1. (10 points) In Individual HW0, we worked with the following operations: For $L$ a set of strings, we can define the following associated sets

$$L \circ L = \{uw \mid u \in L \text{ and } w \in L\}$$

$$DOUBLE(L) = \{vv \mid v \in L\}$$

$$STUTTER(L) = \{w_1w_2w_3\cdots w_nw_n \mid n \geq 0, w_i \in \Sigma \text{ for each } i \text{ where } 1 \leq i \leq n, \text{ and } w_1w_2\cdots w_n \in L\}$$

Note: $\varepsilon$ is in each of these sets iff it is an element of $L$ itself.

Consider the language $E_1 = \{w \in \{0,1\}^* \mid w \text{ ends in } 1\} = \{u1 \mid u \in \{0,1\}^*\}$. $E_1$ is recognized by each of the machines below.

(The machine on the left may be either a DFA or an NFA because it has exactly one outgoing arrow from each state for each input symbol, and has no spontaneous transitions. The machine on the right is an NFA.)

a. Which of the following DFAs, NFAs, and regular expressions recognizes / describes $E_1 \circ E_1$? List the identifiers of the correct choices; if none do, write “None”.

b. Which of the following DFAs, NFAs, and regular expressions recognizes / describes $DOUBLE(E_1)$? List the identifiers of the correct choices; if none do, write “None”.

c. Which of the following DFAs, NFAs, and regular expressions recognizes / describes $STUTTER(E_1)$? List the identifiers of the correct choices; if none do, write “None”.

d. For each of the four machines listed, label it “DFA” or “NFA” or “either”.

Regexp 1
\[((00)^{*}(11)^{+})^{*}(11)\]

Regexp 2
\((0 \cup 1)^{+}1(0^{*}1)^{+}1\)

Machine 3

Machine 4

Machine 5

Machine 6

| Note: You can download the .jff files for all these machines, open them in JFLAP, and test them on specific inputs and/or use the ‘test equivalence’ feature of JFLAP to check your work. |
| 2. Consider an arbitrary DFA \(M = (Q, \{0, 1\}, \delta, q_0, F)\) and call the language of this DFA \(L\). Fill in the blanks in the definition of a new DFA whose language is the result of taking each string in \(L\) and replacing each 0 in the string with \(a\) and each 1 in the string with \(b\). For example, if \(L = \{0, 001\}\), then the new language is \(\{a, aab\}\). The new machine is \(M' = (Q', \{a, b\}, \delta', q'_0, F')\) where |
| \(Q' = \) This will be the set of states for your new machine. |
| \(\delta'(r, x) = \) For each possible input to the transition function, specify the output. Notice that \(r\) is a state in \(Q'\) and \(x \in \{a, b\}\). |
| \(q'_0 = \) What is the initial state of \(M'\)? Make sure you choose an element of \(Q'\). |
| \(F' = \) What is the set of accepting states of \(M'\)? Choose a subset of \(Q'\). |
3. Consider an arbitrary DFA $M = (Q, \{0, 1\}, \delta, q_0, F)$ and call the language of this DFA $L$. We will define an NFA whose language is the result of taking each string in $L$ and replacing each 0 in the string with $aa$ and each 1 in the string with $ab$. For example, if $L = \{0, 001\}$, then the new language is $\{aa, aaaaab\}$. The idea for this construction is to have two copies of each state to allow us to group the bits read in as pairs. The new machine is

$$M'' = (Q'', \{a, b\}, \delta'', q'', F'')$$

where

$$Q'' = Q \times \{1st, 2nd\}$$

$$\delta''( ((r, 1st), a) ) = \{(r, 2nd)\}$$

$$\delta''( ((r, 1st), b) ) = \emptyset$$

$$\delta''( ((r, 2nd), a) ) = \{\delta((r, 0)), 1st)\}$$

$$\delta''( ((r, 2nd), b) ) = \{\delta((r, 1)), 1st)\}$$

$$q'' = (q_0, 1st)$$

$$F'' = F \times \{1st\}$$

a. Apply this construction to the DFA from question 1:

Draw the resulting state diagram in JFLAP, export the image as a png or jpg file, and include it as part of your submission.

b. Give an example of a string of length 3 for which the computation of $M''$ on this string gets “stuck”.

c. If you were to use Theorem 1.39 on page 55 in the book (the “subset construction”) to write an equivalent DFA to the NFA you produced in a., how many states would the DFA have? Your answer should take into account unreachable states. Extension: (not to submit) is this number optimal?