Fetch Unit Operation

CSE 141L Lab 2
Hot off the presses!

La Jolla, CA -- Thirty-three dedicated processor design teams have successfully created new design concepts for highly optimized Supergarbage execution engines.

JHacker90, speaking for the group of happy hackers, said "It was a heroic effort, but nothing that couldn't be handled with perseverance and a large pizza." When queried about the rapidly exploding processor design industry in San Diego, executives at Intel corporation answered curtly "We don't comment on our competition, no matter how awesome."

Meanwhile, the US Patent Office has decided to open a new branch in San Diego to handle the imminent wave of patent filings that are sure to ensue.

Upon receipt of the news, Jake Ivring, of UnificationITY, Inc, was seen raising his fist to the sky and proclaiming: "Finally! We will crush that meddlesome Boatorola Corp! To Supergarbage and Beyond!"
Lab 2A Status

• Great progress by everyone
  • Almost halfway to a fully working processor!

• Amnesty policy:
  • By end of today (4/13): 0% penalty
  • By end of tomorrow (4/14): 5% penalty
Front-End vs Back-End

- Front-end *fetches* instructions (Lab 2)
  - Reads from instruction memory
  - Stores fetched instructions in FIFO
  - Predicts which addresses to fetch
- Back-end *executes* instructions (Lab 3)
  - Pulls instructions from front-end’s FIFO
  - Reads/writes to data memory
  - Must check correctness of front-end prediction.

*Why separate fetch from execution?*
Pipelined Fetch Unit

- Fetch unit split into two separate stages
- Address Calculation
  - Determines what address goes into SRAM
- FIFO storage
  - Stores instruction address and SRAM output into FIFO
Fetch Pipeline In Action

Cycle # | 0 | 1 | 2 | 3 | 4 | 5
---|---|---|---|---|---|---
0x0: add $1, $2, $3 | AC | FS |
0x1: sub $3, $3, $4 | AC | FS |
0x2: bez $3, 0x10 | AC | FS |
0x12: add $7, $5, $9 | AC | FS |
0x13: ... | AC | ... |
# Restarting Fetch

<table>
<thead>
<tr>
<th>Cycle #</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>5</th>
<th>6</th>
<th>7</th>
</tr>
</thead>
<tbody>
<tr>
<td>0x2: bez $3, 0x10</td>
<td>AC</td>
<td>FS</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>0x12: add $7, $5, $9</td>
<td>AC</td>
<td>FS</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>0x13: mul $7, $7, $4</td>
<td>AC</td>
<td>FS</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>0x3: add $3, $6, $7</td>
<td>AC</td>
<td>FS</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>0x4: ...</td>
<td>AC</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

The back-end detects a mispredict! Restart @ 0x3

Clear FIFO. Why?
### Stalling for loads/stores

<table>
<thead>
<tr>
<th>Cycle #</th>
<th>0</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>5</th>
</tr>
</thead>
<tbody>
<tr>
<td>0x5: add $1, $2, $3</td>
<td>AC</td>
<td>FS</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>0x6: sub $3, $3, $4</td>
<td></td>
<td></td>
<td>AC</td>
<td>FS</td>
<td></td>
<td></td>
</tr>
<tr>
<td>0x7: muli $3, 0x8</td>
<td></td>
<td></td>
<td></td>
<td>Load addr 0x30 asserted</td>
<td>AC</td>
<td>FS</td>
</tr>
<tr>
<td>0x8: add $7, $5, $9</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>AC</td>
<td>FS</td>
</tr>
</tbody>
</table>
Stalled on Full FIFO

Cycle #  0  1  2  3  4  5

0x5: add $1, $2, $3
   AC  FS

0x6: sub $3, $3, $4
   AC  FS  FS

0x7: mulli $3, 0x8
   AC  AC  FS

0x8: add $7, $5, $9
   AC  FS

FIFO Full asserted
FIFO Full deasserted
Lab 2B: Fetch Unit Control
Control logic implementation

- Fetch unit is pipelined
- Not modeled as a finite state machine
- Each pipeline stage depends on a subset of the control signals
- Need to identify which signals those are
Simplified Implementation Strategy

• For each pipeline stage, ask yourself the following questions:
  • What control signals are needed for the datapath during this stage?
  • What does each of these signals do?
  • What inputs do I need to correctly set those control signals?
  • Encode relationship between control input and output using Verilog
A Simple Example

- Address Calculation stage:
  - What are the control signals associated with that stage?
    - sel_mux[3:0], ...
  - What does sel_mux[0] do?
    - Selects between PC+1 and PC+offset
  - What is the value of sel_mux[0] based on?
    - “P” bit: dictates whether we branch or not
  - Anything else?