
LMgr: A Low-Memory Global Router with Dynamic Topology Update and Bending- Aware Optimum Path Search

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Outline

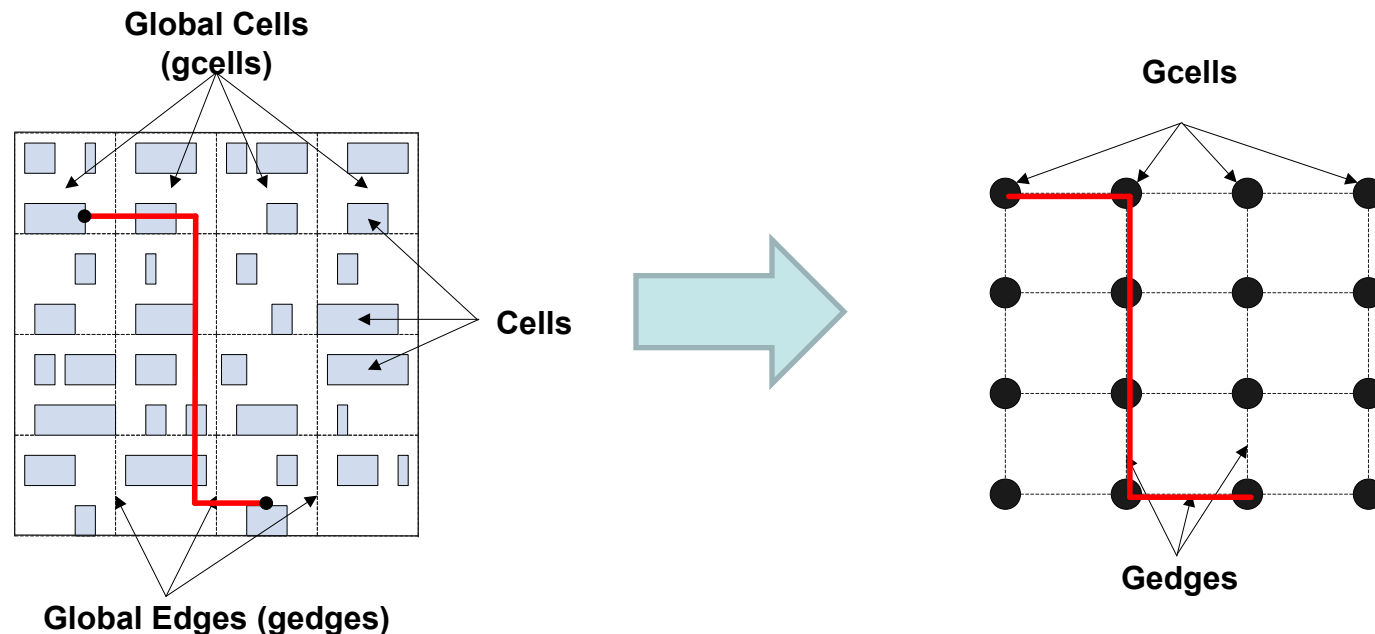
- **Introduction**
- **Problem formulation**
- **Existing problems and our solutions**
- **A new global router → LMgr**
- **Experimental Results**
- **Conclusion**

Introduction

- **Fundamental research problem**
 - Performance dominance
 - Impacts on timing and power closure
 - Bottleneck on design efficiency
- **ISPD 2007-2008 global routing contest** [ISPD08]
 - Bookshelf-format benchmark suite from real VLSI design
 - Lots of related research works have been proposed
- **State-of-the-art approaches**
 - Integer linear programming [Cho07, Shojaei11]
 - Iterative rip-up and rerouting [Chang08, chen09]

Problem Formulation

- Huge complexity of routing problem
 - Global routing + detail routing
 - Decomposition & inter-gcell routing



Obj → min. congestion and wirelength

Existing Problems

- **Memory overhead**
 - **Cause** → avoiding of overlap
 - *Our solution* → *low-memory routing by graph coloring*

