

# CSE 20

# DISCRETE MATH

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<https://goo.gl/forms/1o1KOd17RMoURxjn2>

Fall 2017

<http://cseweb.ucsd.edu/classes/fa17/cse20-ab/>

# Today's learning goals

- Explain the steps in a proof by mathematical and/or structural induction
- Use a recursive definition of a set to determine the elements of the set
- Write a recursive definition of a given set
- Use structural induction to prove properties of a recursively defined set
- Define functions from  $\mathbf{N}$  to  $\mathbf{R}$  using recursive definitions

# Mathematical induction

*Rosen p. 311*

To show that some statement  $P(k)$  is true about **all** nonnegative integers  $k$ ,

1. Show that it's true about 0      i.e.  $P(0)$
2. Show  $\forall k ( P(k) \rightarrow P(k + 1) )$  Hence conclude  $P(1), \dots$



# Mathematical induction\*

*Rosen p. 311*

To show that some statement  $P(k)$  is true about **all** nonnegative integers  $k \geq b$ ,

1. Show that it's true about  **$b$**  i.e.  $P(b)$
2. Show  $\forall k ( P(k) \rightarrow P(k + 1) )$  Hence conclude  $P(b+1), \dots$



# Sizes of (finite) sets

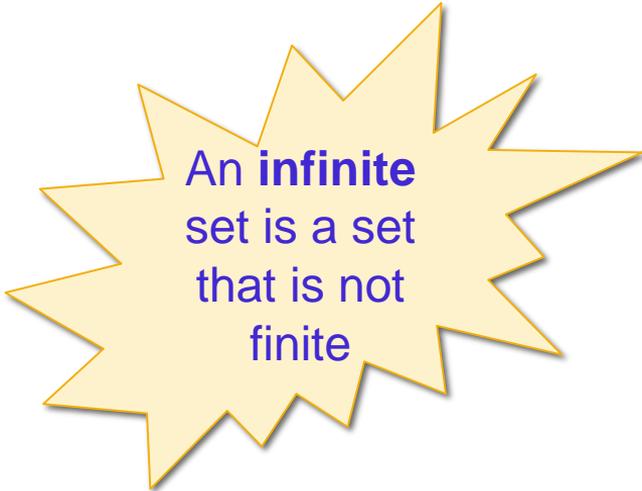
If  $S$  is a set with exactly  $n$  **distinct** elements, with  $n$  a nonnegative integer, then  $S$  is **finite set** and  $|S| = n$ .

$$|\{\}\| = 0 \quad |\{1\}| = 1 \quad |\{1,1,2,3,8/4\}| = 3$$

For  $A, B$  finite sets:

$$|A \cup B| = |A| + |B| - |A \cap B|$$

- $|A \times B| = |A| \cdot |B|$
- if  $A$  has size  $n$ , then the power set of  $A$  has size  $2^n$ .



An **infinite** set is a set that is not finite

# Recursive definitions, part 1

*Rosen Sec 5.3*

To specify a **function**,  
we need to specify its

1. domain – input
2. codomain – type of outputs
3. assignment / rule – formula, table of values, induction

Rosen p. 138

# Recursive definitions, part 1

*Rosen Sec 5.3*

For functions  $f: \mathbf{N} \rightarrow X$

- define by a (closed-form) formula

f(0)      f(1)      f(2)      f(3)      f(4)

1. domain –  $\mathbf{N}$
2. codomain –  $X$
3. assignment / rule – formula, table of values, induction

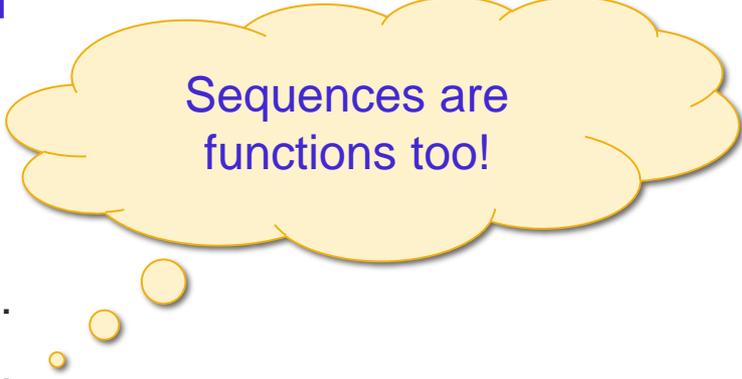
# Recursive definitions, part 1

*Rosen Sec 5.3*

For functions  $f: \mathbf{N} \rightarrow X$

- define by a (closed-form) formula

$f(0)$	$f(1)$	$f(2)$	$f(3)$	$f(4)$	...
$a_0$	$a_1$	$a_2$	$a_3$	$a_4$	...



