

CSE 127

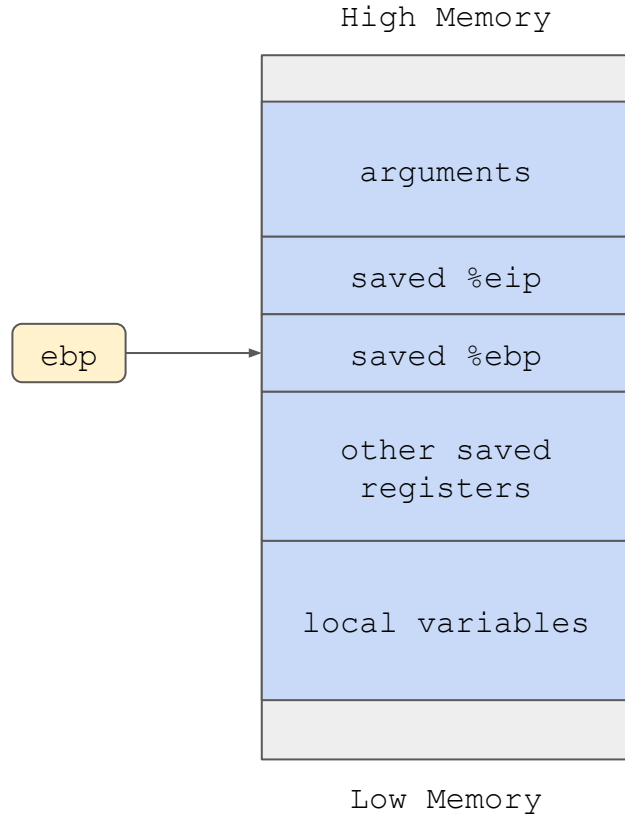
# Week 3 Discussion

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# Stack Layout (Revision)



`%eip` is a register pointing to the instruction that CPU will execute in next cycle

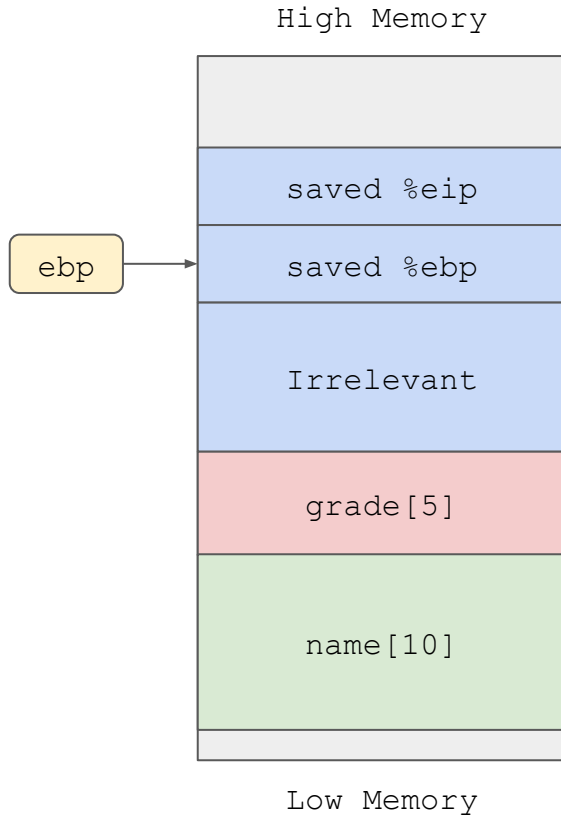
saved `%eip` references to a 4-byte address value stored on stack

saved `%eip` is stored on stack when a function call is made. It has the address of where to resume execution in the caller function