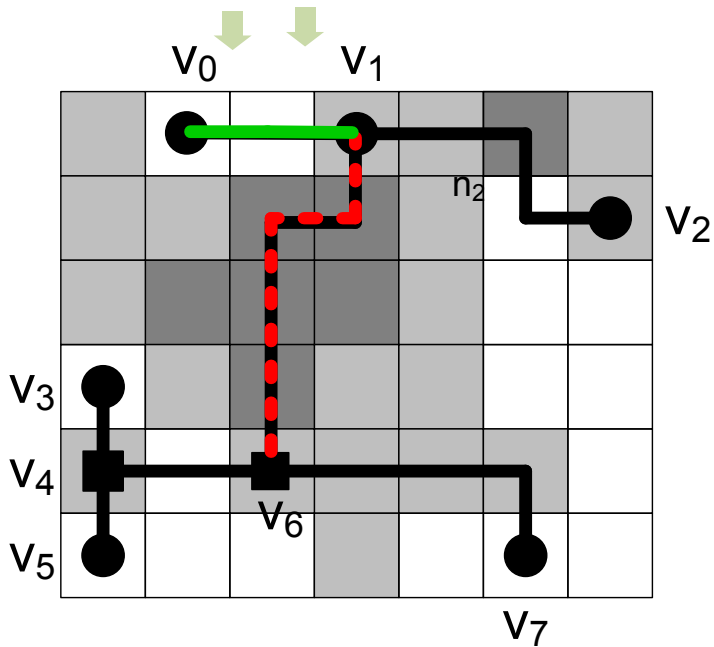
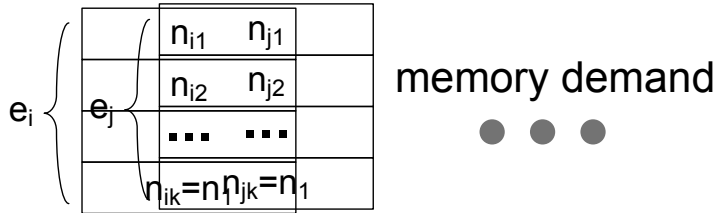
- **Topology corruption**
 - **Cause** → subnet rip-up and rerouting
 - *Our solution* → *dynamic topology update*

- **Suboptimality of Lee's Maze routing method**
 - **Cause** → vertex expansion based method
 - *Our solution* → *bending-aware optimum maze routing*

Memory Overhead

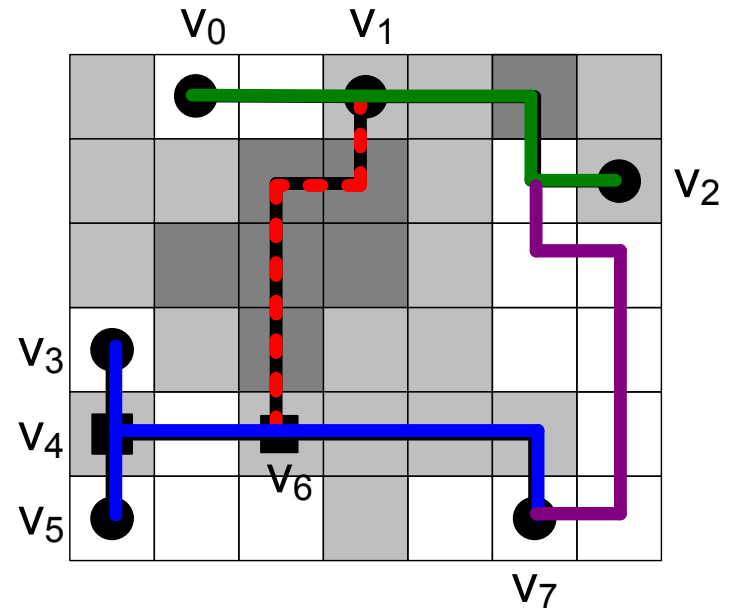
- **Observation from experiment**
 - Routers are consuming huge memory
- **Why low memory?**
 - Cache miss, page fault and thrashing → efficiency
 - Easier for multi-threading
 - Negligible penalty on quality and efficiency
- **Cause of memory demand**
 - Avoid routing overlap
 - Recording list of nets using each edge

Improvement on Memory



Previous method

$O(N)$ Memory $N = \text{size of net}$



Our method

Zero Memory

Improvement on Efficiency

- **Pre-routing time cost**
 - Previous methods $\rightarrow 0$
 - Our method $\rightarrow O(N)$, due to net traversal
 - Negligible due to $O(N^2)$ routing time cost

- **Routing time cost $\rightarrow O(N^2)$**
 - Need to search over all grids around the net
 - Other methods $\rightarrow k \times N^2$ ($k \gg 1$)
 - Hash table operation cost
 - Other data structure will induce higher complexity
 - Our method \rightarrow exactly N^2
 - One operation \rightarrow color checking

