

mlin: Rethinking and Rebooting `gprof` for the Multicore Era

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Activating a “gprof for parallelization”

How effective are programmers at picking the right parts of a program to parallelize?

User study* we performed at UC San Diego (UCSD IRB #100056)

First and second year CS graduate students

Users parallelize their programs and submit to job queue for timing

32-core AMD machine, Cilk++, access to gprof

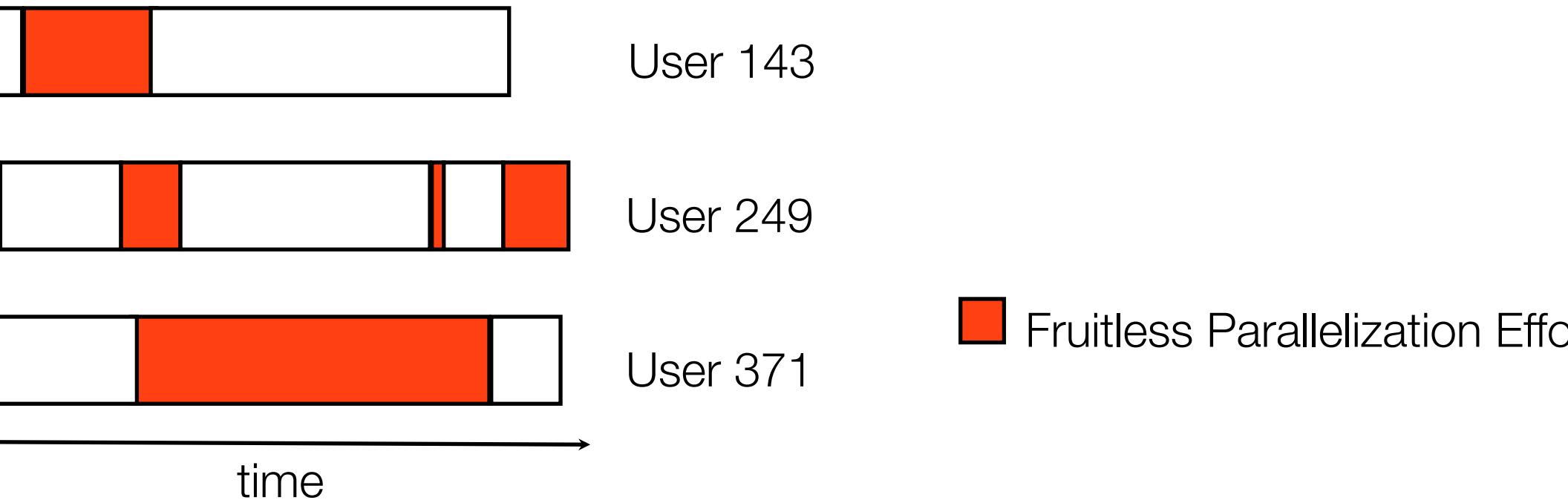
Students were graded based on effectiveness of their parallel speedup

Students told serial optimization would not help their grade

**Disclaimer: No graduate students w*

er Study: Results

Examined student's activities to determine result of efforts



significant fraction of fruitless effort because of three basic problems

Low Parallelism: Region was not parallel enough

Low Coverage: Region's execution time was too small

gprof answers the question:

*“What parts of this program should I spend time **optimizing**?”*

Kremlin answers the question:

*“What parts of this program should I spend time **parallelizing**?”*

kremlin's Usage Model

Usage model inspired by gprof

```
make CC=kremlin-cc
```

```
/tracking lolcats
```

```
kremlin tracking --personality=openmp
```

1. Compile instrumented binary

2. Profile with sample input

3. Run analysis tool to create

file (lines)	Self-P	Cov (%)
imageBlur.c (49-58)	145.3	9.7
imageBlur.c (37-45)	145.3	8.7
getInterpPatch.c (26-35)	25.3	8.9
calcSobel_dX.c (59-68)	126.2	8.1
calcSobel_dX.c (46-55)	126.2	8.1

mlin's Key Components

al
c
de

Parallelism Discovery
“What’s the potential parallel speedup of each part of this program?”

- *Hierarchical Critical Path Analysis (HCPA)*
 - ▶ Quantifies self-parallelism in each program region
- *Self-Parallelism*
 - ▶ Estimates ideal parallel speedup of a specific region

Parallelism Planning
“What regions must I parallelize to get the maximum benefit on this system?”

