Trees/Intro to counting

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Thanks to Janine Tiefenbruck

http://cseweb.ucsd.edu/classes/sp16/cse21-bd/
April 29, 2016
**Goal (2):** There is always some way to put directions on the edges of an unrooted tree to make it a rooted tree.

Using the subgoals to achieve the goal:

**Root(\(T\): unrooted tree with \(n\) nodes)**
1. If \(n=1\), let the only vertex \(v\) be the root, set \(h(v):=0\), and return.
2. Find a vertex \(v\) of degree 1 in \(T\), and let \(u\) be its only neighbor.
3. **Root(\(T-\{v\}\)).**
4. Set \(p(v):=u\) and \(h(v):=h(u)+1\).
Example
Counting
1, 2, 3, 4, ...
What do we mean by counting?

How many arrangements or combinations of objects are there of a given form?

How many of these have a certain property?
Why is counting important?

For computer scientists:

- **Hardware**: How many ways are there to arrange components on a chip?
- **Algorithms**: How long is this loop going to take? How many times does it run?
- **Security**: How many passwords are there?
- **Memory**: How many bits of memory should be allocated to store an object?

4 digit pin code? 10,000
In some video games, each player can create a character with custom facial features.

How many distinct characters are possible?
In some video games, each player can create a character with custom facial features. How many distinct characters are possible?

Considering only these 12 hairstyles and 8 hair colors, how many different characters are possible?

A. $8 + 12 = 20$
B. $8 \times 12 = 96$
C. $8^{12} = 68719476736$
D. $12^8 = 429981696$
E. None of the above
Product rule

For any sets, A and B: \(|A \times B| = |A| \cdot |B|\)

In our example:

\(A = \{ \text{hair styles} \} \quad |A| = 12\)

\(B = \{ \text{hair colors} \} \quad |B| = 8\)

\(A \times B = \{ (s, c) : s \text{ is a hair style and } c \text{ is a hair color} \}\)

\(|A \times B| = \text{the number of possible pairs of hair styles & hair colors}\)

\(|A \times B| = \text{the number of different ways to specify a character}\)
For any sets, A and B: \(|A \times B| = |A| \cdot |B|\)

More generally:

Suppose that a procedure can be broken down into a sequence of two tasks. If there are \(n_1\) ways to do the first task and for each of these ways of doing the first task, there are \(n_2\) ways to do the second task, then there are \(n_1n_2\) ways to do the procedure.
For any sets, $A$ and $B$: $|A \times B| = |A| \times |B|$

More generally:

To count the number of pairs of objects:
* Count the number of choices for selecting the first object.
* Count the number of choices for selecting the second object.
* Multiply these two counts.

CAUTION: this will only work if the number of choices for the second object doesn't depend on which first object we choose.
Other than the 96 possible custom Miis, a player can choose one of 10 preset characters.

How many different characters can be chosen?

A. 96
B. 10
C. 106
D. 960
E. None of the above.
For any disjoint sets, A and B: \(|A \cup B| = |A| + |B|\)

In our example:

- \(A = \{ \text{custom characters} \} \) \(|A| = 96\)
- \(B = \{ \text{preset characters} \} \) \(|B| = 10\)

\(A \cup B = \{ m : m \text{ is a character that is either custom or preset} \}\)

\(|A \cup B| = \text{the number of possible characters}\)
For any disjoint sets, A and B: $|A \cup B| = |A| + |B|$

More generally:

If a task can be done either in one of $n_1$ ways or in one of $n_2$ ways, where none of the set of $n_1$ ways is the same as any of the set of $n_2$ ways, then there are $n_1 + n_2$ ways to do the task.
For any disjoint sets, A and B: $|A \cup B| = |A| + |B|$

More generally:

To count the number of objects with a given property:
* Divide the set of objects into mutually exclusive (disjoint/nonoverlapping) groups.
* Count each group separately.
* Add up these counts.
Select which method lets us count the number of length n binary strings.