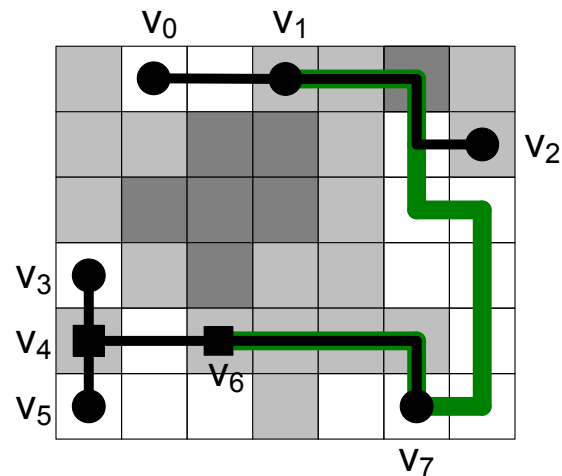
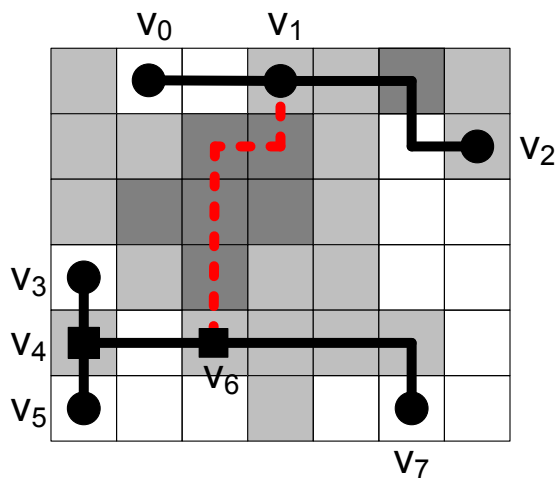
Topology Corruption

- **Steiner tree topology**
 - Widely used in modern routers

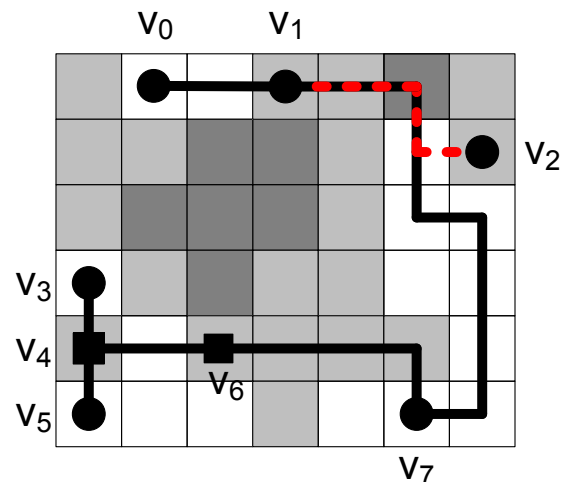
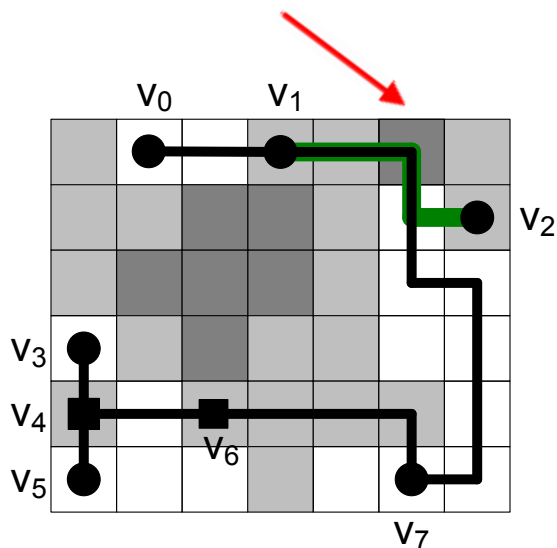
- **Topology corruption**
 - Due to rerouting of congested subnets

- **Previous solutions**
 - BFS to rebuild the topology
 - Usually at the end of iteration
 - Impede congestion reduction

