

CSE 222

Graduate Networking

Winter 2001

Lecture 2: Internet Design

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Overview

- Current Internet architecture at a high level
 - Should seem like a review
- Clark Internet design philosophy paper

Network Components

- Hosts: Computer, PDA, light switch...
- Link: Transmission media
 - ♦ Wired or wireless
 - ♦ Broadcast or switched
- Switch: Crossroads, move bits between links
 - ♦ Packet switching: stateless store and forward (Internet)
 - ♦ Circuit switching: stateful, cut through (POTS)
- Protocol: Agreement on how information is to be formatted and transmitted

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Internetwork

- An internetwork is a network of networks
 - ♦ Could do this in more than one way
 - ♦ IP is just one, POTS is another
- Routers (gateways) move packets between networks
 - ♦ Packet switched or circuit switched
 - ♦ Could also translate
- The Internet
 - ♦ Routers move IP packets from one network to another
 - » IP "dial tone"
 - ♦ Minimal requirements on underlying network
 - ♦ Can use almost any potential network or link layer
 - » Modem, Ethernet, token ring, ATM, ADSL, cell phone, cable...

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How do we communicate in the Internet?

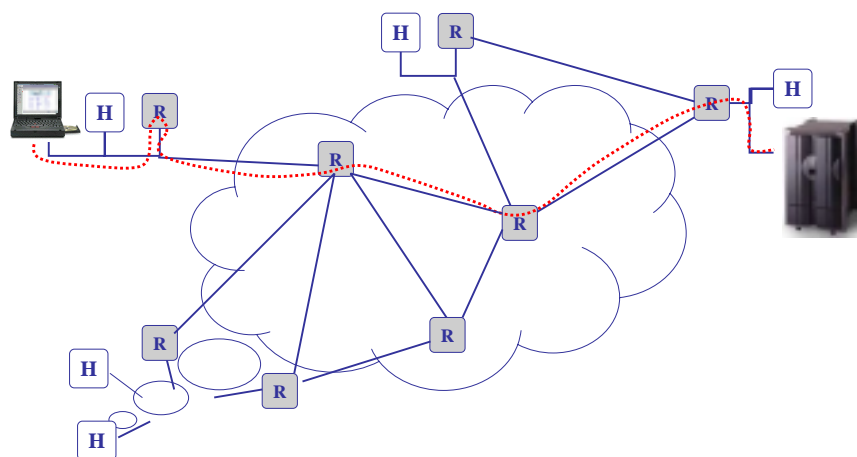
- Need to specify how to handle
 - ♦ Addressing
 - ♦ Routing
 - ♦ Service models
 - ♦ Failures
 - ♦ Management
 - ♦ Accounting
 - ♦ Etc.

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Using the Web



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Addressing

- Domain name: www.ucsd.edu
 - ◆ Global (across networks), human readable
- IP: 132.239.50.184
 - ◆ Global, what is actually used in routing
- Ethernet: 00-90-27-BD-BC-F7
 - ◆ Local, used within a particular network
- When making a connection...
 - ◆ Domain name is converted to an IP address
 - ◆ IP address used in packets
 - ◆ Packets sent on Ethernet to Ether address of the gateway

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Routing

- When a gateway receives a packet, it needs to decide what to do with it
 - ◆ If it is to itself, deliver it
 - ◆ Otherwise, figure out which link to send it to (packet switching)
- Routing tables
 - ◆ Map IP prefixes to output links using forwarding tables
 - ◆ Per-hop, router does not need to know final destination
 - ◆ Automatically updated in response to failures, changes
- Internet routing done at two levels
 - ◆ Within a domain, across domains
- Much more later...

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Service Models

- What can you expect from a particular protocol?
 - ♦ Unreliable (UDP)
 - ♦ Reliable (TCP)
 - ♦ Timely (RTP)
 - » No guarantees on Internet, though!

Failures

- Internet delivery is best effort – no guarantees
- Routers go down
 - ♦ Use a different route (how do routers figure out alternates?)
- Data corruption
 - ♦ Can happen in many places – any piece of hardware a packet travels over (link, router I/O bus, router memory, etc.)
 - ♦ Hardware CRCs, software checksums to detect
- Reordering
 - ♦ Sequence numbers (how large?)
 - ♦ Buffer packets at end points (how much?)
- Drops
 - ♦ Router link buffers can fill up (congestion)
 - ♦ Need to drop (which ones?)

Clark: Design Philosophy of the DARPA Internet Protocols

- Great paper
 - ♦ Not many papers explaining the motivation and reasoning that went into the design of systems that we take for granted
- Note that this was written 15 years after the project began
 - ♦ And the paper itself is already 13 years old!
- The setting
 - ♦ Multiple research and military networks
 - ♦ How do we connect them so that they can talk to each other?
 - ♦ Hard to imagine, but this was *before* LANs

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Before going into the text...

- Meta-points
 - ♦ The Internet, TCP/IP, etc., were designed and engineered – there is no natural law that says the Internet had to look the way it does now
 - » It could well have been done differently
 - ♦ “The Internet”, “TCP/IP”, etc., continually evolve
 - » The Internet today is not the same Internet as 1988, 1973
 - » TCP/IP have changed considerably over the years
 - » We’re using IPv4, with IPv6 in the works
 - ♦ Seemingly straightforward decisions can have very subtle correctness and performance implications
 - » EOL and PSH
 - » Acknowledging bytes vs. packets (debates still rage...)