- *Planning Personalities*
 - ▶ Incorporates target systems constraints in parallelization

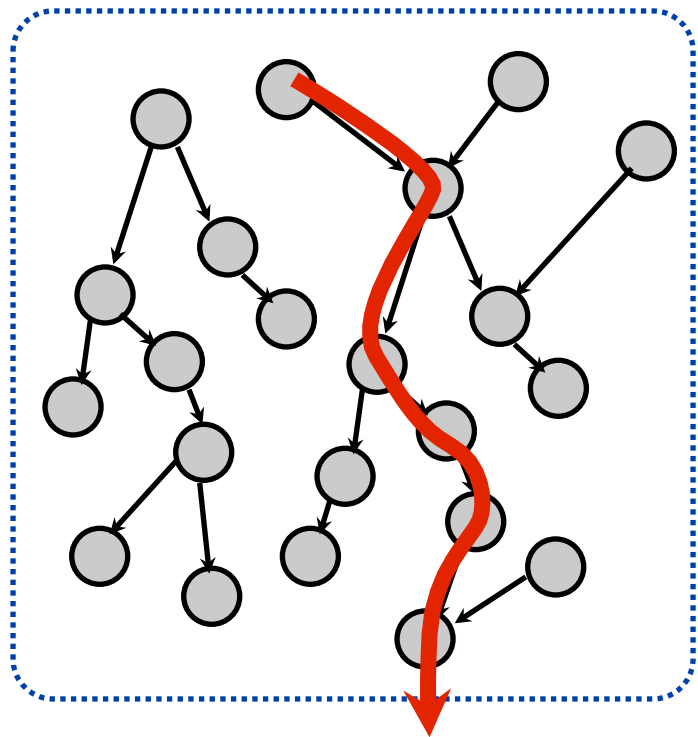
Parallelization

Developing an Approach for Parallelism Discovery

Existing Technique: 1980's-era Critical Path Analysis (CPA)

Finds critical path through the dynamic execution of a program

Mainly used in research studies to quantify limits of parallelism



critical path (cp)

$$\text{parallelism} = \frac{\text{work}}{\text{critical path length}}$$

work \sim # of instrs

○ instruction

↘ data or control dependence

Benefits of CPA as a Basis for a Parallelism Discovery

Estimates program's potential for parallelization under relatively optimistic assumptions

Closer approximation to what human experts can achieve

versus pessimistic static analysis in automatic parallelizing compilers

It is predictive of parallelism after typical parallelization transformations

e.g., Loop interchange, loop fission, locality enhancement

Improving CPA with Hierarchical CPA (HCPA)

A is typically run on an entire program

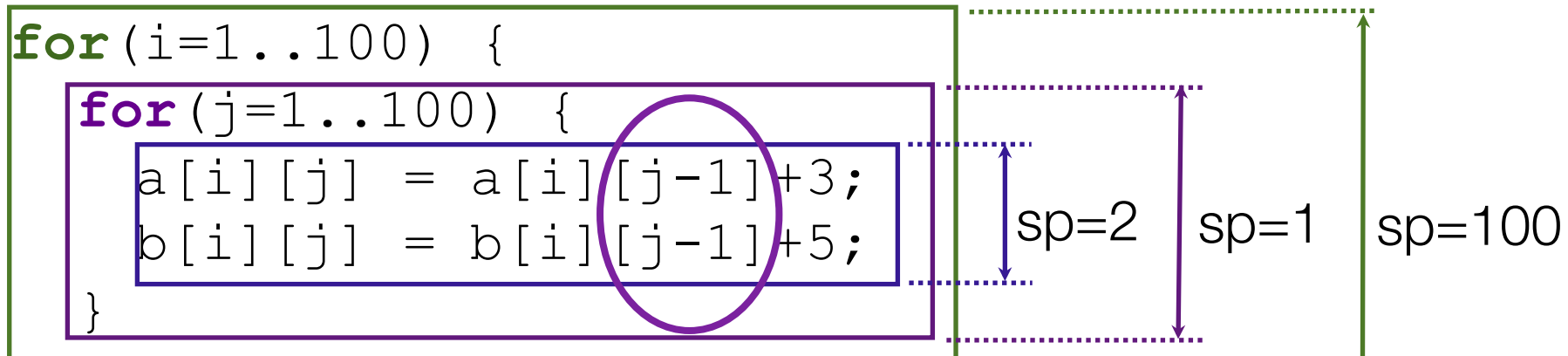
Not helpful for identifying specific regions to parallelize

Doesn't help evaluate execution time of a program if only a subset of the program is parallelized