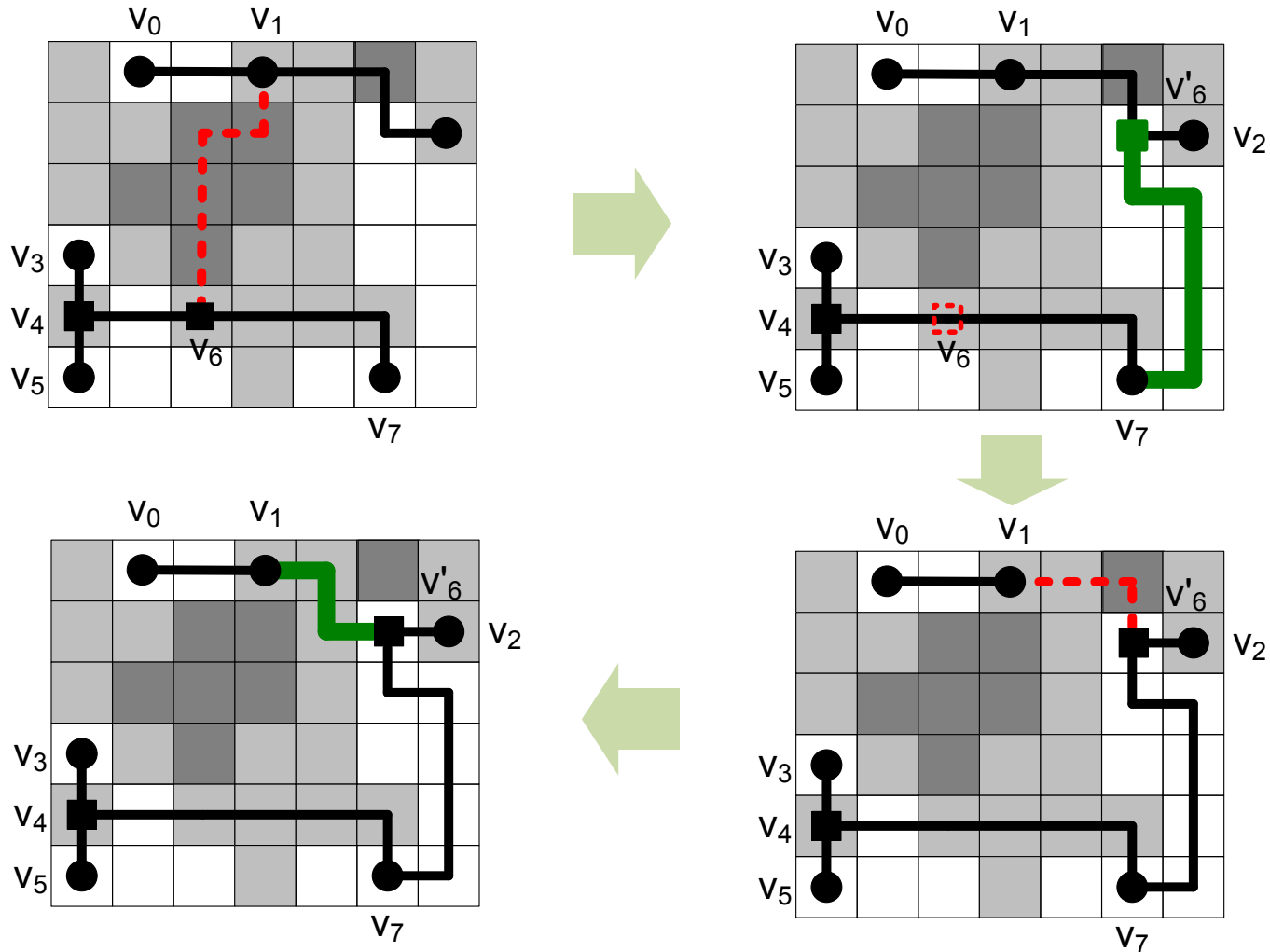
Topology Corruption – Previous Methods



Congestion Problem

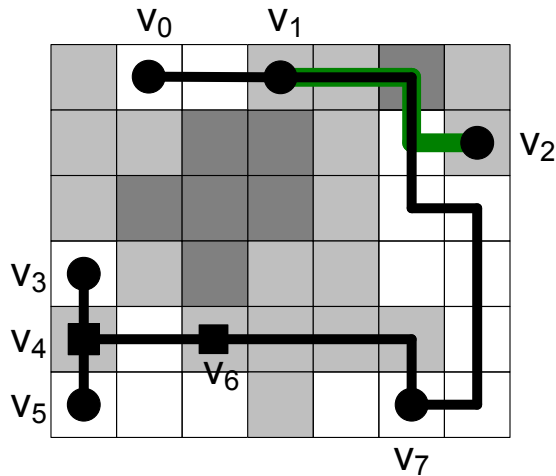


Our Solution – Dynamic Topology Update

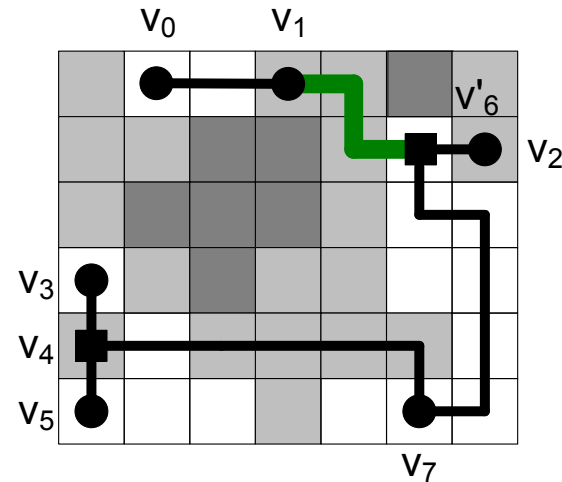


Comparison

Previous Methods



Our Method



Improved routing congestion

$O(N)$ overhead on runtime
(negligible)

