

HW#3, Solution

PROBLEM 2 (Universal gates)

For showing if a set of gates is universal or not, we need to check if we can build the AND, OR and NOT functionality using the set.

1. {XOR, NOT}, is **not universal**.
 - NOT : $X \text{ xor } 1 = X' = \text{NOT } X$
 - AND : AND gate cannot be built using only {XOR, NOT}
 - OR : OR gate cannot be built using only {XOR, NOT}
2. {NAND, NOT}, is **universal**.
 - NOT : is given
 - AND : $(X \text{ NAND } Y)' = X \text{ AND } Y$
 - OR : $(X' \text{ NAND } Y') = (X'Y')' = X + Y = X \text{ OR } Y$
3. {XNOR}, is **not universal**.
 - NOT : $X \text{ xnor } 0 = X' = \text{NOT } X$
 - AND : AND gate cannot be built using only {XNOR}
 - OR : OR gate cannot be build using only {XNOR}
4. {f(a, b)}, where $f(a, b) = a' + b'$, is **universal**.
 - NOT : $f(X, 1) = X' + 0 = X' = \text{NOT } X$
 - AND : $f(X, Y)' = (X' + Y')' = XY = X \text{ AND } Y$
 - OR : $f(X', Y') = (X + Y) = X \text{ OR } Y$
5. {f(a, b, c)}, where $f(a, b, c) = (a + b)(b + c)(a' + b' + c')$, is **universal**.
 - NOT : $f(a, 1, 1) = (a + 1)(1 + 1)(a' + 0 + 0) = a' = \text{NOT } a$
 - AND : $f(a, 1, c)' = (a + 1)(1 + c)(a' + 0 + c') = ac = a \text{ AND } c$
 - OR : $f(a', 1, c)' = (a' + 1)(1 + c')(a + 0 + c) = a + c = a \text{ OR } c$
6. {f(a, b, c)}, where $f(a, b, c) = ab + ac + bc$, is **not universal**.
 - NOT : NOT gate cannot be build with given function
 - AND : $f(a, b, 0) = ab + a0 + b0 = ab = a \text{ AND } b$
 - OR : $f(a, b, 1) = ab + a1 + b1 = a + b = a \text{ OR } b$

PROBLEM 3 (Other types of gates)

I. simplify expressions

$$1. f(a,b) = a \oplus (a + b') \oplus ab \oplus a' b$$

We can simplify the equation using shannon theorem:

shannon formula : $f(a,b) = a.f(1,b) + a'f(0,b)$

$$f(1,b) = 1 \oplus (1 + b') \oplus 1b \oplus 0 b$$

$$= 1 \oplus 1 \oplus b \oplus 0$$

$$\begin{aligned}
&= 0 \oplus b \\
&= b \\
f(0,b) &= 0 \oplus (0 + b') \oplus 0b \oplus 1b \\
&= 0 \oplus b' \oplus 0 \oplus b = b \oplus b' = 1
\end{aligned}$$

substituting in shannon formula :

$$\begin{aligned}
f(a,b) &= a.f(1,b) + a'.f(0,b) = ab + a' \\
&= ab + a' + b && //consensus \\
&= b(a+1) + a' \\
&= b + a'
\end{aligned}$$

$$2. f(a, b, c) = abc \oplus (a + b' + c) \oplus (b + c)b \oplus a' b \oplus (b + c')$$

We can simplify the equation using shannon theorem:

Shannon formula: $f(a,b,c) = b.f(a,1,c) + b'.f(a,0,c)$

$$\begin{aligned}
f(a,1,c) &= abc \oplus (a + b' + c) \oplus (b + c)b \oplus a' b \oplus (b + c') \\
&= a1c \oplus (a + 0 + c) \oplus (1 + c)1 \oplus a' 1 \oplus (1 + c') \\
&= ac \oplus (a + c) \oplus 1 \oplus a' \oplus 1 \\
&= ac \oplus (a + c) \oplus a' \\
&= a(1c \oplus (1 + c) \oplus 0) + a'(0c \oplus (0 + c) \oplus 1) \\
&= a(c \oplus 1 \oplus 0) + a'(0 \oplus c \oplus 1) \\
&= a(c \oplus 1) + a'(c \oplus 1) \\
&= a(c') + a'(c') \\
&= ac' + a'c' \\
&= c'(a + a') \\
&= c'
\end{aligned}$$

$$//f(a,c) = a f(1,c) + a' f(0,c)$$

$$\begin{aligned}
f(a,0,c) &= a0c \oplus (a + 1 + c) \oplus (0 + c)0 \oplus a' 0 \oplus (0 + c') \\
&= 0 \oplus 1 \oplus 0 \oplus 0 \oplus c' \\
&= 1 \oplus c' \\
&= c
\end{aligned}$$

$$\begin{aligned}
f(a,b,c) &= b.f(a,1,c) + b'.f(a,0,c) \\
&= bc' + b'c \\
&= b \oplus c
\end{aligned}$$

