

CSE 20

DISCRETE MATH

Fall 2020

<http://cseweb.ucsd.edu/classes/fa20/cse20-a/>

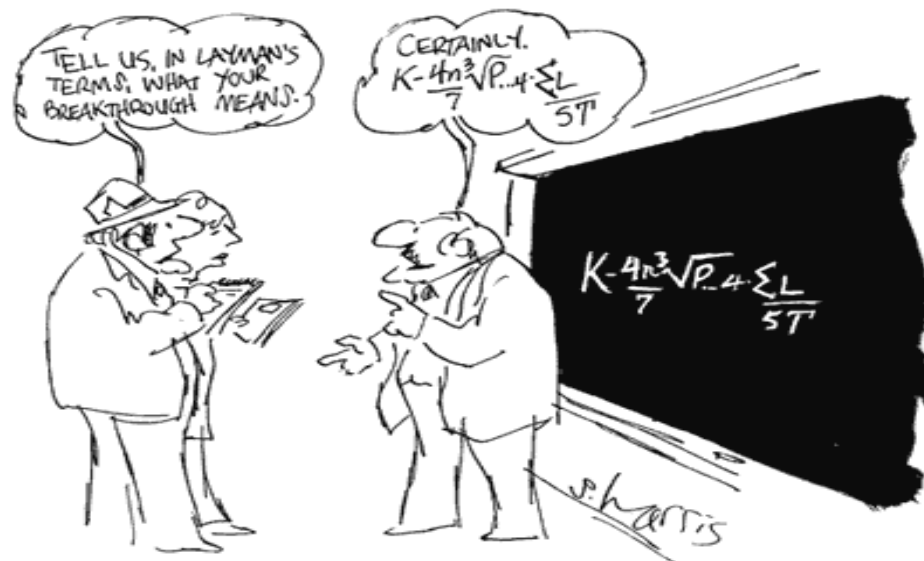
Today's learning goals

- Trace an algorithm specified in pseudocode
- Define the base expansion of a positive integer, specifically decimal, binary, hexadecimal, and octal.
- Convert between expansions in different bases of a positive integer.
- Define and use the **div** and **mod** operators.

Learning goals

In the past two classes, when have we used numbers?

Multiple Representations



Integer representations

Different contexts call for different representations.



Base 10



Base 2

Base b expansion of n

Rosen p. 246

Also known as **positional representation** of positive integers

Definition (Rosen p. 246) For b an integer greater than 1 and n a positive integer, the **base b expansion of n** is

$$(a_{k-1} \cdots a_1 a_0)_b$$

where k is a positive integer, a_0, a_1, \dots, a_{k-1} are nonnegative integers less than b , $a_{k-1} \neq 0$, and

$$n = a_{k-1}b^{k-1} + \cdots + a_1b + a_0$$

Using the terminology from last class: the base b expansion of n is a string over the alphabet $\{x \in \mathbb{N} \mid x < b\}$ and whose leftmost character is nonzero.

Base b expansion

In what base **could** this expansion be
 $(1401)_?$

- A. Binary (base 2)
- B. Octal (base 8)
- C. Decimal (base 10)
- D. Hexadecimal (base 16)
- E. More than one of the above

Base b expansion

In what base **could** this expansion be

$$(1401)_?$$

A. Binary (base 2)

B. Octal (base 8)

$$(1401)_8 = 1 \cdot 8^3 + 4 \cdot 8^2 + 1 = (769)_{10}$$

C. Decimal (base 10)

$$(1401)_{10} = 1 \cdot 10^3 + 4 \cdot 10^2 + 1 = 1401$$

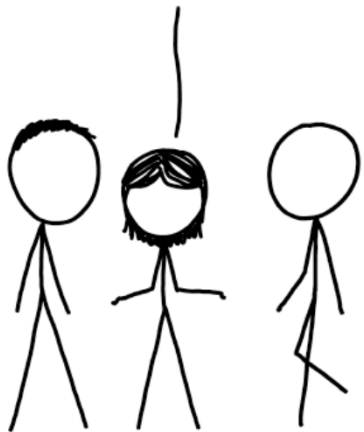
D. Hexadecimal (base 16)