Bending-Aware Cost Function

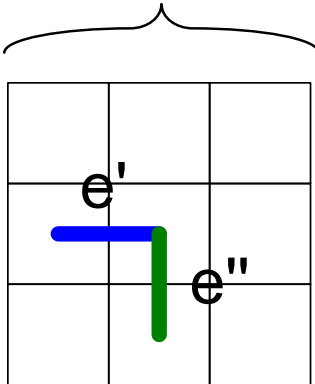
routing cost on $e_{i,j}$

$$c_{i,j} = c_{i,j}^{cong} + c_{i,j}^{len}$$

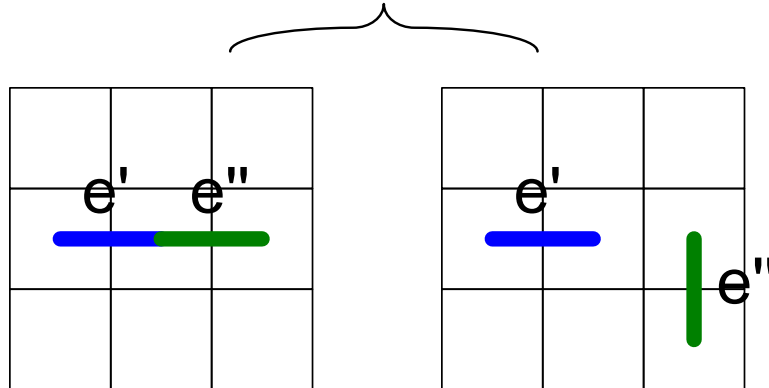
total cost of solution r_i

$$C_i = \sum_{e_{u,v} \in r_i} c_{u,v} + c_{bnd} \left(\sum_{e', e'' \in r_i} IsBend(e', e'') \right)$$

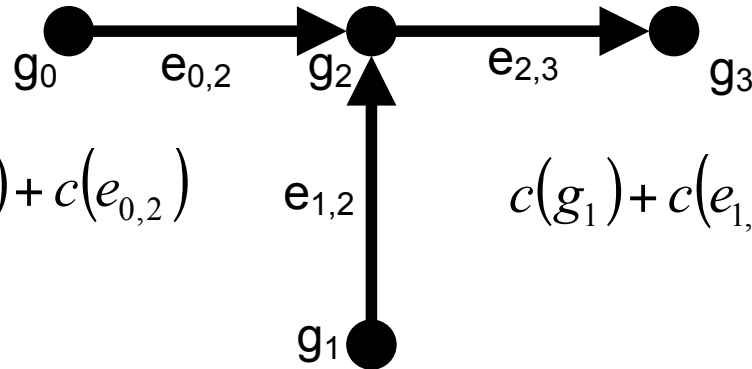
$$IsBend(e', e'') = 1$$



$$IsBend(e', e'') = 0$$



Sub-Optimality of Lee's Algorithm [Lee61]



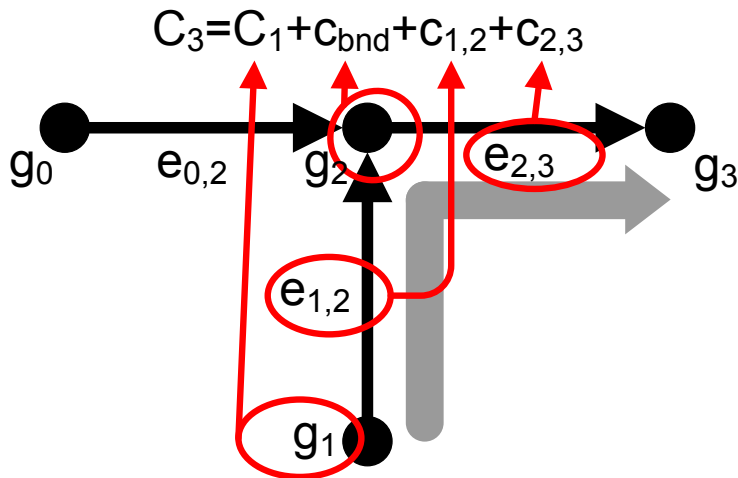
$$c(g_1) + c(e_{1,2}) < c(g_0) + c(e_{0,2})$$