Hierarchical CPA is a region-based analysis

Self-Parallelism (sp) identifies parallelism in specific regions

Provides basis for estimating parallel speedup of individual regions



CPA Step 1: Hierarchically Apply CPA

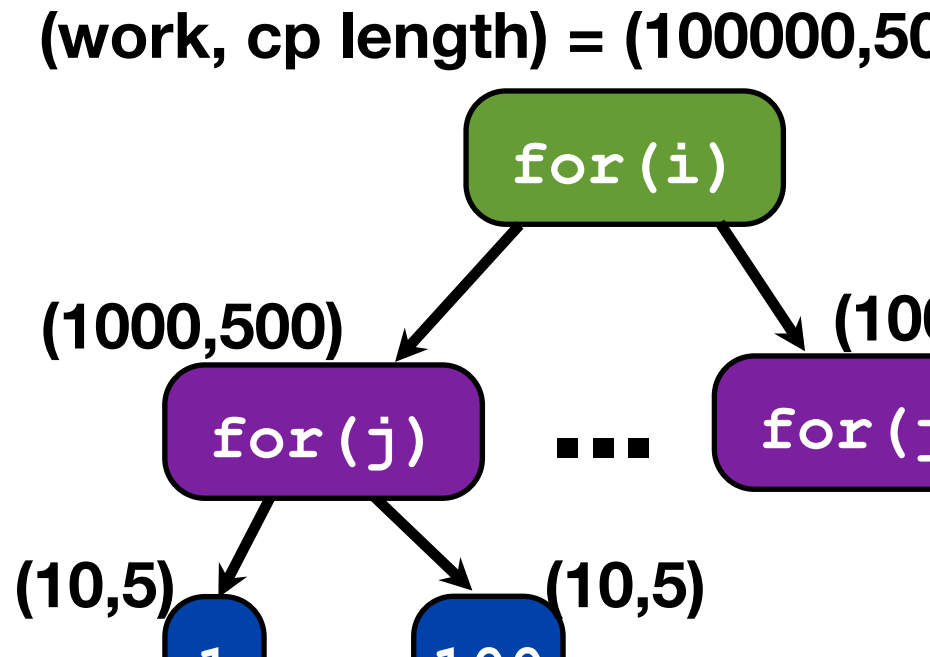
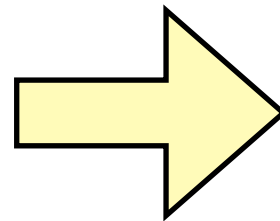
Goal: Introduce localization through *region*-based analysis

Window-memory based implementation

Performs CPA analysis on every program region

Single pass: Concurrently analyzes multiple nested regions

```
for (i=1..100) {  
  for (j=1..100) {  
    a[i][j] = a[i][j-1]+3;  
    b[i][j] = b[i][j-1]+5;  
  }  
}
```



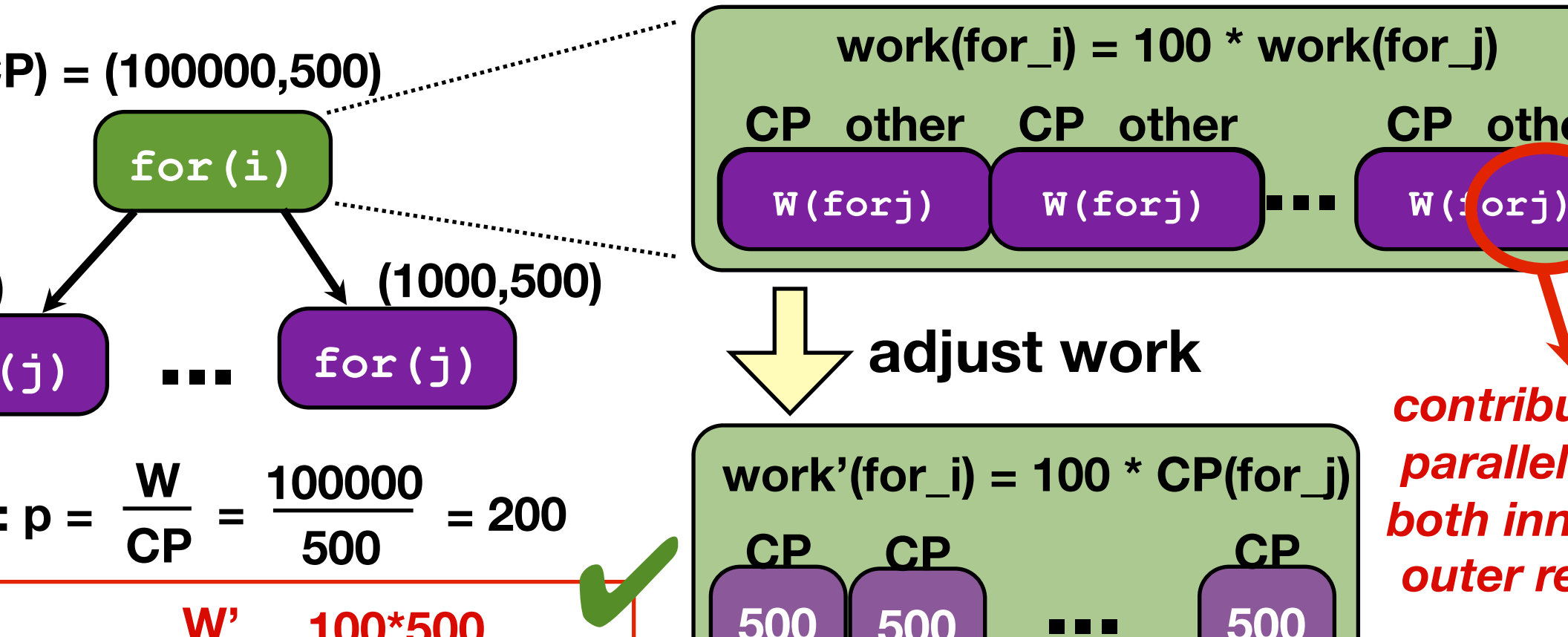
$$p = \frac{W}{CP} = \frac{100000}{500} = 200 \quad \text{X}$$