$$(1401)_{16} = 1 \cdot 16^3 + 4 \cdot 16^2 + 1 = 5121$$

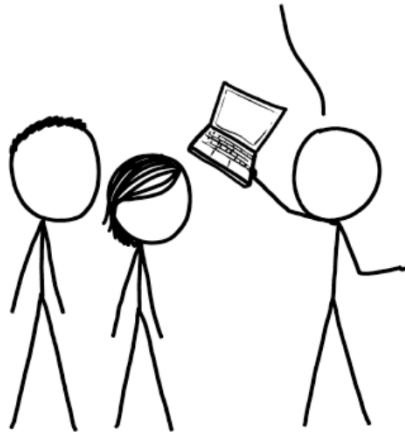
E. More than one of the above

Converting between bases

OUR FIELD HAS BEEN STRUGGLING WITH THIS PROBLEM FOR YEARS.



STRUGGLE NO MORE!
I'M HERE TO SOLVE
IT WITH *ALGORITHMS!*



SIX MONTHS LATER:

WOW, THIS PROBLEM
IS REALLY HARD.

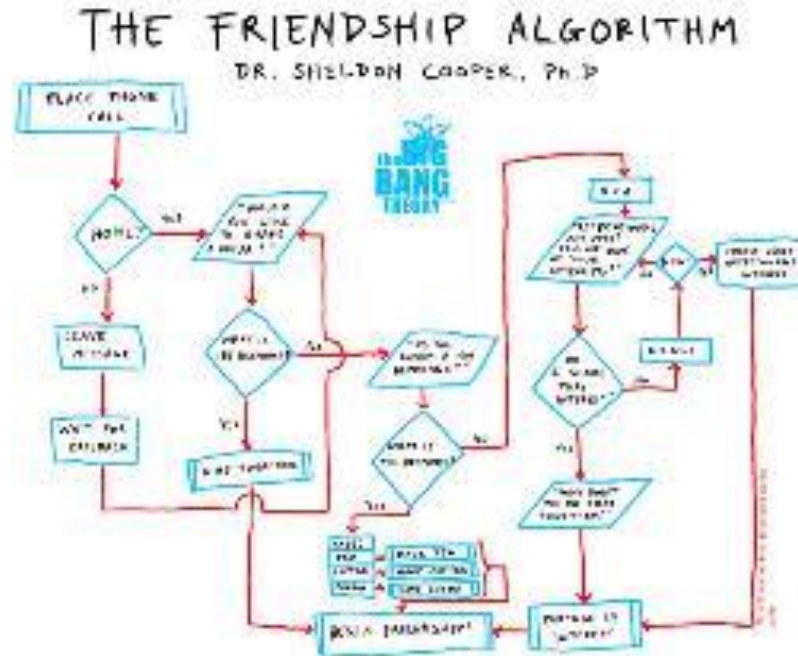
YOU DON'T SAY.



Algorithm?

Rosen 3.1 p. 191

Finite sequence of precise instructions for solving problem.



Algorithm: Pseudocode

Appendix

Finite sequence of precise instructions for solving problem.

```
1 procedure log(n: a positive integer)
2   r := 0
3   while n > 1
4     r := r + 1
5     n := n div 2
6   return r {r holds the result of the log operation}
```

At the end of running *log*(6) what values are in the variables *r* and *n*?

- A. $r = 6, n = 0$
- B. $r = 6, n = 6$
- C. $r = 2, n = 0$
- D. $r = 2, n = 1$
- E. None of the above.

Algorithm: constructing base b expansion

Input n, b **Output** k , coefficients in expansion

- English description.

- Pseudocode.

Algorithm 1: constructing base b expansion

Input n, b **Output** k , coefficients in expansion

- English description.

Initialize value remaining to be n

Find biggest power of b that is less than or equal to value remaining.

Increment appropriate coefficient.

Update value remaining by subtract this power of b from it.

Repeat until value remaining is 0.