Sequences are functions too!

1. domain –  $\mathbf{N}$
2. codomain –  $X$
3. assignment / rule – formula, table of values, induction

# BA Recursive definitions, part 1

Rosen Sec 5.3

For functions  $f: \mathbf{N} \rightarrow X$

- define by a (closed-form) formula or ...

*Basis step: Specify the value of the function at 0*

*Recursive step: Give a rule for finding its value at an integer from its values at smaller integers*

Which of the following functions / sequences have recursive definitions?

A.  $n!$

B.  $2^n$

C. 2, -8, 32, -128, 512, ...

D.  $\sum_{i=1}^n (i^2 + i)$

E. All of the above.

# Recursive definitions, part 2

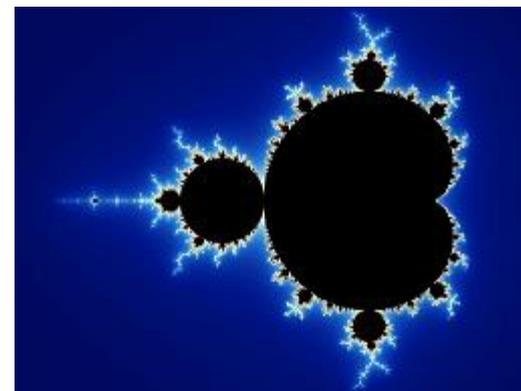
Rosen Sec 5.3

$$n! \quad 0! = 1, \quad n! = n (n-1)! \quad \text{for } n > 0.$$

$$2^n \quad 2^0 = 1, \quad 2^n = 2 (2^{n-1}) \quad \text{for } n > 0.$$

2, -8, 32, -128, 512, ...

$$a_0 = 2, \quad a_n = -4 a_{n-1} \quad \text{for } n > 0.$$



# terms

$$\sum_{i=1}^n g(i) = g(1) + g(2) + \dots + g(n)$$

$\sum_{i=1}^0 g(i) = 0$ ,  $\sum_{i=1}^1 g(i) = g(1)$ , and

Basis

new term

$$\sum_{i=1}^n g(i) = g(n) + \sum_{i=1}^{n-1} g(i) \quad \text{for } n > 1.$$

Rec.

# Extra mathematical induction practice

①

$n^2 < n!$  for all  $n$  greater than or equal to ...

②

$$1 + 2 + 3 + \dots + n = \sum_{i=1}^n i = \frac{n(n+1)}{2}$$

Video

③

$$\sum_{i=1}^n \frac{1}{i(i+1)} = \frac{n}{n+1}$$

# Recursive definitions, part 2

Rosen Sec 5.3

Define a set S by

$$\mathbb{N} = \{0, 1, 2, \dots\}$$

$\{\dots\}$  roster

$\{x \mid P(x)\}$  set builder

## Recursive definition:

**Basis step** – Specify initial collection of elements.



**Recursive step** – Provide rules for forming new elements in the set from those already known to be in the set.

$$\text{If } a \in \mathbb{N} \text{ then } a+1 \in \mathbb{N}$$

Example:  $S = \{ x \mid x \text{ is an int, } x > 3 \}$

Basis Step:

$4 \in S$   
 $17 \in S$

Rec Step:

If  $n \in S$   
then  $n+1 \in S$

# Recursive definitions, part 2

Rosen Sec 5.3

Let  $S$  be the subset of the set of **integers** defined recursively by

*Basis step:*  $1 \in S$

*Recursive step:* If  $a \in S$ , then  $a + 2 \in S$ .

$$S = \{1, 3, 5, \dots\}$$

What's an equivalent description of this set?