$$c(g_1) + c(e_{1,2}) + c_{bnd} > c(g_0) + c(e_{0,2})$$

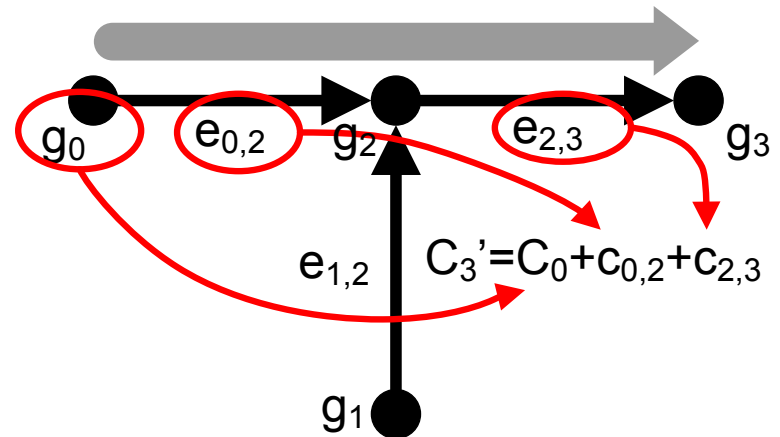


Actual solution for g_3 [Lee61]

Another solution for g_3



Sub-optimum



Optimum

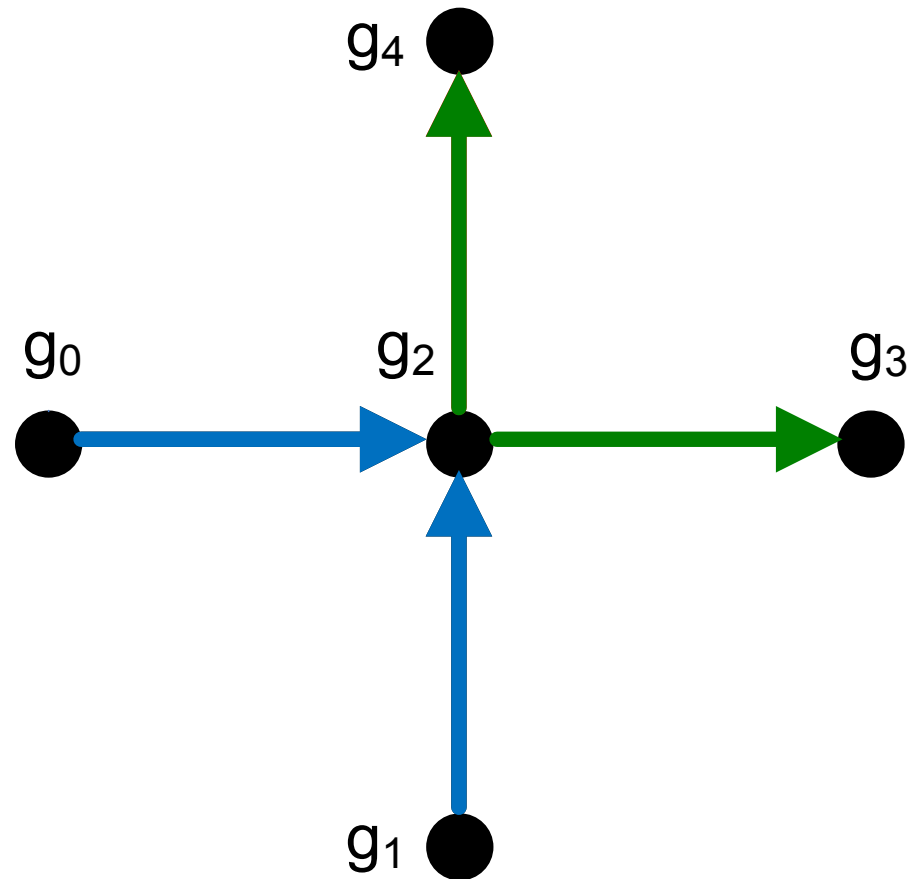
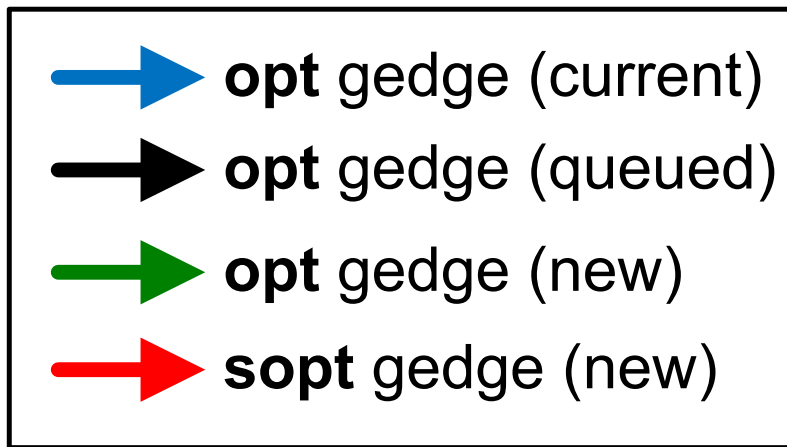
Our Optimum Solution