HCPA Step 2: Calculate Self-Parallelism

Goal: Eliminate effect of nested parallelism in parallelism calculation

Approximate self-parallelism using only HCPA output

Subtracts "nested parallelism" from overall parallelism



PA Step 3: Compute Static Region Data

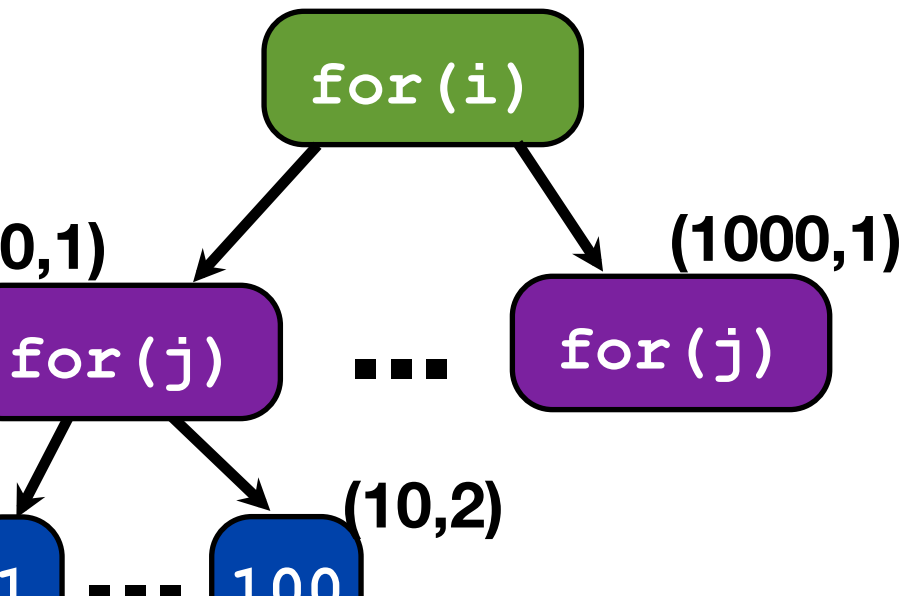
Goal: Convert dynamic region data to static region output

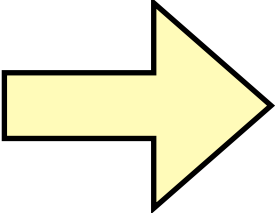
Task: Merge dynamic nodes associated with same static region

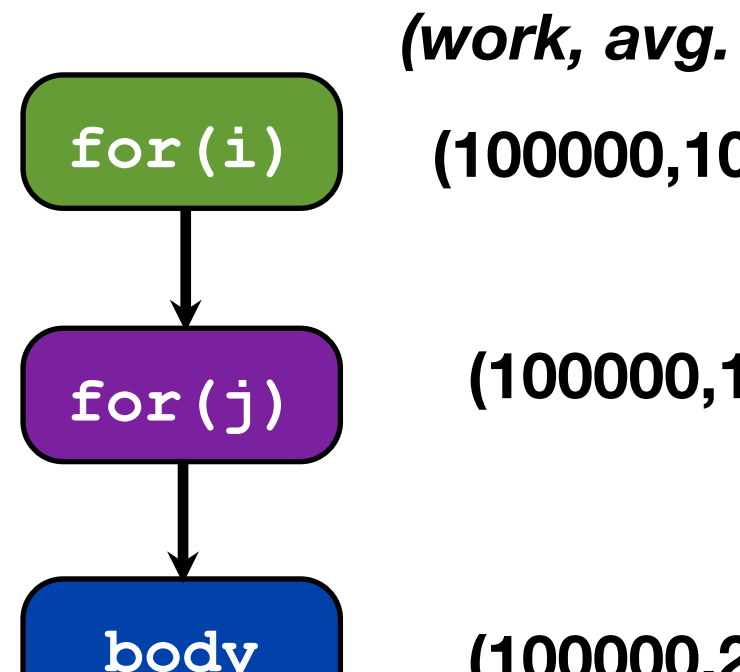
Work: *Sum* of work across dynamic instances

Self-Parallelism: *Weighted Average* across dynamic instances

$(work, sp) = (100000, 100)$




merge
dynamic



Other Details on Discovery in Our Paper

mlin handles much more complex structures than just nested loops: finds parallelism in arbitrary code including recursion

f-parallelism metric is defined and discussed in detail in the paper

mpression technique used to reduce size of HCPA output

Generating a Parallelization Plan

Goal: Use HCPA output to select best regions for target system

Planning personalities allow user to incorporate system constraints

Software constraints: What types of parallelism can I specify?

Hardware constraints: Synchronization overhead, etc.