A. The product rule.
B. The sum rule.
C. Either rule works.
D. Neither rule works.
Select which method lets us count the number of length n binary strings.

A. The product rule. Select first bit, then second, then third …
B. The sum rule. \( \{0\ldots\} \cup \{1\ldots\} \) gives recurrence \( N(n) = 2N(n-1), N(0)=1 \)
C. Either rule works.
D. Neither rule works.
Memory: storing length $n$ binary strings

How many binary strings of length $n$ are there?

How many bits does it take to store a length $n$ binary string?
Memory: storing length $n$ binary strings

How many binary strings of length $n$ are there? $2^n$

How many bits does it take to store a length $n$ binary string? $n$

**General principle:** number of bits to store an object is $\lceil \log_2(\text{number of objects}) \rceil$

Why the ceiling function?
Scenario: We want to store a non-negative integer that has at most n digits. How many bits of memory do we need to allocate?

A. \( n \)
B. \( 2^n \)
C. \( 10^n \)
D. \( n \log_2 10 \)
E. \( n \log_{10} 2 \)

\[
\log_2 (10^n) = \left\lfloor n \log_2 (10) \right\rfloor
\]
At an ice cream parlor, you can choose to have your ice cream in a bowl, cake cone, or sugar cone. There are 20 different flavors available.

How many single-scoop creations are possible?

A. 20
B. 23
C. 60
D. 120
E. None of the above.
At an ice cream parlor, you can choose to have your ice cream in a bowl, cake cone, or sugar cone. There are 20 different flavors available.

You can convert your single-scoop of ice cream to a sundae. Sundaes come with your choice of caramel or hot-fudge. Whipped cream and a cherry are options. How many desserts are possible?

A. $20 \cdot 3 \times 2 \times 2 \times 2$
B. $20 \cdot 3 \times 2 \times 2 \times 2 \times 2$
C. $20 \cdot 3 + 20 \cdot 3 \times 2 \times 2$
D. $20 \cdot 3 + 20 \cdot 3 \times 2 \times 2 \times 2$
E. None of the above.
A scheduling problem

In one request, four jobs arrive to a server: J1, J2, J3, J4.

The server starts each job right away, splitting resources among all active ones.

Different jobs take different amounts of time to finish.

How many possible finishing orders are there?

A. $4^4$
B. $4+4+4+4$
C. $4 \times 4$
D. None of the above.
A scheduling problem

In one request, four jobs arrive to a server: J1, J2, J3, J4.

The server starts each job right away, splitting resources among all active ones.

Different jobs take different amounts of time to finish.

How many possible finishing orders are there?

Product rule analysis

- 4 options for which job finishes first.
- Once pick that job, 3 options for which job finishes second.
- Once pick those two, 2 options for which job finishes third.
- Once pick first three jobs, only 1 remains.

$$(4)(3)(2)(1) = 4! = 24$$

Which options are available will depend on first choice; but the number of options will be the same.
Permutations

Permutation:
rearrangement / ordering of n distinct objects so that each object appears exactly once

Theorem 1: The number of permutations of n objects is

\[ n! = n(n-1)(n-2) \ldots (3)(2)(1) \]

Convention: 0! = 1

Rosen p. 407
Traveling salesperson

Planning a trip to

- New York
- Chicago
- Baltimore
- Los Angeles
- San Diego
- Minneapolis
- Seattle

Must **start in New York** and **end in Seattle**.

How many ways can the trip be arranged?

- A. $7!$
- B. $2^7$
- C. None of the above.
Traveling salesperson

Planning a trip to

- New York
- Chicago
- Baltimore
- Los Angeles
- San Diego
- Minneapolis
- Seattle

Must **start in New York** and **end in Seattle**.
Must also **visit Los Angeles immediately after San Diego**.

*How many ways can the trip be arranged now?*
Planning a trip to

New York
Chicago
Baltimore
Los Angeles
San Diego
Minneapolis
Seattle

Must **start in New York** and **end in Seattle**.
Must also **visit Los Angeles immediately after San Diego**.