- **Edge expansion based maze search**
 - Instead of traditional grid expansion
 - All gedges are directed

- **Edges classified into three sets**
 - opt → Optimum solution obtained
 - sopt → SubOptimum solution obtained
 - nopt → No solution obtained

- **Priority queue storing optimized edges**
 - Pop out min-cost edge for expansion

Our Optimum Solution



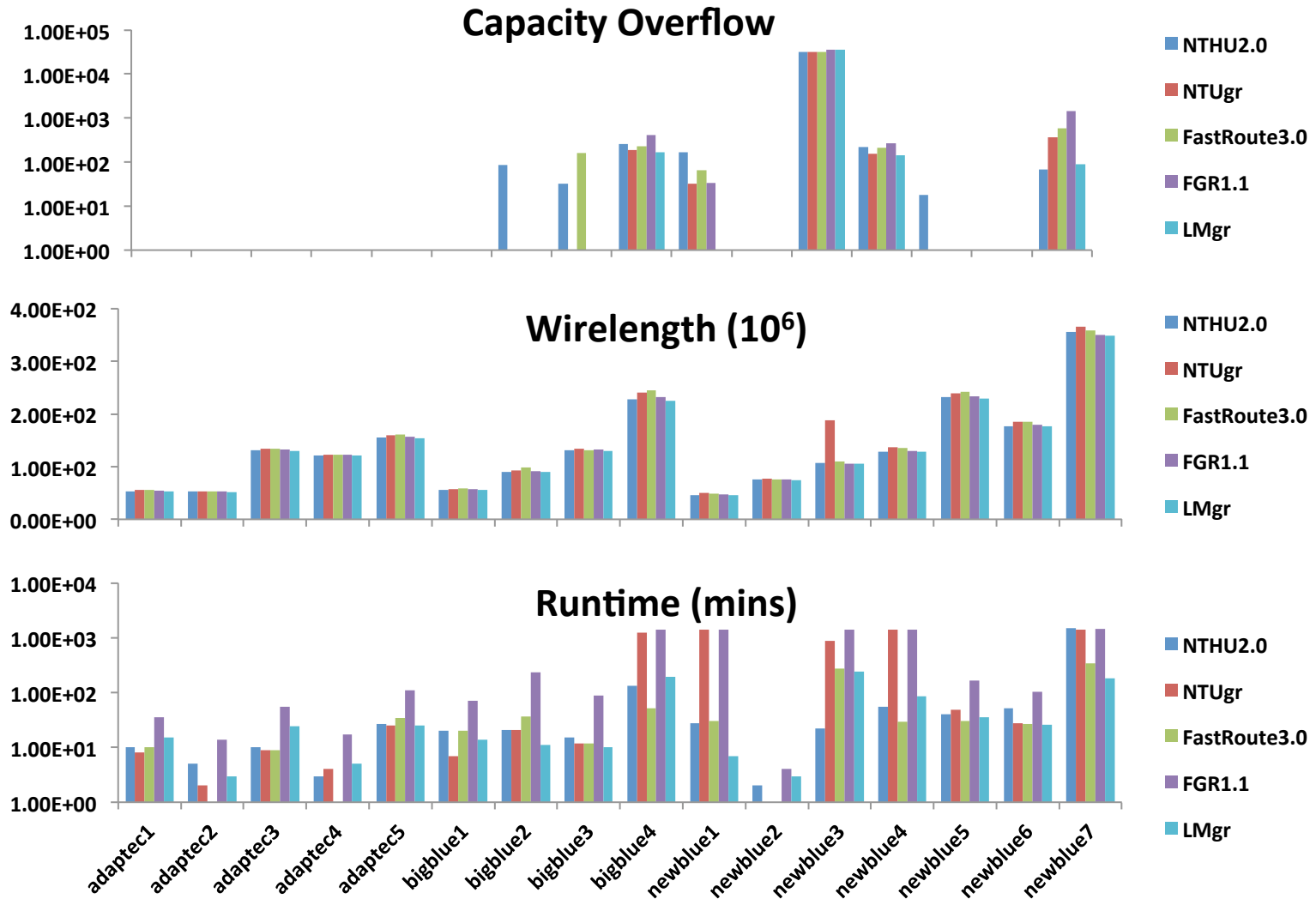
Overview of LMgr

- **2D global routing (LMgr) + 3D layer assign** [Dai09]
- **RSMT topology generation by FLUTE** [Chu08]
- **Pattern-based initial routing**
 - **L-shape** [Kastner02] & **monotonic** [Zhang08]
- **Maze-based iterative rip-up & rerouting**
 - (1) **Congestion- and wirelength-driven**
 - (2) **Congestion-driven (most aggressive)**
 - (3) **Wirelength-driven (zero congestion penalty)**