Planning algorithm can change based on constraints

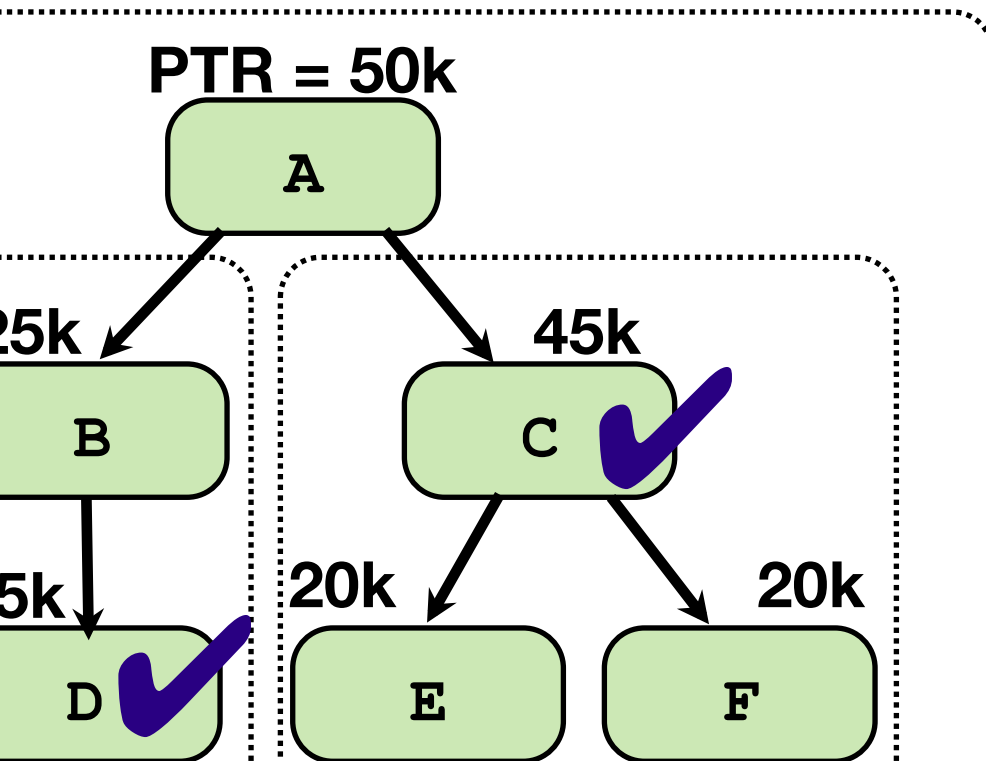
OpenMP Planner

Based on OpenMP 2.0 specification

Focused on loop-level parallelism

Disallows nested parallelism because of overhead

Planning algorithm based on dynamic programming



parallelized time reduction = $W - (W/P)$

Region	Work	SP
A	100k	2
B	50k	2
C	50k	10
D	50k	10
F	25k	5

Evaluation

Methodology:

Ran Kremlin on serial versions; targeting OpenMP

Parallelized according to Kremlin's plan

Gathered performance results on 8 socket AMD 8380 Quad-core

Compared against third-party parallelized versions (3rd Party)

Benchmarks: NAS OpenMP and SpecOMP

Have both serial and parallel versions

Wide range of parallel speedup (min: 1.85x, max: 25.89x) on 32 cores

How much effort is saved using Kremlin?