When a function returns, the saved `%eip` value will be popped into the register `%eip` → control will transfer to where saved `%eip` points to

return address == saved return address == saved `%eip`  
== `%ebp+4`

# Target 0



- How does `gets` copy `stdin` into `name`?
- When does `gets` stop copying?
- Will `gets` check the boundary between `name` and `grade`?

```
int _main(int argc, char *argv[])
{
    char grade[5];
    char name[10];

    strcpy(grade, "nil");

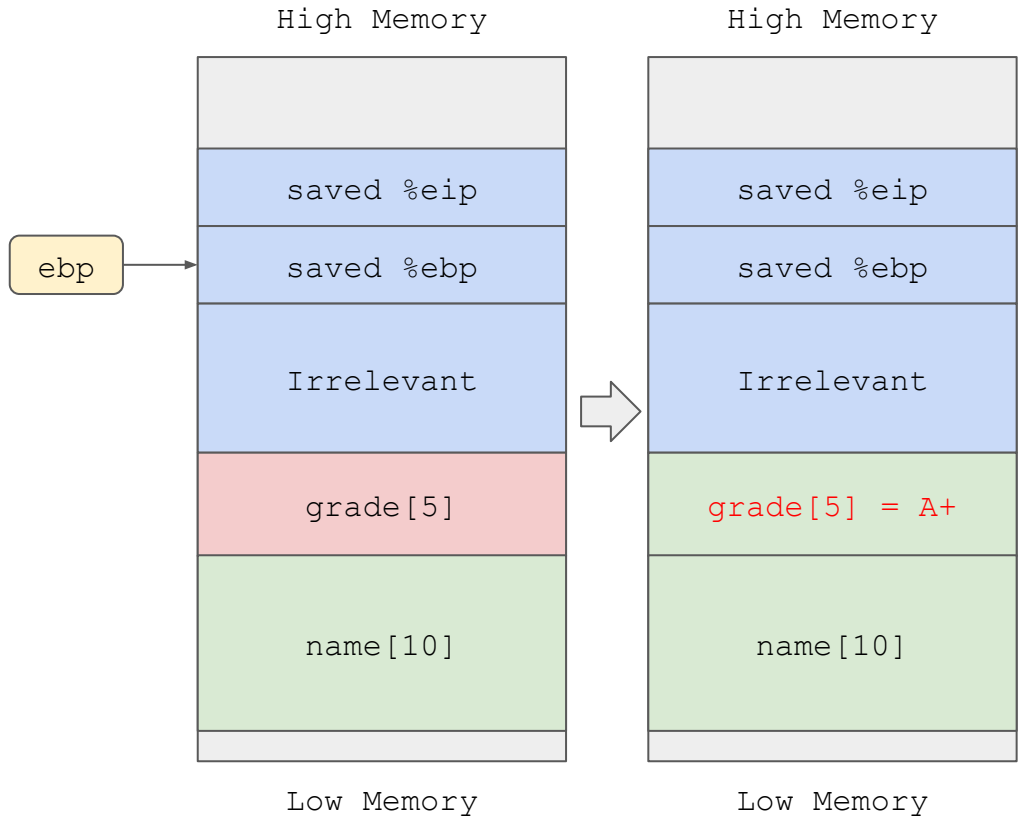
    gets(name);

    printf("Hi %s! Your grade is %s.\n", name,
grade);

    exit(0);
}
```

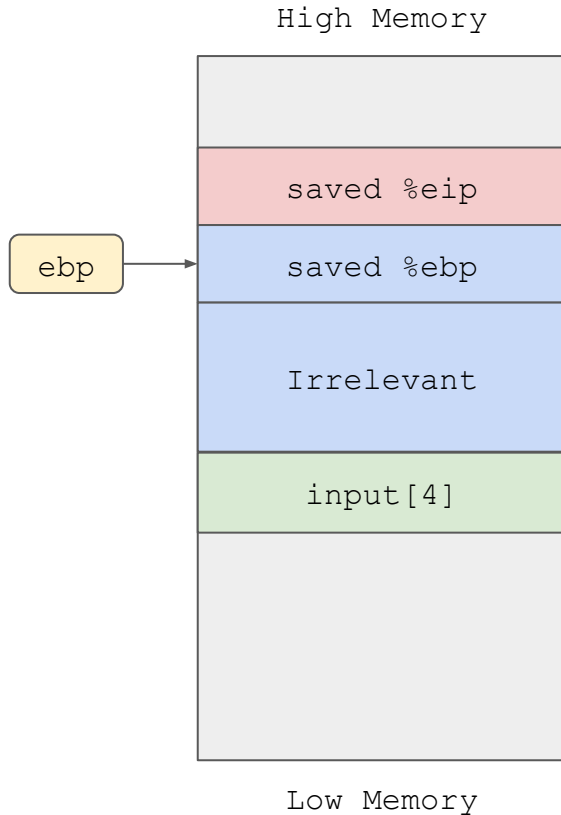
Be careful, when `name` is printed, no extra garbage content should be printed. Think about when `printf` stops prints given a pointer.

