

CSE20

Lecture 2: Number Systems: Binary Numbers, Gray Code, and Negative Numbers

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Number Systems

1. Introduction
2. Binary Numbers
3. Gray code
4. Negative Numbers
5. Residual Numbers

2. Binary Numbers

b_2	b_1	b_0	Value
0	0	0	0
0	0	1	1
0	1	0	2
0	1	1	3
1	0	0	4
1	0	1	5
1	1	0	6
1	1	1	7

Examples:

$$\begin{array}{r}
 \mathbf{3 + 5 = 8} \\
 \begin{array}{cccc}
 & 8 & 4 & 2 & 1 \\
 0 & 0 & 1 & 1 & (3) \\
 + & 0 & 1 & 0 & 1 & (5) \\
 \hline
 1 & 0 & 0 & 0 & (8)
 \end{array}
 \end{array}$$

$$\begin{array}{r}
 \mathbf{3 + 6 = 9} \\
 \begin{array}{cccc}
 & 8 & 4 & 2 & 1 \\
 0 & 0 & 1 & 1 & (3) \\
 + & 0 & 1 & 1 & 0 & (6) \\
 \hline
 1 & 0 & 0 & 1 & (9)
 \end{array}
 \end{array}$$

This is a non-redundant number system

2. Binary Cont.

a	b	Carry	Sum
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

RULE:

$$2 \times \text{Carry} + \text{Sum} = a + b + c$$

id	a	b	c	Carry	Sum
0	0	0	0	0	0
1	0	0	1	0	1
2	0	1	0	0	1
3	0	1	1	1	0
4	1	0	0	0	1
5	1	0	1	1	0
6	1	1	0	1	0
7	1	1	1	1	1

$$2 \cdot 0 + 0 = 0 \quad 0 \quad 0 \quad \text{id 0}$$

$$2 \cdot 0 + 1 = 0 \quad 0 \quad 1 \quad \text{id 1}$$

$$2 \cdot 1 + 0 = 1 \quad 1 \quad 0 \quad \text{id 6}$$

$$2 \cdot 1 + 1 = 1 \quad 1 \quad 1 \quad \text{id 7}$$

3. Gray Code

ld	b3	b2	b1	b0
0	0	0	0	0
1	0	0	0	1
2	0	0	1	1
3	0	0	1	0
4	0	1	1	0
5	0	1	1	1
6	0	1	0	1
7	0	1	0	0
8	1	1	0	0
9	1	1	0	1
10	1	1	1	1
11	1	1	1	0
12	1	0	1	0
13	1	0	1	1
14	1	0	0	1
15	1	0	0	0

} reflection

Low power (reliability) when the numbers are consecutive in series.

The idea is to only change ONE bit at a time.

e.g. addresses, analog signals

NOTE: Not for arithmetic operations (the rule is too complicated)

4. Negative Numbers

Given a positive integer x , represent the negative integer $-x$ in (b_{n-1}, \dots, b_0)

(i) Signed bit system

$b_{n-1}=1$: negative, $(b_{n-2}, \dots, b_0)=x$.

(ii) One's Complement

Present $2^n - 1 - x$ in binary.

(iii) Two's Complement

Present $2^n - x$ in binary. Ignore bit b_n .

4. Negative Numbers

NOTE: Back to binary system

Deriving One's and Two's

Reverse Derivation

(i) Signed bit - x
b3: negative

(ii) One's Complement
 $2^n - 1 - x$

(iii) Two's Complement
 $2^n - x$

n is the number of bits (in this case n=4)

Use the above formulas to solve for x when number is negative

One's Complement Two's Complement

$$8 = 16 - 1 - x$$

$$8 = 16 - x$$

$$9 = 16 - 1 - x$$

id	b ₃	b ₂	b ₁	b ₀	Signed	One's	Two's	Two's	(b ₄)	b ₃	b ₂	b ₁	b ₀
0	0	0	0	0	0	0	0	7	1	0	1	1	1
1	0	0	0	1	1	1	1	6	1	0	1	1	0
2	0	0	1	0	2	2	2	5	1	0	1	0	1
3	0	0	1	1	3	3	3	4	1	0	1	0	0
4	0	1	0	0	4	4	4	3	1	0	0	1	1
5	0	1	0	1	5	5	5	2	1	0	0	1	0
6	0	1	1	0	6	6	6	1	1	0	0	0	1
7	0	1	1	1	7	7	7	0	1	0	0	0	0
8	1	0	0	0	-0	-7	-8	-1	0	1	1	1	1
9	1	0	0	1	-1	-6	-7	-2	0	1	1	1	0
10	1	0	1	0	-2	-5	-6	-3	0	1	1	0	1
11	1	0	1	1	-3	-4	-5	-4	0	1	1	0	0
12	1	1	0	0	-4	-3	-4	-5	0	1	0	1	1
13	1	1	0	1	-5	-2	-3	-6	0	1	0	1	0
14	1	1	1	0	-6	-1	-2	-7	0	1	0	0	1
15	1	1	1	1	-7	-0	-1	-8	0	1	0	0	0