II. Prove or disprove the following equalities.

$$1. a \oplus bc = (a \oplus b)(a \oplus c)$$

we can show the equality does not hold by a counterexample :

if $b=1, c=0$

$$\text{LHS} = a \oplus bc = a \oplus 0 = a$$

$$\text{RHS} = (a \oplus b)(a \oplus c) = (a \oplus 1)(a \oplus 0) = (a') (a) = 0$$

$$\text{RHS} \neq \text{LHS}.$$

$$\text{LHS} = a(bc) + a'(bc)$$

$$\begin{aligned} \text{RHS} &= (ab' + a'b) (ac' + a'c) \\ &= ab'c' + a'bc \\ \text{RHS} &\neq \text{LHS}. \end{aligned}$$

So, we cannot distribute XOR over AND.

$$2. a(b \oplus c) = ab \oplus ac$$

$$\begin{aligned} \text{RHS} &= ab \oplus ac = (ab)'(ac) + (ab)(ac)' \\ &= (a'+b')(ac) + (ab)(a'+c)' \\ &= a'ac + b'ac + aba' + abc' \\ &= 0 + b'ac + 0 + abc' \\ &= abc' + ab'c \\ &= a(bc' + b'c) \\ &= a(b \oplus c) \\ &= \text{LHS} \end{aligned}$$

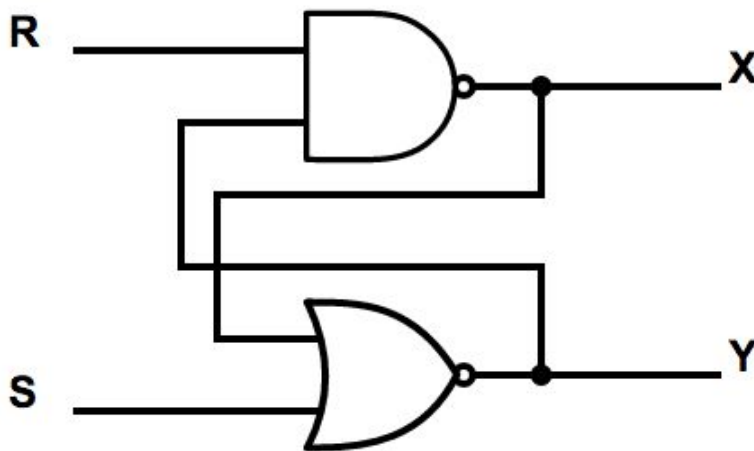
So, we can distribute AND over XOR.

PROBLEM 4 (Latch)

A. Replacing the top NOR gate with a NAND gate

$$X(t+1) = (R Y(t))'$$

$$Y(t+1) = (S + X(t))'$$



id	X(t)	Y(t)	S	R	X(t1)	Y(t1)	X(t2)	Y(t2)	X(t3)	Y(t3)
0	0	0	0	0	1	1	1	0	1	0
1	0	0	0	1	1	1	0	0	1	1