# Target 0



```
int _main(int argc, char *argv[])  
{  
    char grade[5];  
    char name[10];  
  
    strcpy(grade, "nil");  
  
    gets(name);  
  
    printf("Hi %s! Your grade is %s.\n", name,  
grade);  
  
    exit(0);  
}
```

# Target 1



- What is the distance between `input` and `saved %eip`?
- How do you fill the gap in between?
- Where do you want to redirect control to?

```

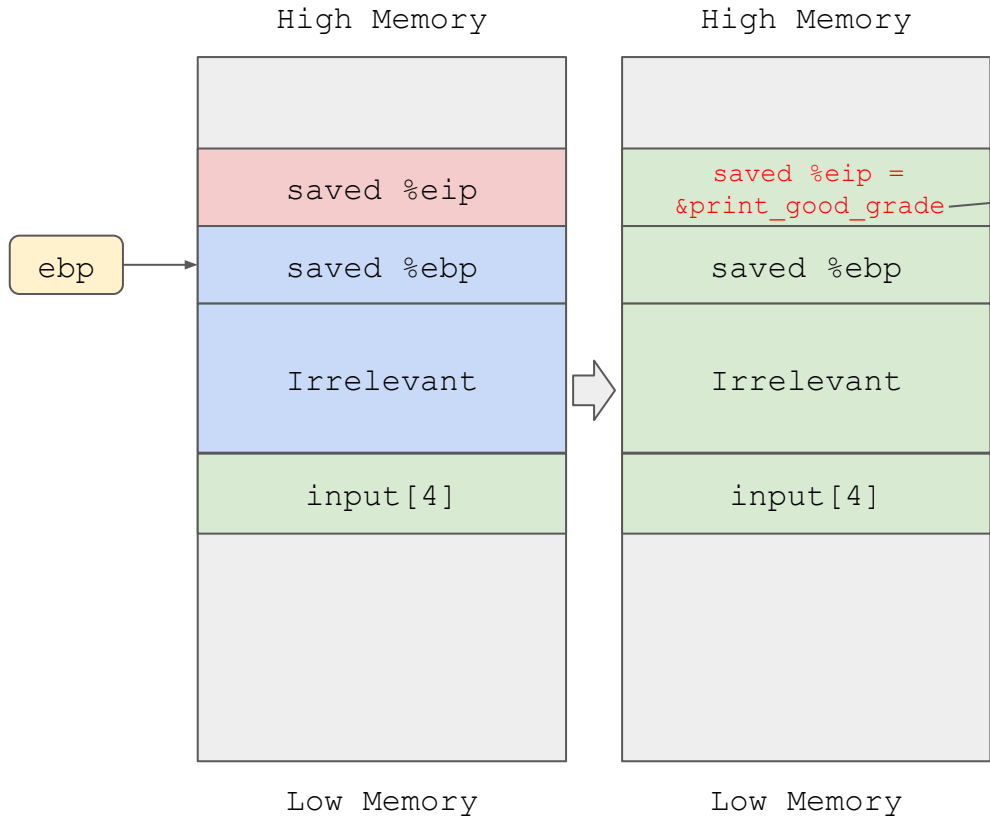
void print_bad_grade(void)
{
    puts("Your grade is nil.");
    exit(0);
}

void print_good_grade(void)
{
    puts("Your grade is perfect.");
    exit(1);
}

void vulnerable()
{
    char input[4];
    gets(input);
}

int _main()
{
    vulnerable();
    print_bad_grade();
    return 0;
}
  