Ternary representation of 17

- A. $(17)_3$
- B. $(211)_3$
- C. $(122)_3$
- D. $(221)_3$
- E. $(112)_3$

Algorithm 1: constructing base b expansion

Calculating base b expansion, from left

```
1 procedure baseb1( $n, b$ : positive integers with  $b > 1$ )
2  $v := n$ 
3  $k := \text{log}_b(n, b) + 1$ 
4 for  $i := 1$  to  $k$ 
5    $a_{k-i} := 0$ 
6   while  $v \geq b^{k-i}$ 
7      $a_{k-i} := a_{k-i} + 1$ 
8      $v := v - b^{k-i}$ 
9 return  $(a_{k-1}, \dots, a_0)$   $\{(a_{k-1} \dots a_0)_b$  is the base  $b$  expansion of  $n\}$ 
```

a_{k-1} is coefficient of biggest power of b that is less than n
Thus: k is 1 more than integer part of $\log_b n$

Algorithm 2: constructing base b expansion

Input n, b **Output** k , coefficients in expansion

Idea: Find smallest digit first, then next smallest, etc.

.... **but how?**

Bases and Divisibility

Rosen p. 237-239

Theorem: For n an integer and d a positive integer, there are unique integers q and r with $0 \leq r < d$ and $n = dq + r$. **Notation:** $q = n \text{ div } d$ and $r = n \text{ mod } d$

When $k > 1$

$$n = a_{k-1}b^{k-1} + \dots + a_1b + a_0$$

$$n = b(a_{k-1}b^{k-2} + \dots + a_1) + a_0$$

d

$q = n \text{ div } d$

$r = n \text{ mod } d$

Algorithm 2: constructing base b expansion

Input n, b **Output** k , coefficients in expansion

Idea: Use $n \bmod b$ to compute least significant digit.

Use $n \operatorname{div} b$ to compute new integer whose expansion we need. Repeat.

Algorithm 2: constructing base b expansion

Calculating base b expansion, from right

```
1 procedure base2( $n, b$ : positive integers with  $b > 1$ )
2    $q := n$ 
3    $k := 0$ 
4   while  $q \neq 0$ 
5      $a_k := q \bmod b$ 
6      $q := q \operatorname{div} b$ 
7      $k := k + 1$ 
8   return  $(a_{k-1}, \dots, a_0)\{(a_{k-1}, \dots, a_0)_b$  is the base  $b$  expansion of  $n\}$ 
```

n	b	q	k	a_k	$q \neq 0?$

Representing more

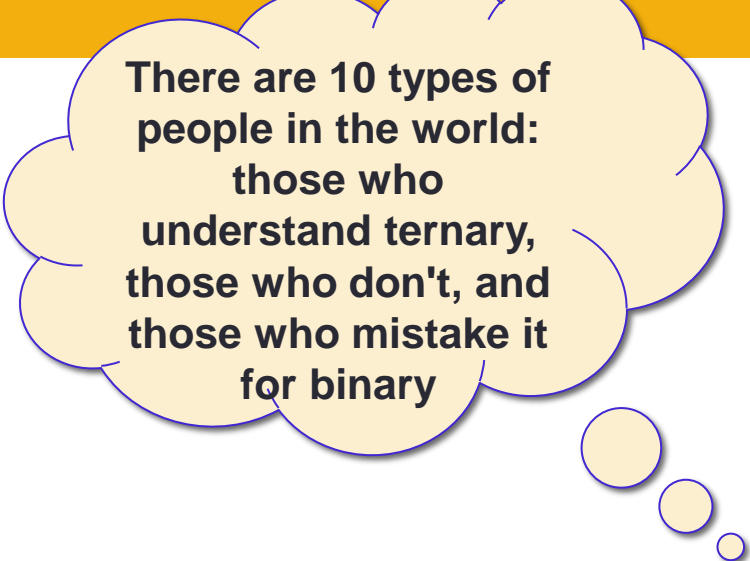
- Base b expansions can express any **positive integers**
- What about
 - Zero?
 - negative integers?
 - rational numbers?
 - other real numbers?

For next time

- Read website carefully

<http://cseweb.ucsd.edu/classes/fa20/cse20-a/>

- No pre-class reading for next lecture



**There are 10 types of
people in the world:
those who
understand ternary,
those who don't, and
those who mistake it
for binary**