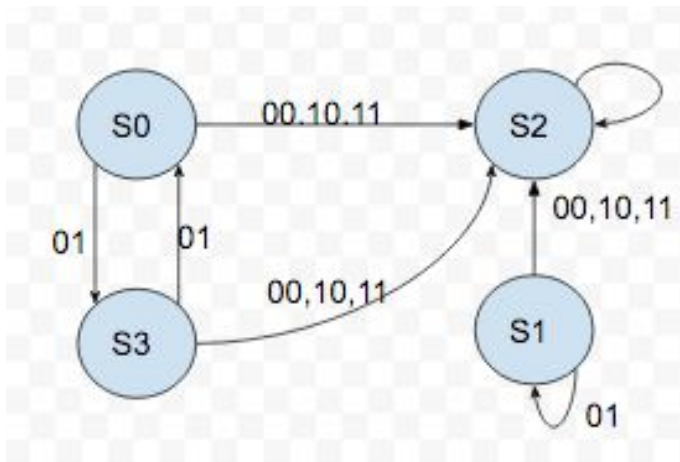
2	0	0	1	0	1	0	1	0	1	0
3	0	0	1	1	1	0	1	0	1	0
4	0	1	0	0	1	1	1	0	1	0
5	0	1	0	1	0	1	0	1	0	1
6	0	1	1	0	1	0	1	0	1	0
7	0	1	1	1	0	0	1	0	1	0
8	1	0	0	0	1	0	1	0	1	0
9	1	0	0	1	1	0	1	0	1	0
10	1	0	1	0	1	0	1	0	1	0
11	1	0	1	1	1	0	1	0	1	0
12	1	1	0	0	1	0	1	0	1	0
13	1	1	0	1	0	0	1	1	0	0
14	1	1	1	0	1	0	1	0	1	0
15	1	1	1	1	0	0	1	0	1	0

using S0 = 00, S1= 01, S2 = 10, S3=11

id	Current state	S	R	state (t1)	state X(t2)	state X(t3)
0	S0	0	0	S3	S2	S2
1	S0	0	1	S3	S0	S3
2	S0	1	0	S2	S2	S2
3	S0	1	1	S2	S2	S2
4	S1	0	0	S3	S2	S2
5	S1	0	1	S1	S1	S1
6	S1	1	0	S2	S2	S2
7	S1	1	1	S0	S2	S2

8	S2	0	0	S2	S2	S2
9	S2	0	1	S2	S2	S2
10	S2	1	0	S2	S2	S2
11	S2	1	1	S2	S2	S2
12	S3	0	0	S2	S2	S2
13	S3	0	1	S0	S3	S0
14	S3	1	0	S2	S2	S2
15	S3	1	1	S0	S2	S2

state diagram:



we can avoid the input "01" and ignore states S0 and S3 to prevent the unstable loop.

next state:

current state/ SR	00	01	11	10
01 (s1)	s2	x	s2	s2
10 (s2)	s2	x	s2	s2

current state/ SR	00	01	11	10
01 (s1)	10	x	10	10
10 (s2)	10	x	10	10

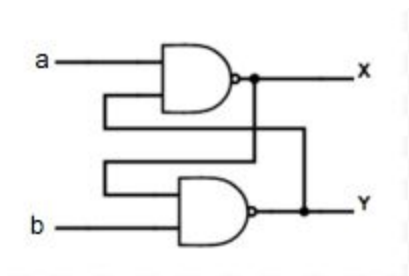
$$X(t+1) = 1$$

current state/ SR	00	01	11	10
01 (s1)	1	x	1	1
10 (s2)	1	x	1	1

$$y(t+1) = 0$$

current state/ SR	00	01	11	10
01 (s1)	0	x	0	0
10 (s2)	0	x	0	0

B. Replacing both NOR gates with NAND gates



$$X(t+1) = (a Y(t))'$$

$$Y(t+1) = (b X(t))'$$

id	X(t)	Y(t)	b	a	X(t1)	Y(t1)	X(t2)	Y(t2)	X(t3)	Y(t3)
0	0	0	0	0	1	1	1	1	1	1
1	0	0	0	1	1	1	0	1	0	1

2	0	0	1	0	1	1	1	0	1	0
3	0	0	1	1	1	1	0	0	1	1
4	0	1	0	0	1	1	1	1	1	1
5	0	1	0	1	0	1	0	1	0	1
6	0	1	1	0	1	1	1	0	1	0
7	0	1	1	1	0	1	0	1	0	1
8	1	0	0	0	1	1	1	1	1	1
9	1	0	0	1	1	1	0	1	0	1
10	1	0	1	0	1	0	1	0	1	0
11	1	0	1	1	1	0	1	0	1	0
12	1	1	0	0	1	1	1	1	1	1
13	1	1	0	1	0	1	0	1	0	1
14	1	1	1	0	1	0	1	0	1	0
15	1	1	1	1	0	0	1	1	0	0

