

# Week 10 Discussion

# FAQ

- Q: When and where is the final exam?
- A: 8 am, 11/12 on Wednesday. Center Hall 113.
  
- Q: Are we allowed to bring cheat sheets?
- A: Yes, one 8.5x11 doubled sided cheat sheet is allowed.
  
- Q: What will be covered in the final exam?
- A: Everything in lecture slides and assigned readings.

# FAQ

- Q: Will there be a final review session?
- A: Yes, during the last lecture on Friday.
  
- Q: Will there be any TA office hours in the final week?
- A: Yes, we will update the schedule on piazza soon.
  
- Q: Do we have practice finals?
- A: Yes, we will post practice finals later this week.

# Aloha Protocol

Goal: distributed access control over a shared broadcast channel

Fact: only one transmission at a time is allowed in shared broadcast channel

Algorithm:

If a node has data to send:

    try to send the data

    while failed:

        delay for some random time interval

        resend the data

# Aloha Protocol

Suppose there are two nodes A and B in a channel. All frames transmitted by A and B have identical length  $L$ , and take the same amount of time  $T$  to transmit from one node to the other.

a. If node A starts to transmit a frame  $F1$  at time  $t1 = 0$ , and node B starts to transmit a frame  $F2$  at time  $t2 = 0.5T$ . Which transmission(s) will be successful?

b. If node A wants to successfully transmit a frame  $F1$  starting from  $t1 = T$ , give a time range that node B cannot start to transmit any frames.

# Slotted Aloha

- (1) Time is divided into equal size slots
- (2) Time slots should be synchronized across all nodes in the channel
- (3) Each node can transmit only at the beginning of a time slot.
- (4) Retransmissions occur at the nearest next slot

# Slotted Aloha

- Suppose there are  $n$  nodes in a channel, and each node has a transmission probability  $p$ .
- The probability that node  $n_1$  successfully transmits data in its first try is:  
$$S(n_1 \text{ success}) = p (1-p)^{(n-1)}$$
- The probability that there exists one node (among  $n$  nodes) that successfully transmits data in the first try is:  
$$S(\text{any one node in } n \text{ success}) = np(1-p)^{(n-1)}$$

# Slotted Aloha Example

Suppose two stations A and B use the slotted Aloha protocol to send frames in the same channel. For each time slot of  $T$  ms, station A and B have transmission probability  $P_a$  and  $P_b$ , respectively. Assume station A and B have infinite number of frames to transmit, and at  $t = 0$ , station A sends a frame  $F_a$  and simultaneously station B sends a frame  $F_b$ ; both  $F_a$  and  $F_b$  have size less than maximum transmission unit (MTU).

- a. What is the probability that station A successfully transmits  $F_a$  at  $t = 0$ ?
- b. What is the probability that station A successfully transmits  $F_a$  at  $t = T$ ?
- c. What is the probability that station A successfully transmits  $F_a$  at  $t = 2T$ ?