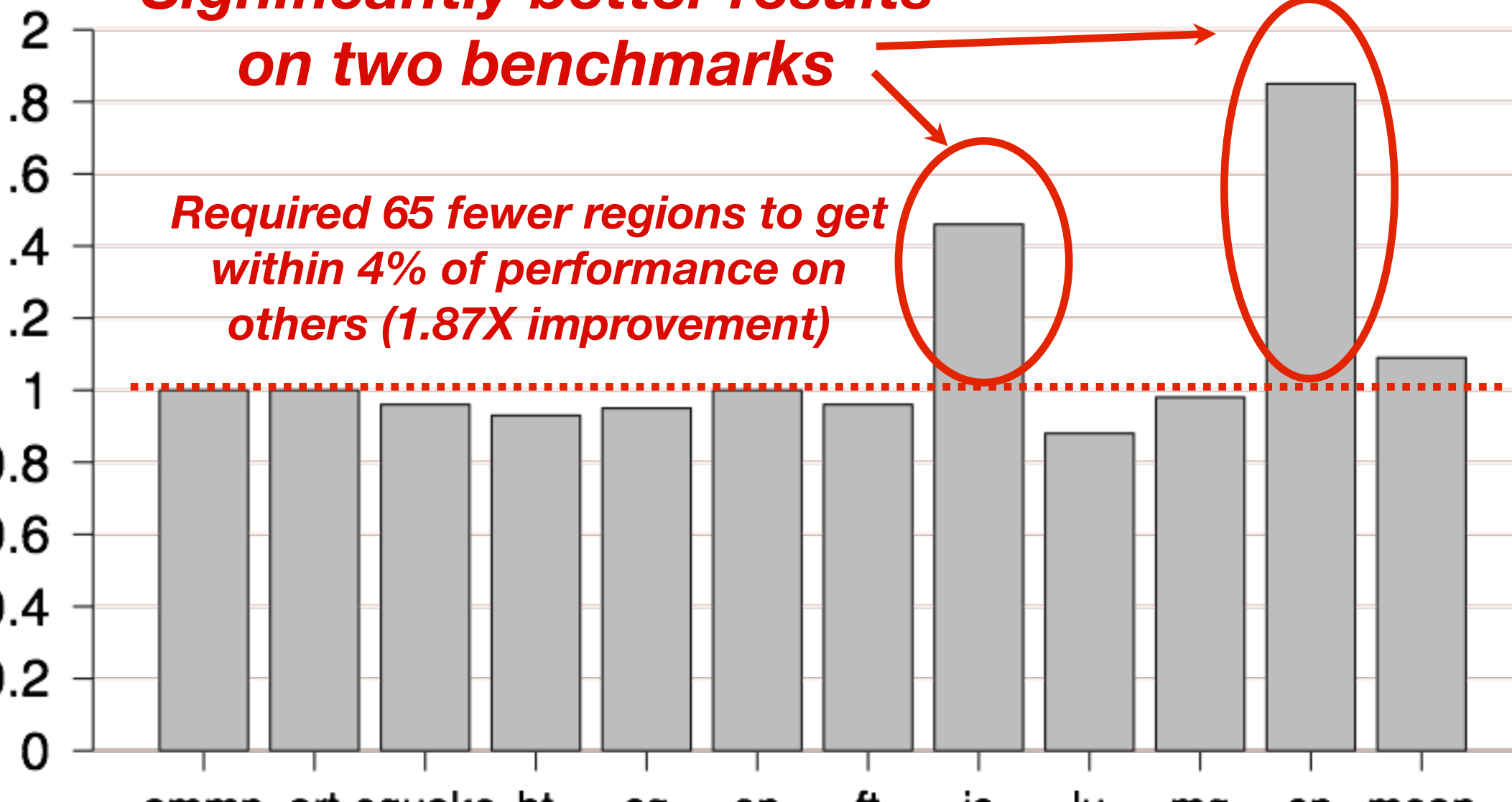
		# of Regions Parallelized		
Suite	Benchmark	3rd Party	Kremlin	Reduction
ecOMP	art	3	4	0.75x
	ampp	6	3	2.00x
	equake	10	6	1.67x
NPB	ep	1	1	1.00x
	is	1	1	1.00x
	ft	6	6	1.00x
	mg	10	8	1.25x
	cg	22	9	2.44x
	lu	28	11	2.55x
	bt	54	27	2.00x
	sp	70	58	1.21x
	Overall	211	134	1.57x

How good is Kremlin-guided performance?

Compared performance against expert, third-party version

**Significantly better results
on two benchmarks**


**Required 65 fewer regions to get
within 4% of performance on
others (1.87X improvement)**



Does Kremlin pick the best regions first?

Determined what % of speedup comes from first {25,50,75,100} recommended regions

	Fraction of Kremlin Plan Applied			
	First 25% of regions	Second 25% of regions	Third 25% of regions	Last 25% regions
Marginal benefit (% max speedup) (avg)	56.2%	30.2%	9.2%	4.4%



Red circles highlight the values 56.2% and 4.4%. Red arrows point from these circles to the green arrow below the table.

Conclusion

Kremlin helps a programmer determine:

What parts of this program should I spend time parallelizing?

Three key techniques introduced by Kremlin

Hierarchical CPA: How much total parallelism is in each region?

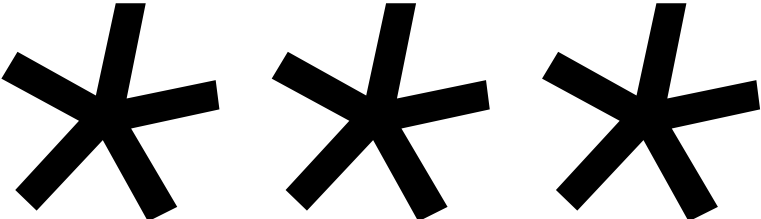
Self-Parallelism: How much parallelism is only in this region?

Planning Personalities: What regions are best for my target system?