```

# Target 1



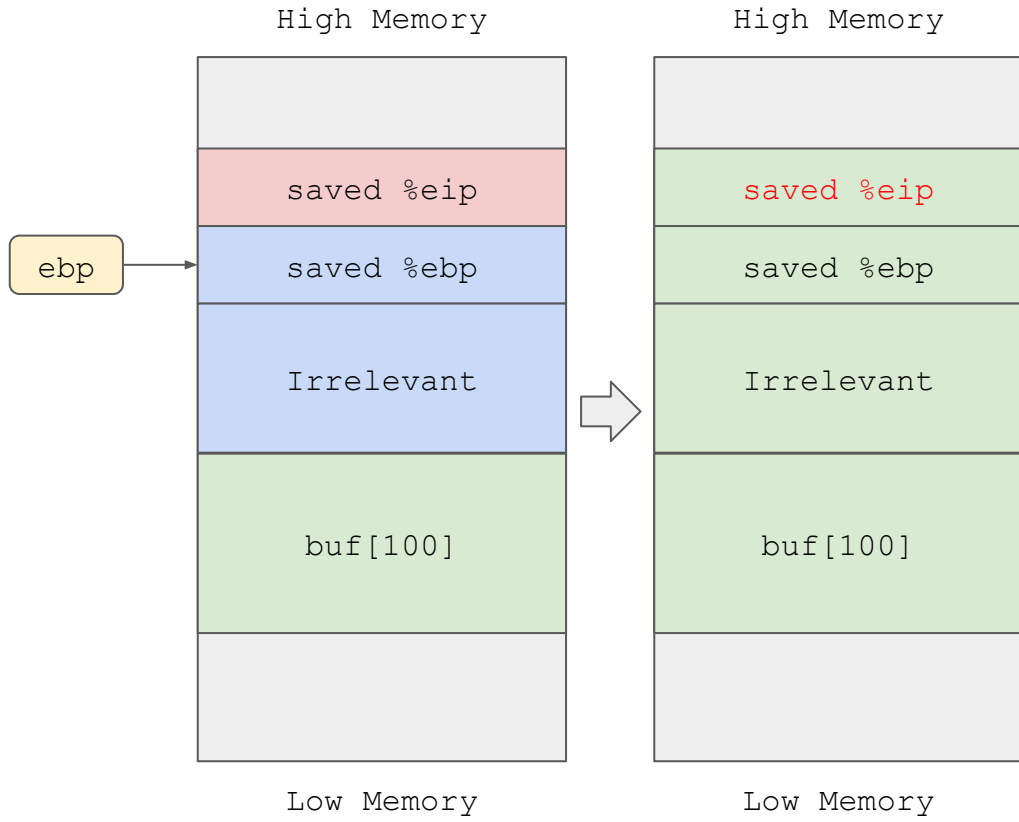
```
void print_bad_grade(void)
{
    puts("Your grade is nil.");
    exit(0);
}

void print_good_grade(void)
{
    puts("Your grade is perfect.");
    exit(1);
}

void vulnerable()
{
    char input[4];
    gets(input);
}

int _main()
{
    vulnerable();
    print_bad_grade();
    return 0;
}
```

# Target 2



```

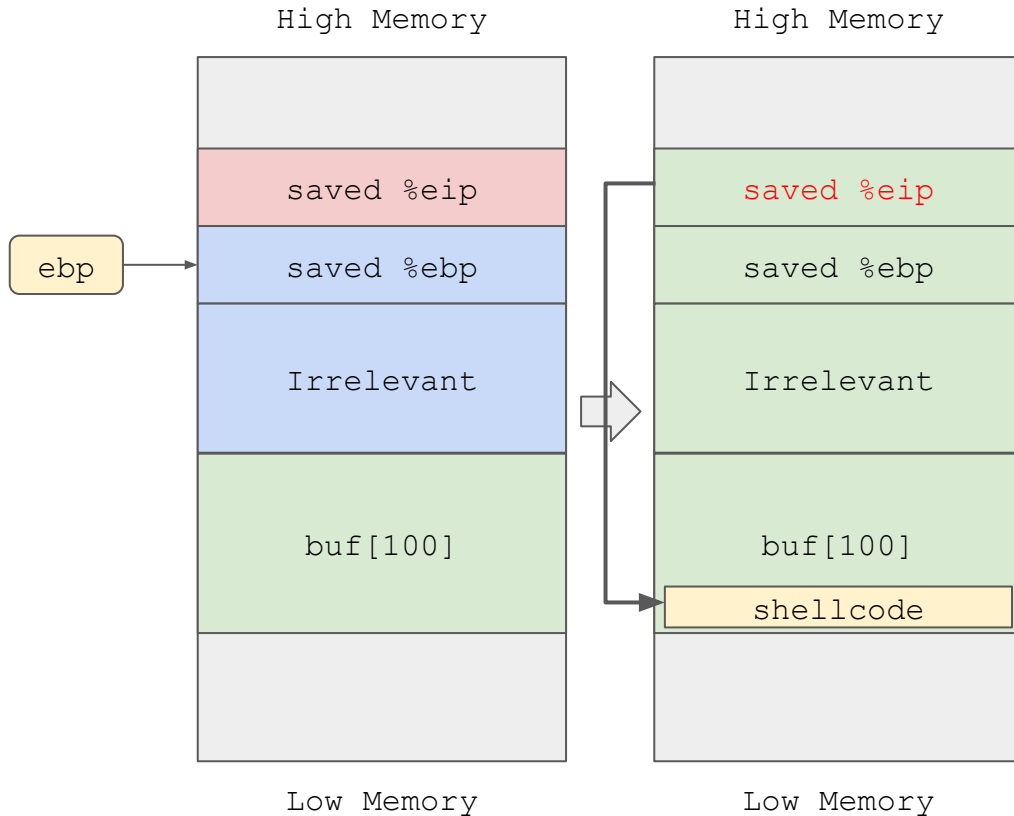
void vulnerable(char *arg)
{
    char buf[100];
    strcpy(buf, arg);
}

int _main(int argc, char **argv)
{
    if (argc != 2) {
        fprintf(stderr, "Error: need a command-line
argument\n");
        return 1;
    }
    vulnerable(argv[1]);
    return 0;
}