# Slotted Aloha Example

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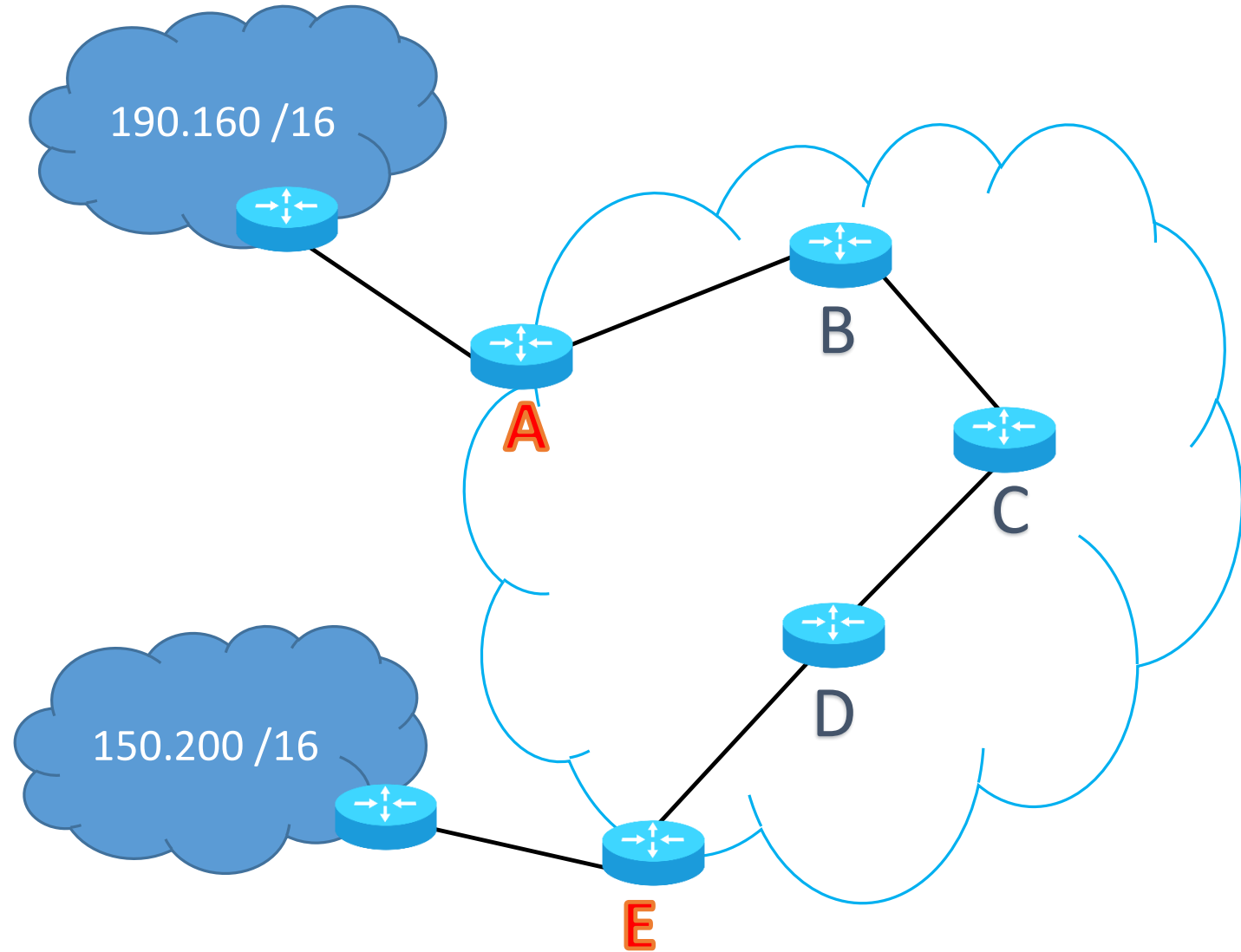
Now suppose  $P_a = 2P_b$ , and at  $t = T$ , the probability of either  $F_a$  or  $F_b$  is successfully transmitted is 0.56. What is the value of  $P_a$ ?

# Border Gateway Protocol (BGP)

## True or False

- a. One vulnerability of BGP is that packets might be lost during transmission.
- b. BPG routing decisions are made based on shortest AS path (number of hops).
- c. Routers running BPG should frequently exchange keepalive messages with each other.
- d. Internal BGP (iBGP) is established between routers that are adjacent, but within the same Autonomous System.