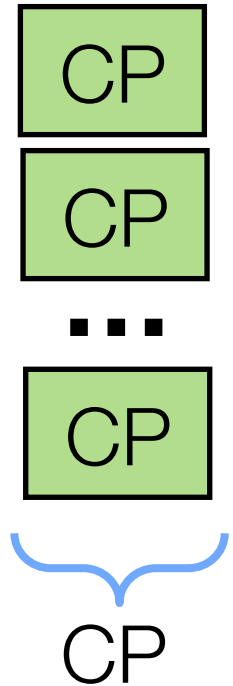
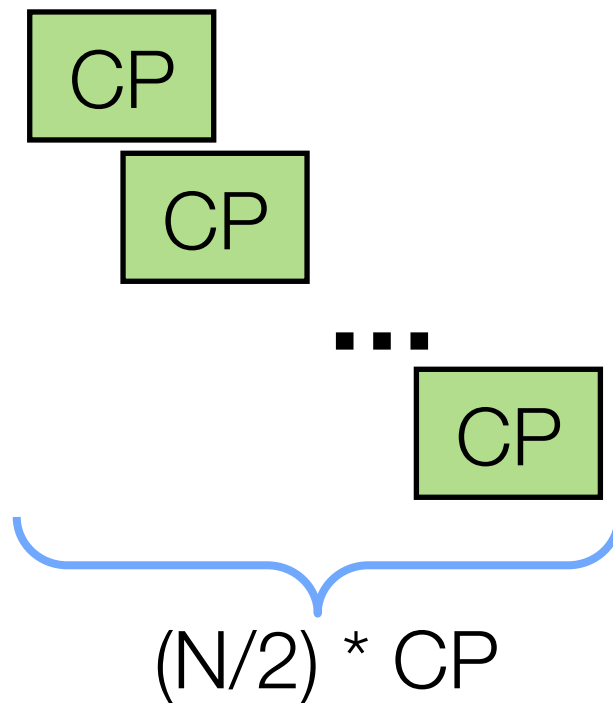
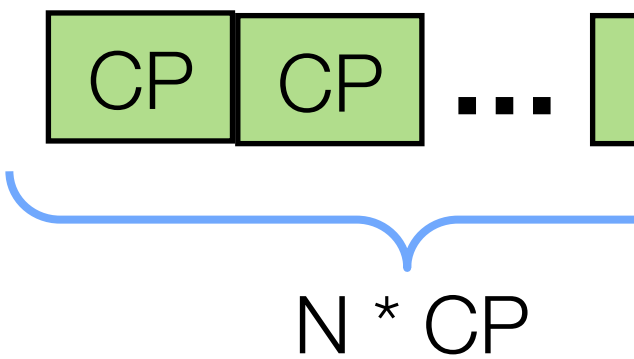
Compelling results

1.57x average reduction in number of regions parallelized

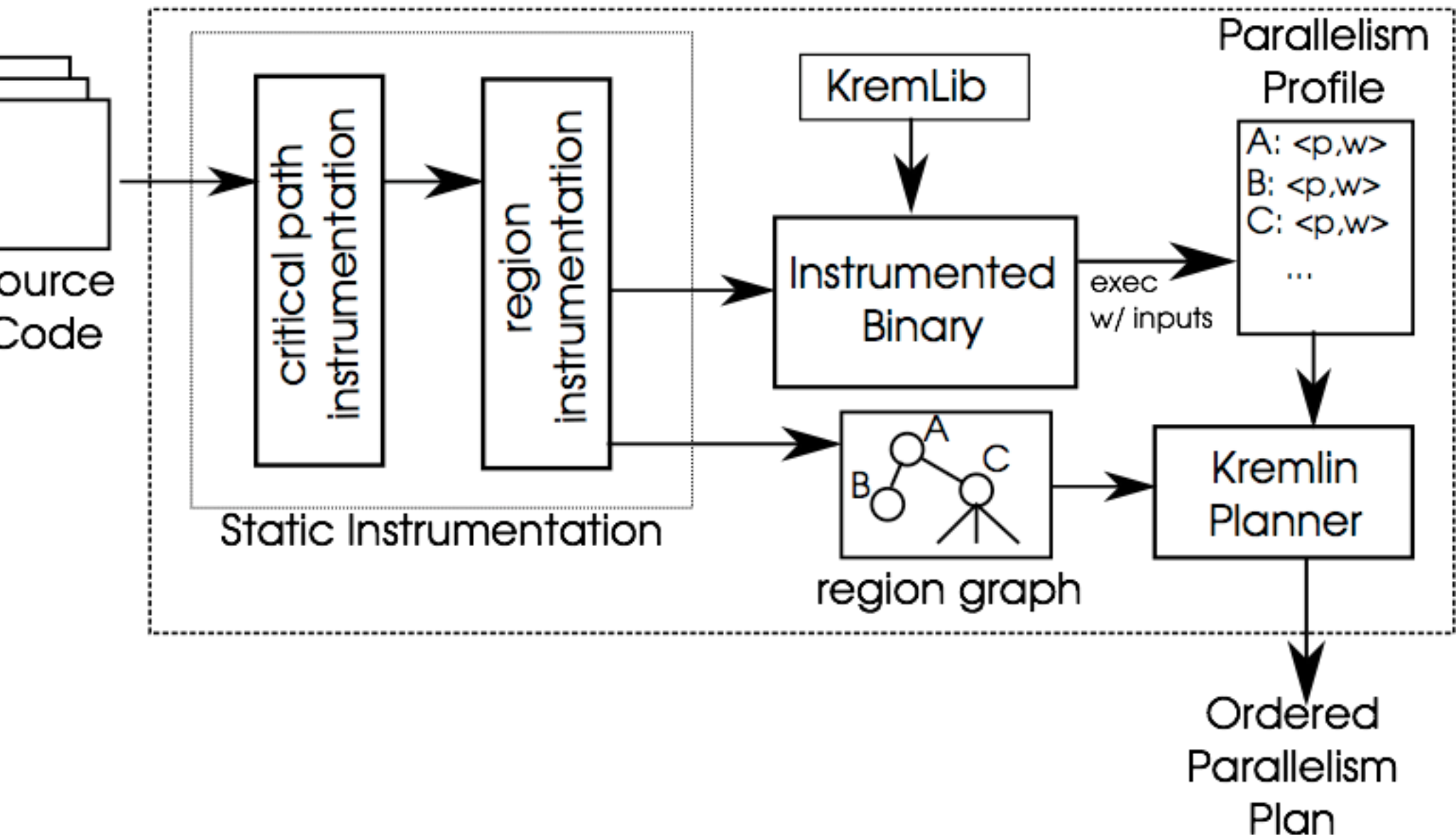
Greatly improved performance on 2 of 11 benchmarks; very close to others