*Treat LA & SD as a single stop.*

\[(1)(4!)(1) = 24\] arrangements.

**How many ways can the trip be arranged now?**
Traveling salesperson

Planning a trip to

New York
Chicago
Baltimore
Los Angeles
San Diego
Minneapolis
Seattle

Must **start in New York** and **end in Seattle**.
Must also **visit Los Angeles and San Diego immediately after each other (in any order)**.

*How many ways can the trip be arranged now?*
Traveling salesperson

Planning a trip to

New York
Chicago
Baltimore
Los Angeles
San Diego
Minneapolis
Seattle

Must start in New York and end in Seattle.
Must also visit Los Angeles and San Diego immediately after each other (in any order).

How many ways can the trip be arranged now?

Break into two disjoint cases:
Case 1: LA before SD 24 arrangements
Case 2: SD before LA 24 arrangements
Traveling salesperson

Planning a trip to

New York
Chicago
Baltimore
Los Angeles
San Diego
Minneapolis
Seattle

Must start in New York and end in Seattle.
Must also visit Los Angeles and San Diego immediately after each other (in any order).

How many ways can the trip be arranged now?
Traveling salesperson

Planning a trip to

New York  
Chicago  
Baltimore  
Los Angeles  
San Diego  
Minneapolis  
Seattle

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<th>NY</th>
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<td>1700</td>
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Is there an order of visiting the cities that stops at each city exactly once and minimizes the distance traveled?
Planning a trip to

New York
Chicago
Baltimore
Los Angeles
San Diego
Minneapolis
Seattle

Is there an order of visiting the cities that stops at each city exactly once and minimizes the distance traveled?

Want a Hamiltonian tour
Traveling salesperson

Developing an algorithm which, given a set of cities and distances between them, computes a shortest distance path between all of them is $\text{NP-hard}$ (considered intractable, very hard).

Is there any algorithm for this question?

A. No, it's not possible.
B. Yes, it's just very slow.
C. ?
Traveling salesperson

Exhaustive search algorithm

List all possible orderings of the cities.
For each ordering, compute the distance traveled.
Choose the ordering with minimum distance.

How long does this take?

Want a Hamiltonian tour
Traveling salesperson

Exhaustive search algorithm: given \( n \) cities and distances between them.

List all possible orderings of the cities. For each ordering, compute the distance traveled. Choose the ordering with minimum distance. \( O(\text{number of orderings}) \)

How long does this take?
Traveling salesperson

Exhaustive search algorithm: given $n$ cities and distances between them.

List all possible orderings of the cities. For each ordering, compute the distance traveled. Choose the ordering with minimum distance.

$O(\text{number of orderings})$

How long does this take?

A. $O(n)$
B. $O(n^2)$
C. $O(n^n)$
D. $O(n!)$
E. None of the above.
Traveling salesperson

Exhaustive search algorithm: given $n$ cities and distances between them.

List all possible orderings of the cities.
For each ordering, compute the distance traveled.
Choose the ordering with minimum distance.

How long does this take?

A. $O(n)$
B. $O(n^2)$
C. $O(n^n)$
D. $O(n!)$
E. None of the above.

Moral: counting gives upper bound on algorithm runtime.
A **complete bipartite graph** is an undirected graph whose vertex set is partitioned into two sets $V_1$, $V_2$ such that

- there is an edge between each vertex in $V_1$ and each vertex in $V_2$
- there are no edges both of whose endpoints are in $V_1$
- there are no edges both of whose endpoints are in $V_2$

**Is this graph Hamiltonian?**

A. Yes
B. No
A **complete bipartite graph** is an undirected graph whose vertex set is partitioned into two sets $V_1$, $V_2$ such that

- there is an edge between each vertex in $V_1$ and each vertex in $V_2$
- there are no edges both of whose endpoints are in $V_1$
- there are no edges both of whose endpoints are in $V_2$

Is every complete bipartite graph Hamiltonian?

