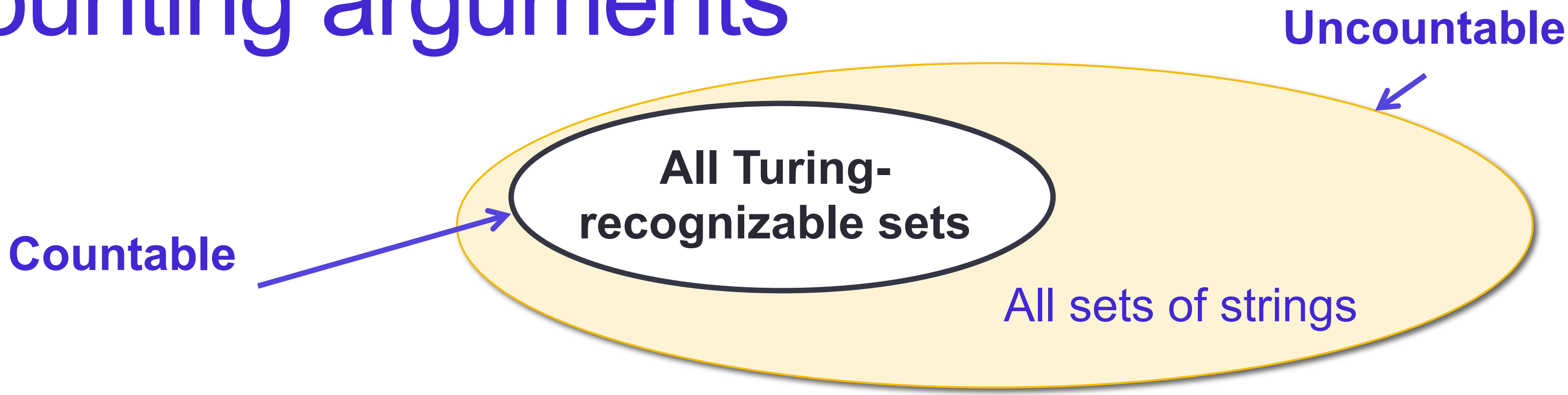
B. It's a subset of the set of all TMs, which we showed is countable.

C. Each Turing-recognizable language is associated with a TM, so there can be no more Turing-recognizable languages than TMs.

D. More than one of the above.

E. I don't know.

Counting arguments



Is the set of Turing-decidable sets countable?

Satisfied?

- Maybe not …

- What's a specific example of a language that is not Turing-recognizable? or not Turing-decidable?

- *Idea: consider a set that, were it to be Turing-decidable, would have to "talk" about itself, and contradict itself!*

A_{TM}

Recall $A_{DFA} = \{\langle B,w\rangle \mid B \text{ is a DFA and } w \text{ is in } L(B)\}$

$A_{TM} = \{\langle M,w\rangle \mid M \text{ is a TM and } w \text{ is in } L(M)\}$

What is A_{TM} ?

A. A Turing machine whose input is codes of TMs and strings.
B. A set of pairs of TMs and strings.
C. A set of strings that encode TMs and strings.
D. Not well defined.
E. I don't know.

A_{TM}

A_{TM} = {<M,w> | M is a TM and w is in L(M) }

Define the TM N = "On input <M,w>:

1. Simulate M on w.
2. If M accepts, accept. If M rejects, reject."

Which of the following statements is true?

A. N decides A_{TM}

B. N recognizes A_{TM}

C. N always halts

D. More than one of the above.

E. I don't know

A_{TM}

A_{DFA} = {<B,w> | B is a DFA and w is in L(B) }

Decider for this set simulates arbitrary DFA

A_{TM} = {<M,w> | M is a TM and w is in L(M) }

Decider for this set simulates arbitrary TMs.

What happens when it simulates itself?

Diagonalization proof: A_{TM} not decidable *Sipser p. 207*

Assume, towards a contradiction, that it is.

Call $\mathbf{M_{ATM}}$ the decider for A_{TM}:

$\mathbf{M_{ATM} \neq N}$

For every TM M and every string w,

- Computation of M_{ATM} on <M,w> halts and accepts if w is in L(M).
- Computation of M_{ATM} on <M,w> halts and rejects if w is not in L(M).

Diagonalization proof: A_{TM} not decidable *Sipser 4.11*

Assume, towards a contradiction, that M_{ATM} decides A_{TM}

Define the TM D = "On input <M>:

1. Run M_{ATM} on <M, <M>>.
2. If M_{ATM} accepts, reject; if M_{ATM} rejects, accept."

Diagonalization proof: A_{TM} not decidable *Sipser 4.11*

Assume, towards a contradiction, that M_{ATM} decides A_{TM}

Define the TM D = "On input <M>:

1. Run M_{ATM} on <M, <M>>.
2. If M_{ATM} accepts, reject; if M_{ATM} rejects, accept."

Which of the following computations halt?

A. Computation of D on <X> where X is TM which accepts string w if first character is 0 and loops otherwise.
B. Computation of D on <Y> where Y is TM with L(Y) =Σ*
C. Computation of D on <D>
D. All of the above.

Diagonalization proof: A_{TM} not decidable *Sipser 4.11*

Assume, towards a contradiction, that M_{ATM} decides A_{TM}

Define the TM D = "On input <M>:

1. Run M_{ATM} on <M, <M>>.
2. If M_{ATM} accepts, reject; if M_{ATM} rejects, accept."

Consider **running D on input <D>**. Because D is a decider:

- **either computation halts and accepts** …
- **or computation halts and rejects** …

Diagonalization proof: A_{TM} not decidable *Sipser 4.11*

Assume, t

Def

1. R

2. If

Diagonalization???

Self-reference

"Is <D> an element of L(D)?"

Consider **running** **on input <D>**. Because D is a decider:

- **either comp** **ation halts and accepts** …
- **or computation halts and rejects** …

A_{TM}

- Recognizable
- Not decidable

Fact (from discussion section): A language is decidable iff it and its complement are both recgonizable.

Corollary: The complement of A_{TM} is **unrecognizable**.

Do we have to diagonalize?

- Next time (after exam): comparing difficulty of problems.

Next time: Exam 2

Bring ID, note card

Check seat maps before coming to class