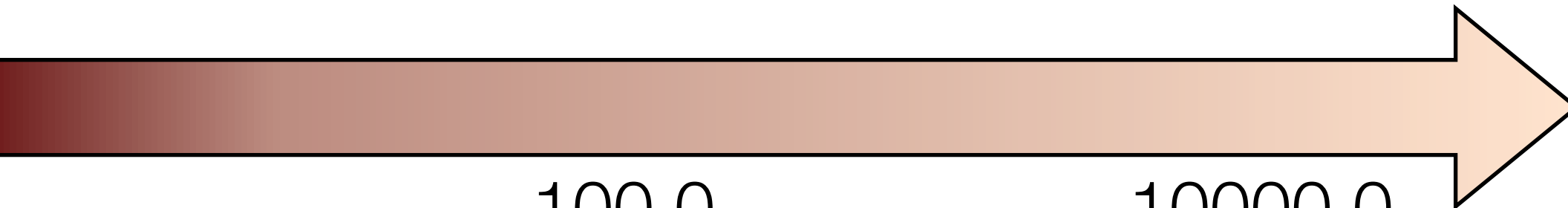
Parallelism for Three Common Loop Types

Type	DOALL	DOACROSS	Serial
Path (CP)			
CP)	$N * CP$	$N * CP$	$N * CP$
m	$\frac{N * CP}{CP} = \mathbf{N}$	$\frac{N * CP}{(N/2) * CP} = \mathbf{2.0}$	$\frac{N * CP}{N * CP} = \mathbf{1.}$

Kremlin System Architecture



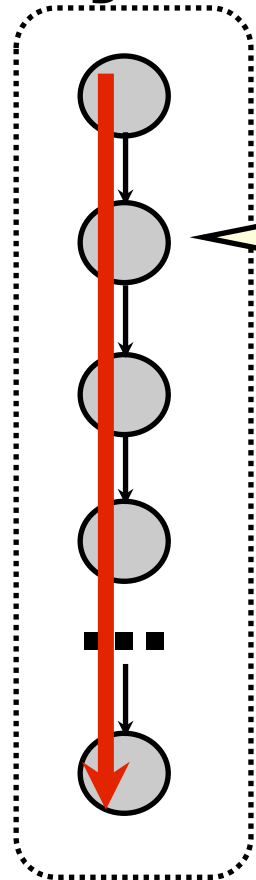
Interpreting the Parallelism Metric



100.0

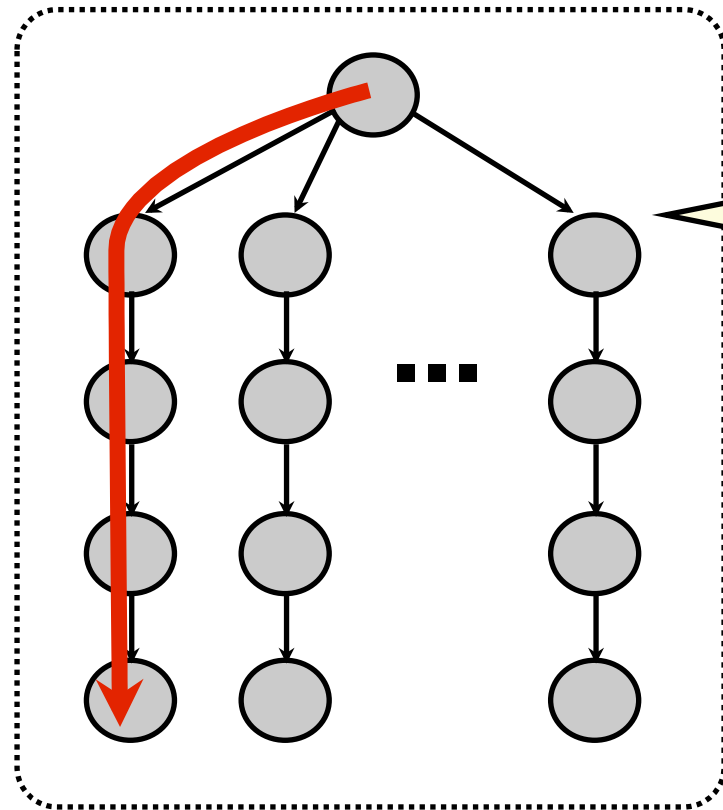
10000.0

Completely Serial



All work is
on critical
path
($ET == CP$)

Highly Parallel



Most work
off critical
path
($ET >$

Parallelism is a result of execution time