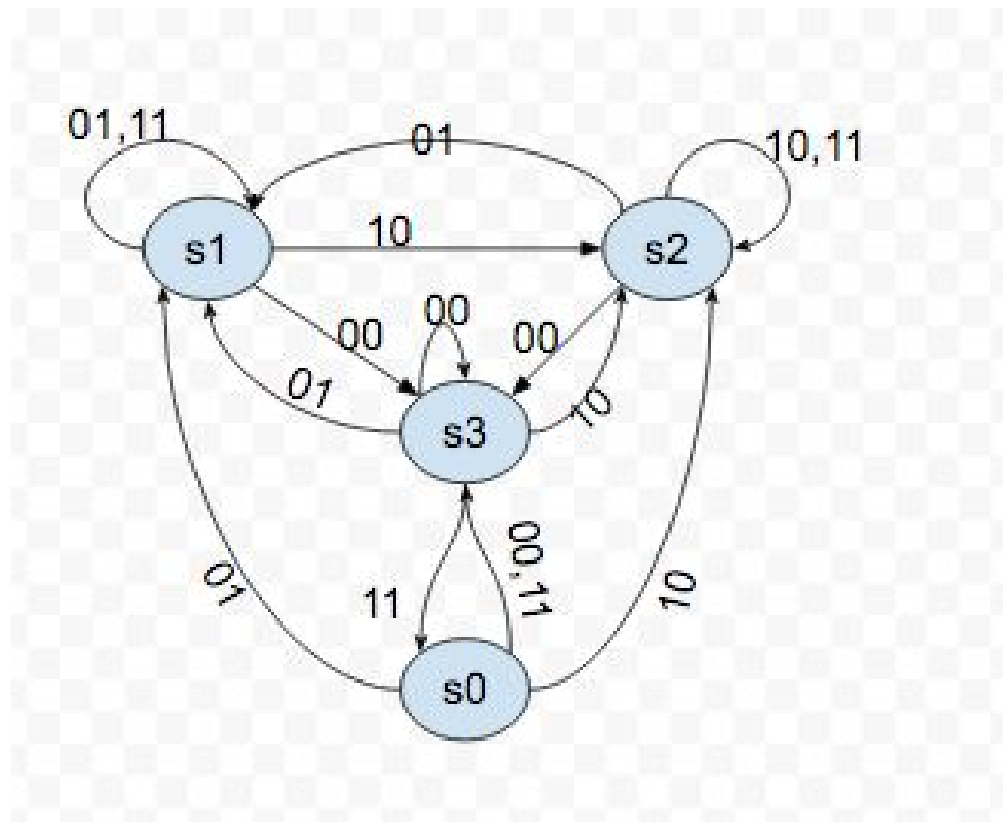
using S0 = 00, S1= 01, S2 = 10, S3=11

id	current state	b	a	state (t1)	state (t2)	state (t3)
0	S0	0	0	S3	S3	S3
1	S0	0	1	S3	S1	S1
2	S0	1	0	S3	S2	S2
3	S0	1	1	S3	S0	S3
4	S1	0	0	S3	S3	S3
5	S1	0	1	S1	S1	S1
6	S1	1	0	S3	S2	S2
7	S1	1	1	S1	S1	S1

8	S2	0	0	S3	S3	S3
9	S2	0	1	S3	S1	S1
10	S2	1	0	S2	S2	S2
11	S2	1	1	S2	S2	S2
12	S3	0	0	S3	S3	S3
13	S3	0	1	S1	S1	S1
14	S3	1	0	S2	S2	S2
15	S3	1	1	S0	S3	S0

state diagram:

You can see that the diagram has the same structure as the NOR SR latch from the lecture slide.



00 is the invalid input in this case, by replacing it with dont cares we will have just two state S1 and S2:

$X(t) \backslash ba$	00	01	11	10
0	x	0	0	1
1	x	0	1	1

$$X(t+1) = a' + bX(t)$$

rubric:

- question 2: 18 pts
 - 3pts each
 - give full credit to final answer
 - if final answer in wrong, give 2 points for good reasoning
- 3 I . 1 10 pts for attempt
- 3 I . 2 10 pts for attempt
- 3 II . 1 10 pts for attempt
- 3 || . 2 16 pts (-10 if its wrong)
- 4 I . 1 8pts race condition in table
- 4 I . 2 5pts for attempt
- 4 I . 3 5pts for attempt
- 4 II . 1 8pts race condition in table
- 4 || 2 5pts for attempt
- 4 || 3 5pts for attempt