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Primary Goal

- “Effective technique for multiplexed utilization of existing interconnected networks”
 - ♦ Minimal assumptions about underlying networks
 - » No support for broadcast, multicast, real-time, reliability, etc.
 - » Nets could support it, but not necessary
 - » Extra support could actually get in the way (X.25 example)
 - ♦ Packet switched, store and forward
 - » Matched application needs, nets already packet switched
 - » Enables fine-grained resource sharing
 - ♦ “Gateways” interconnect networks
 - » Routers today

Why is this hard?

- Every network characteristic is different between two arbitrary networks
 - ♦ Addressing
 - » Each network media has a different addressing scheme
 - ♦ Bandwidth
 - » Modems to terabits
 - ♦ Latency
 - » Seconds to nanoseconds
 - ♦ Packet size
 - » Dozens to thousands of bytes
 - ♦ Loss rates
 - » Differ by many orders of magnitude

Possibilities

- All nets communicate using a common format
 - Internet: IP over everything
 - To talk across networks, you send IP packets
 - Internal to a network, can use whatever you want
 - » Raw Ethernet, ATM, etc.
- Translate packets from one network to another
 - Convert Ethernet to ATM
 - Convert IP to OSI
 - X.25

Secondary Goals

- Survivability
- Multiple communication services
- Accommodate a variety of networks
- Distributed management of resources
- Cost-effective
- Minimal effort to attach hosts to net
- Resource accounting

- These are in order of priority -- a different priority ordering would likely result in a different design

Survivability

- Internet
 - ♦ Assume anything can fail between two end points
 - » Fate-sharing (state is lost only when end point is lost)
 - ♦ Designed to tolerate 1% drop rate
 - » High or low? Depends on the app, assumptions, goals, etc.
 - » 10% things break down
 - ♦ Routing is simpler, does not have to adapt to failure
 - » For a given datagram
- POTS (the *other* global network)
 - ♦ Ultra reliable switches
 - ♦ Self-healing
 - » Hardware switch over in the middle of a phone call

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Survivability Implications

- End points maintain all essential state
 - ♦ Routers are stateless ("soft state")
 - ♦ End points responsible for recovering from failures
- Host machines are trusted
 - ♦ Have to rely upon hosts to implement the protocols correctly
 - » For performance as well as correctness
 - ♦ Easy to be malicious
 - » Ex: source addresses (everything in an IP packet) are trusted (IP spoofing)
- Can be difficult to determine source of failures
 - ♦ Not much feedback from network back to end point
 - ♦ Makes performance optimizations more difficult

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Types of Service

- Bottom line: Best effort datagram service
- Building block
 - ♦ Unreliable (UDP)
 - ♦ Reliable service (TCP)
 - » Even two types of this: what are they?
 - ♦ Real-time
 - » Can you do real-time without help from network?
 - ♦ Multicast
- Reason why TCP and IP became separate protocols
 - ♦ IP basis for all other protocols
 - ♦ Originally were combined

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Network Variety

- Fundamental goal was to interconnect networks...
- Internet successful in part because its design met this goal extremely well
 - ♦ Topology: point-to-point, bus, ring, radio, satellite, etc.
 - ♦ Characteristics: modem to Tbit speeds, us to sec delays
- Does not mean that IP on a given network is efficient
 - ♦ X.25 supports reliable delivery
 - ♦ ATM uses 53 byte cells, poor fragmentation for IP packets
 - » IP gets layered on top of cells, more effort at end points

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Other Goals

- Distributed management
 - ♦ Different parts of network owned, controlled, managed by separate, Difficult to support, though
 - » Hard to do things across entities (e.g., optimize routes)
 - » Not sure what optimize means (to an ISP, customer?)
 - ♦ Problem in 88, more of a problem today
- Cost-effective (compared to?)
 - ♦ IP routers cheap compared to POTS switches
 - ♦ But higher drop rates, inefficient routing, end-to-end reliability all impose an overhead/inefficiency on the network
 - ♦ Still a source of religious debate

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Other Goals (2)

- Attachment cost
 - ♦ \$100M/year on protocol stacks by major OS vendors
 - » Is that cheap or expensive given 100s of millions of users?
 - ♦ Bugs still result in misbehaving hosts!
 - » Both correctness and performance
 - » There are tools that can tell you what OS, TCP/IP stack you are using based upon the bug signatures in implementations
- Accountability
 - ♦ Who pays for all of this?
 - ♦ What is the economic model of the future?
 - ♦ No one knows (plenty of ideas, though)
 - ♦ Another RD (religious debate)

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Discussion

- What surprised you most about the Internet design and architecture when you read the paper?
- What is different about the Internet today than in '88 when Clark wrote this paper?
- What are limitations to the Internet design?

- When thinking about the Internet, imagine it as the culmination of many iterations of design and trial and error – not as an architecture that was preordained

For Next Time...

- Send me email if you aren't on the list yet
- Read Saltzer84 and Clark90 papers
- Browse Chapter 3, read Chapter 4
- Optional
 - Zimmerman80 – OSI reference model
 - Shenker95 – Should there be a new service model for the Internet to support multimedia applications?