A. Yes  
B. No
Bipartite Graphs

Claim: any complete bipartite graph with $|V_1| = k$, $|V_2| = k+1$ is Hamiltonian.

How many Hamiltonian tours can we find?

A. $k$
B. $k(k+1)$
C. $k!(k+1)!$
D. $(k+1)!$
E. None of the above.

Rosen p. 658
Bipartite Graphs

Rosen p. 658

Claim: any complete bipartite graph with $|V_1| = k$, $|V_2| = k+1$ is Hamiltonian.

How many Hamiltonian tours can we find?

A. $k$
B. $k(k+1)$
C. $k!(k+1)!$
D. $(k+1)!$
E. None of the above.

Product rule!
When product rule fails

How many Hamiltonian tours can we find?

A. 5!
B. 5!4!
C. ?
When product rule fails

Tree Diagrams

Which Hamiltonian tours start at e?
List all possible next moves.
Then count leaves.

Rosen p.394-395
Let A = \{ people who know Java \} and B = \{ people who know C \}

How many people know Java or C (or both)?

A. \(|A| + |B|\)
B. \(|A| |B|\)
C. \(|A|^{|B|}\)
D. \(|B|^{|A|}\)
E. None of the above.
When sum rule fails

Let $A = \{ \text{people who know Java} \}$ and $B = \{ \text{people who know C} \}$

$\# \text{people who know Java or C} = \# \text{people who know Java}$
When sum rule fails

Let $A = \{ \text{people who know Java} \}$ and $B = \{ \text{people who know C} \}$

$\# \text{people who know Java or C} = \# \text{people who know Java} + \# \text{people who know C}$

Double counted!

Rosen p. 392-394
When sum rule fails

Let $A = \{ \text{people who know Java} \}$ and $B = \{ \text{people who know C} \}$

$\# \text{people who know Java or C} = \# \text{people who know Java}$

$\quad + \# \text{people who know C}$

$\quad - \# \text{people who know both}$

Rosen p. 392-394
Inclusion-Exclusion principle

Let $A = \{ \text{people who know Java} \}$ and $B = \{ \text{people who know C} \}$

$$|A \cup B| = |A| + |B| - |A \cap B|$$

Rosen p. 392-394
Inclusion-Exclusion for three sets

\[ |A \cup B \cup C| = ? \]
Inclusion-Exclusion for three sets

|A ∪ B ∪ C| = ?

Rosen p. 392-394
Inclusion-Exclusion for three sets

\[ |A \cup B \cup C| = ? \]
Inclusion-Exclusion for three sets

\[ |A \cup B \cup C| = ? \]
Inclusion-Exclusion for three sets

\[ |A \cup B \cup C| = ? \]
Inclusion-Exclusion for three sets

|A ∪ B ∪ C| =?

Rosen p. 392-394
Inclusion-Exclusion for three sets

\[ |A \cup B \cup C| =? \]
Inclusion-Exclusion for three sets

\[ |A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |B \cap C| - |A \cap C| + |A \cap B \cap C| \]

Rosen p. 392-394
If $A_1, A_2, \ldots, A_n$ are finite sets then

$$|A_1 \cup A_2 \cup \cdots \cup A_n| = \sum_{1 \leq i \leq n} |A_i| - \sum_{1 \leq i < j \leq n} |A_i \cap A_j| + \sum_{1 \leq i < j < k \leq n} |A_i \cap A_j \cap A_k|$$

$$- \cdots + (-1)^{n+1} |A_1 \cap A_2 \cap \cdots \cap A_n|$$
How many four-letter strings have one vowel and three consonants?

There are 5 vowels: AEIOU
and 21 consonants: BCDFGHJKLMNPQRSTVWXYZ.

A. $5 \times 21^3$
B. $26^4$
C. $5+52$
D. None of the above.
How many four-letter strings have one vowel and three consonants?

There are 5 vowels: AEIOU
and 21 consonants: BCDFGHJKLMNPQRSTVWXYZ.