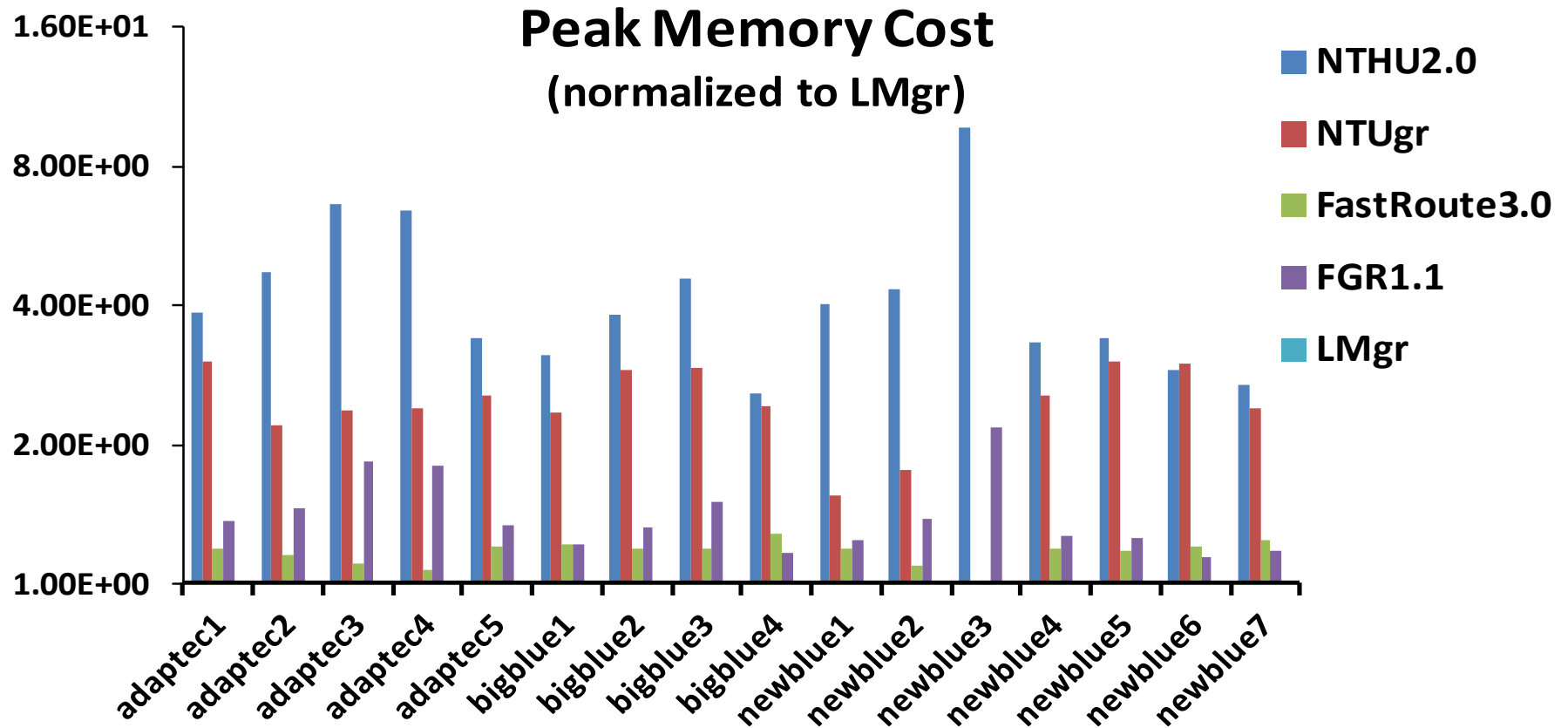
Experiments

- **ISPD08 GR benchmark suite** [ISPD08]
 - Official evaluation script
- **Four sensitivity-driven routers**
 - **NTHU2.0** [Chang08] **NTUgr** [Chen09] **FastRoute3.0** [Zhang08] and **FGR1.1** [Roy07]
 - Winners of ISPD 07-08 contests [ISPD08]
 - Same IRR structure as LMgr
 - Other routers [Xu11, Hu10] are not included due to their incompatibility to our techniques

Routing Solution Quality



Peak Memory Cost



Conclusion

- **Three new techniques**
 - Low-memory global routing
 - Dynamic RSMT update
 - Bending-aware optimum routing
- **Integration in LMgr**
 - Performance validation through experiments
- **Future work**
 - Parallel routing [Liu10]
 - Three-dimensional routing [Chen09]

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Questions?

THANK YOU!