```

**Where do you want to transfer control to?**

# Target 2



```

void vulnerable(char *arg)
{
    char buf[100];
    strcpy(buf, arg);
}

int _main(int argc, char **argv)
{
    if (argc != 2) {
        fprintf(stderr, "Error: need a command-line
argument\n");
        return 1;
    }
    vulnerable(argv[1]);
    return 0;
}
  
```

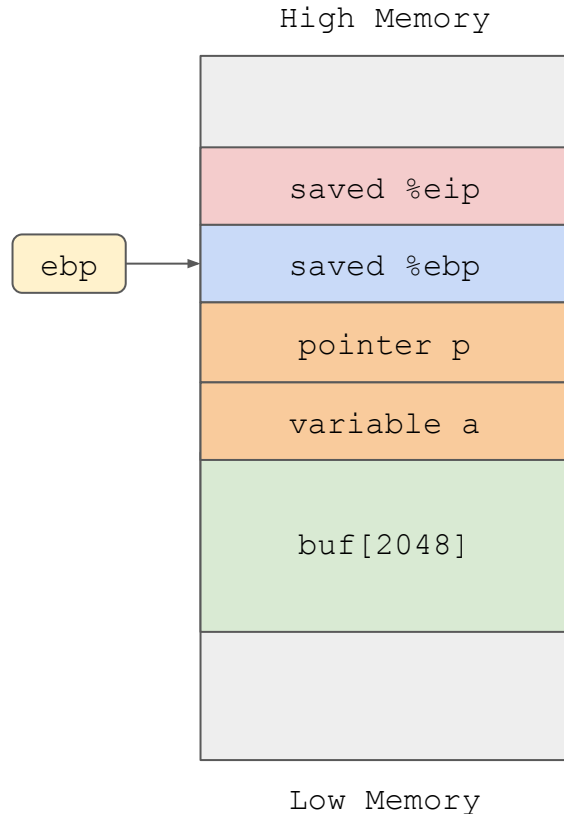
**Where do you want to transfer control to?**



# Shellcode - Simple examples and explanations

- <https://www.youtube.com/watch?v=1S0aBV-Waeo>
- <https://www.youtube.com/watch?v=hJ8lwyhqzD4>

## Target 3 (Hint)



```
void vulnerable(char *arg)
{
    int *p;
    int a;
    char buf[2048];

    strncpy(buf, arg, sizeof(buf) + 8);

    *p = a;
}
```

- Figure out what goes in 'p' and 'a'
- \*p = a is basically dereferencing 'a'
- Think about how dereferencing can be taken advantage of