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<td>21 * 21 * 5 * 21</td>
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</table>

Total: $4 \* 5 \* 21^3$
Counting with categories

If \( A = X_1 \cup X_2 \cup \ldots \cup X_n \) and all \( X_i, X_j \) disjoint and all \( X_i \) have same size, then

\[ |X_i| = \frac{|A|}{n} \]

More generally:

There are \( \frac{n}{d} \) ways to do a task if it can be done using a procedure that can be carried out in \( n \) ways, and for every way \( w \), \( d \) of the \( n \) ways give the same result as \( w \) did.
If $A = X_1 \cup X_2 \cup \ldots \cup X_n$ and all $X_i$, $X_j$ disjoint and all $X_i$ have same size, then

$$|X_i| = |A| / n$$

More generally:

There are $n/d$ ways to do a task if it can be done using a procedure that can be carried out in $n$ ways, and for every way $w$, $d$ of the $n$ ways give the same result as $w$ did.
If $A = X_1 \cup X_2 \cup \ldots \cup X_n$ and all $X_i$, $X_j$ disjoint and all $X_i$ have same size, then

$$|X_i| = |A| / n$$

Or in other words,

If objects are partitioned into categories of equal size, and we want to think of different objects as being the same if they are in the same category, then

$$\text{# categories} = (\text{# objects}) / (\text{size of each category})$$
An ice cream parlor has \( n \) different flavors available. How many ways are there to order a two-scoop ice cream cone (where you specify which scoop goes on bottom and which on top, and the two flavors must be different)?

A. \( n^2 \)
B. \( n! \)
C. \( n(n-1) \)
D. \( 2n \)
E. None of the above.
An ice cream parlor has \( n \) different flavors available. How can we use our earlier answer to decide the number of cones, if we count two cones as the same if they have the same two flavors (even if they're in opposite order)?

A. Double the previous answer.
B. Divide the previous answer by 2.
C. Square the previous answer.
D. Keep the previous answer.
E. None of the above.
An ice cream parlor has $n$ different flavors available. How can we use our earlier answer to decide the number of cones, if we count two cones as the same if they have the same two flavors (even if they're in opposite order)?

**Objects:**

**Categories:**

**Size of each category:**

$$\text{# categories} = \frac{\text{# objects}}{\text{(size of each category)}}$$
Ice cream!

An ice cream parlor has $n$ different flavors available. How can we use our earlier answer to decide the number of cones, if we count two cones as the same if they have the same two flavors (even if they're in opposite order)?

**Objects:** cones  
**Categories:** flavor pairs (regardless of order)  
**Size of each category:**

$$\# \text{ categories} = (\# \text{ objects}) \div (\text{size of each category})$$
An ice cream parlor has $n$ different flavors available. How can we use our earlier answer to decide the number of cones, if we count two cones as the same if they have the same two flavors (even if they're in opposite order)?

**Objects**: cones \( n(n-1) \)

**Categories**: flavor pairs (regardless of order)

**Size of each category**: 2

\[
\text{# categories} = \frac{(n)(n-1)}{2}
\]

Avoiding double-counting
How many different colored triangles can we create by tying these three pipe cleaners end-to-end?

A. $3!$
B. $2^3$
C. $3^2$
D. 1
E. None of the above.
How many different colored triangles can we create by tying these three pipe cleaners end-to-end?

**Objects**: all different colored triangles

**Categories**: physical colored triangles (two triangles are the same if they can be rotated and/or flipped to look alike)

**Size of each category**:

\[
\# \text{ categories} = \frac{\# \text{ objects}}{\text{size of each category}}
\]
How many different colored triangles can we create by tying these three pipe cleaners end-to-end?

**Objects:** all different colored triangles \(3!\)

**Categories:** physical colored triangles (two triangles are the same if they can be rotated and/or flipped to look alike)

**Size of each category:** \((3)(2)\) three possible rotations, two possible flips

\[
\# \text{ categories} = \frac{\# \text{ objects}}{\text{size of each category}} = \frac{6}{6} = 1
\]