

Recall that S is defined as the set of all RNA strands, strings made of the bases in $B = \{\mathbf{A}, \mathbf{U}, \mathbf{G}, \mathbf{C}\}$. Define the functions *mutation*, *insertion*, and *deletion* as described by the pseudocode below:

```

1 procedure mutation( $b_1 \cdots b_n$ : a RNA strand,  $k$ : a positive integer,  $b$ : an element of  $B$ )
2 for  $i := 1$  to  $n$ 
3   if  $i = k$ 
4      $c_i := b$ 
5   else
6      $c_i := b_i$ 
7 return  $c_1 \cdots c_n$  {The return value is a RNA strand made of the  $c_i$  values}

```

```

1 procedure insertion( $b_1 \cdots b_n$ : a RNA strand,  $k$ : a positive integer,  $b$ : an element of  $B$ )
2 if  $k > n$ 
3   for  $i := 1$  to  $n$ 
4      $c_i := b_i$ 
5    $c_{n+1} := b$ 
6 else
7   for  $i := 1$  to  $k-1$ 
8      $c_i := b_i$ 
9    $c_k := b$ 
10  for  $i := k+1$  to  $n+1$ 
11     $c_i := b_{i-1}$ 
12 return  $c_1 \cdots c_{n+1}$  {The return value is a RNA strand made of the  $c_i$  values}

```

```

1 procedure deletion( $b_1 \cdots b_n$ : a RNA strand,  $k$ : a positive integer)
2 if  $k > n$ 
3    $m := n$ 
4   for  $i := 1$  to  $n$ 
5      $c_i := b_i$ 
6 else
7    $m := n-1$ 
8   for  $i := 1$  to  $k-1$ 
9      $c_i := b_i$ 
10  for  $i := k$  to  $n-1$ 
11     $c_i := b_{i+1}$ 
12 return  $c_1 \cdots c_m$  {The return value is a RNA strand made of the  $c_i$  values}

```

Mut with domain $S \times S$ is defined by, for $s_1 \in S$ and $s_2 \in S$,

$$Mut(s_1, s_2) = \exists k \in \mathbb{Z}^+ \exists b \in B (mutation(s_1, k, b) = s_2)$$

Ins with domain $S \times S$ is defined by, for $s_1 \in S$ and $s_2 \in S$,

$$Ins(s_1, s_2) = \exists k \in \mathbb{Z}^+ \exists b \in B (insertion(s_1, k, b) = s_2)$$

Del with domain $S \times S$ is defined by, for $s_1 \in S$ and $s_2 \in S$,

$$Del(s_1, s_2) = \exists k \in \mathbb{Z}^+ (deletion(s_1, k) = s_2)$$

Definition: We say that a RNA strand s_1 is “within one edit” of a RNA strand s_2 to mean

$$Mut(s_1, s_2) \vee Mut(s_2, s_1) \vee Ins(s_1, s_2) \vee Ins(s_2, s_1) \vee Del(s_1, s_2) \vee Del(s_2, s_1)$$

$within1_{TF} : \underline{\hspace{2cm}} \rightarrow \underline{\hspace{2cm}}$

$within1_{\mathcal{P}} : \underline{\hspace{1cm}} \rightarrow \underline{\hspace{2cm}}$

$within1_{TF}(s_1, s_2) = \underline{\hspace{4cm}}$

$within1_{\mathcal{P}}(s_1) = \underline{\hspace{4cm}}$

$$W_1 = \{ \underline{\hspace{4cm}} \}$$

Definition: When A and B are sets, we say any subset of $A \times B$ is a **binary relation**. There are other ways to represent a relation R

- A function $f_{TF} : \text{_____} \rightarrow \text{_____}$ with $f_{TF}(\text{_____}) = \text{_____}$
- A function $f_{\mathcal{P}} : \text{_____} \rightarrow \text{_____}$ with $f_{\mathcal{P}}(\text{_____}) = \text{_____}$

Definition: When A is a set, we say any subset of $A \times A$ is a (binary) **relation** on A .

Definition: Let $R_{(\text{mod } n)}$ be the set of all pairs of integers (a, b) such that $(a \bmod n = b \bmod n)$. Then a is **congruent to $b \bmod n$** means $(a, b) \in R_{(\text{mod } n)}$. A common notation is to write this as $a \equiv b(\bmod n)$.

Some example elements of $R_{(\text{mod } 4)}$ are: _____

Definition: (*Rosen 9.1*) A relation R on a set A is called **reflexive** means $(a, a) \in R$ for every element $a \in A$. A relation R on a set A is called **symmetric** means $(b, a) \in R$ whenever $(a, b) \in R$, for all $a, b \in A$. A relation R on a set A is called **transitive** means whenever $(a, b) \in R$ and $(b, c) \in R$, then $(a, c) \in R$, for all $a, b, c \in A$.

Relation	Reflexive? (why/why not)	Symmetric? (why/why not)	Transitive? (why/why not)
W_1			
$R_{(\text{mod } 4)}$			

Definition: (*Rosen 9.5*) A relation is an **equivalence relation** means it is reflexive, symmetric, and transitive.

Definition: (*Rosen 9.5*) An **equivalence class** of an element $a \in A$ for an equivalence relation R on the set A is the set $\{s \in A | (a, s) \in R\}$. We write this as $[a]_R$.

Some examples of elements of $[5]_{R_{(\text{mod } 4)}}$ are: _____

Some examples of elements of $[9]_{R_{(\text{mod } 4)}}$ are: _____

Some examples of elements of $[6]_{R_{(\text{mod } 4)}}$ are: _____

Definition: A **partition** of a set A is a set of non-empty, disjoint subsets A_1, A_2, \dots, A_n such that $A_1 \cup A_2 \cup \dots \cup A_n = A$.

We can partition the set of integers using equivalence classes of $R_{(\text{mod } 4)}$ using: _____