## Target 4 (Hint)

- count: 32-bit unsigned integer
- What happens when I overflow an unsigned int?

```
int main() {  
    unsigned int a = pow(2, 30) + 127;  
    unsigned int b = pow(2, 31) + 100;  
  
    printf("%u\n", a * 4);  
    printf("%u\n", b * 2);  
}
```

```
void read_elements(FILE *f, int *buf, unsigned int count)  
{  
    unsigned int i;  
    for (i=0; i < count; i++) {  
        if (fread(&buf[i], sizeof(unsigned int), 1, f) < 1) {  
            break;  
        }  
    }  
}  
  
void read_file(char *name)  
{  
    FILE *f = fopen(name, "rb");  
    ....  
  
    unsigned int count;  
    fread(&count, sizeof(unsigned int), 1, f);  
  
    unsigned int *buf = alloca(count * sizeof(unsigned int));  
    ....  
  
    read_elements(f, buf, count);  
}
```

## Helpful links

- <https://devblogs.microsoft.com/oldnewthing/20050107-00/?p=36773>
- <https://www.acunetix.com/blog/web-security-zone/what-is-integer-overflow/>
- <https://stackoverflow.com/questions/9193880/overflowing-of-unsigned-int>

## Target 5 (Hint)

- Remember - Compiled with DEP enabled
- **W^X (Write XOR eXecute)!**
- You cannot put shellcode in buffer and point to it from return address.
- What can you do?
- **"Return-to-libc" style of attack**
- What does `system("/bin/sh");` do?
- Review lecture 6, detailed walkthrough

```
// Compiled with DEP enabled.

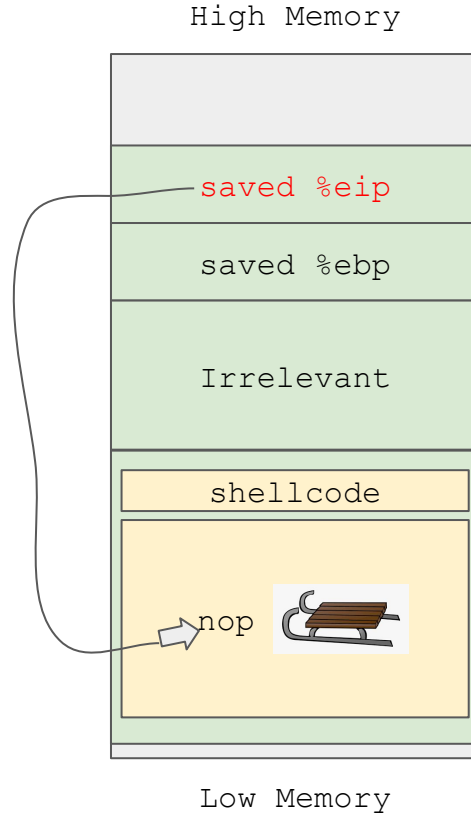
void greetings(void)
{
    system("echo Hello World");
}

void vulnerable(char *arg)
{
    char buf[10];
    strcpy(buf, arg);
}

int _main(int argc, char *argv[])
{
    ....
    setuid(0);
    vulnerable(argv[1]);
    greetings();
    ....
}
```

# Target 6 (Hint)

- ASLR (address-space layout randomization) enabled
- Stack can be placed anywhere in memory (0–255 bytes)
- How can I make sure my shell code always hit?
- Hint: NOP sleds
- NOP = 0X90
- Make sure you can open a root shell 100% of the time



```

// Compiled with -DMINIASLR.

void vulnerable(char *arg)
{
    char buf[1024];
    strcpy(buf, arg);
}

int _main(int argc, char *argv[])
{
    ....
    vulnerable(argv[1]);
    return 0;
}

```

# Target 7

- Return Oriented Programming
- Extra credit (15 points)
- Identical to target2, but it is compiled with DEP enabled
- Implement a ROP-based attack to bypass DEP and open a root shell.
- **ROPgadget.**
- The --binary, --badbytes, --multibr, and --ropchain flags will be particularly helpful.
- Closest to real-world attacks

```
void vulnerable(char *arg)
{
    char buf[100];
    strcpy(buf, arg);
}

int _main(int argc, char **argv)
{
    if (argc != 2) {
        fprintf(stderr, "Error: need a command-line
argument\n");
        return 1;
    }
    vulnerable(argv[1]);
    return 0;
}
```

```
setuid(0); //eax = 23
execve("/bin/sh", 0, 0); //eax = 11
```