# Interior Border Gateway Protocol (iBGP)



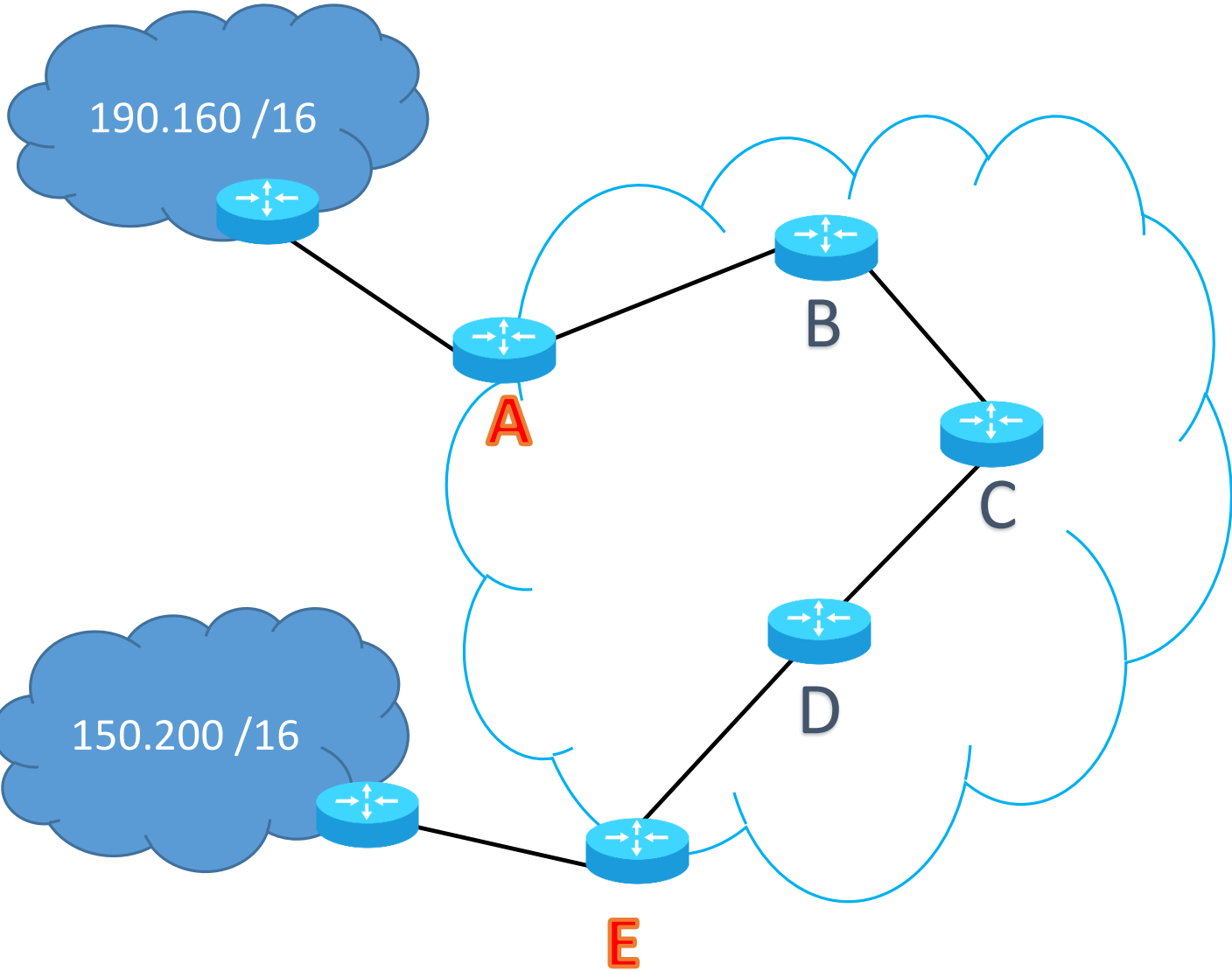
BGP table for the AS

Prefix	BGP Next Hop
190.160 /16	A
150.200 /16	E

IGP table for router C

Router	IGP Path
A	B
B	B
D	D
E	D

# Interior Border Gateway Protocol (iBGP)



Combined table for router C

Prefix	IGP Path
190.160 /16	B
150.200 /16	D

Questions?