- A. All positive multiples of 2.
- B. All positive even integers.
- C. All positive odd integers.
- D. All integers.
- E. None of the above.

$$|\{0,00\}| = 2$$

# Recursive definitions, part 2

Rosen Sec 5.3 page 349

The set of **bit strings**  $\{0,1\}^*$  is defined recursively by

**Basis step:**  $\lambda \in \{0,1\}^*$  where  $\lambda$  is the empty string.

**Recursive step:** If  $w \in \{0,1\}^*$ , then  $w0 \in \{0,1\}^*$  and  $w1 \in \{0,1\}^*$ .

Set of bit strings

$= \{ \lambda, \underline{0}, 1, \underline{00}, 01, 10, 11, \dots, 000 \}$

$w = \lambda$   
 $w0 = 0$

$w = \lambda$   
 $w1 = 1$

$w = 0$   
 $w0 = 00$

Also known as  
**binary sequences**

# Recursive definitions, part 2

Rosen Sec 5.3 page 349

The set of **bit strings**  $\{0,1\}^*$  is defined recursively by

**Basis step:**  $\lambda \in \{0,1\}^*$  where  $\lambda$  is the empty string.

**Recursive step:** If  $w \in \{0,1\}^*$ , then  $w0 \in \{0,1\}^*$  and  $w1 \in \{0,1\}^*$ .

Which of the following are **not** bit strings?

~~A.~~  $\lambda$

B. 1

C. 000

D.  $(01010)1 = ((\dots)0)1 = \dots = (\lambda \dots)$

E. 101100111000111100001111100000111...

Also known as  
**binary sequences**

infinite sequence!

# Recursive definitions in CS

- Strings – encode integers, everything else
- Data structures – linked lists, trees, graphs

*How do we prove a fact is true about **all** strings?*

# Structural induction

*Rosen p. 354*

To show that some statement  $P(k)$  is true about **all** elements of a recursively defined set  $S$ ,

1. Show that it's true for each element specified in the **basis step** to be part of  $S$ .
2. Show that if it's true for each of the elements used to construct **new elements in recursive step**, then it holds for these new elements.

# Structural induction, example

Rosen Sec 5.3

Define the **subset  $S$**  of the set of all bit strings,  $\{0,1\}^*$ , by

*Basis step:*  $\lambda \in S$  where  $\lambda$  is the empty string.

*Recursive step:* If  $w \in S$ , then each of  $\underline{10}w \in S, \underline{01}w \in S$

$\{\lambda, \underbrace{10, 01}_{\substack{1 \text{ appl'n} \\ \text{rec. step}}}, 1010, 0110, \\ 1001, 0101, \dots\}$

# Structural induction, example

Rosen Sec 5.3

Define the **subset S** of the set of all bit strings,  $\{0,1\}^*$ , by

*Basis step:*  $\lambda \in S$  where  $\lambda$  is the empty string.

*Recursive step:* If  $w \in S$ , then each of  $10w \in S, 01w \in S$

**Claim:** Every element in S has an equal number of 0s and 1s.

Pf by Structural Induction

Basis Step WTS  $\lambda$  has equal # of 0s, 1s.  
Note  $\lambda$  has 0 0s, 0 1s, so equal # 

# Structural induction, example

Rosen Sec 5.3

Define the **subset S** of the set of all bit strings,  $\{0,1\}^*$ , by

*Basis step:*  $\lambda \in S$  where  $\lambda$  is the empty string.

*Recursive step:* If  $w \in S$ , then each of  $10w \in S, 01w \in S$

**Claim:** Every element in S has an equal number of 0s and 1s.

**Proof:** *Basis step* – WTS that empty string has equal # of 0s and 1s

*Recursive step* – Let  $w$  be an arbitrary element of S.

Assume, as the **IH** that  $w$  has equal # of 0s and 1s.

WTS that  $10w, 01w$  each have equal # of 0s, 1s.

ex.

# Recursive definitions, part 3

*Rosen Sec 5.3*

Define a function  $f: \mathbf{S} \rightarrow X$  by

- specifying  $f(a)$  when  $a$  is in the basis collection of elements of  $S$
- giving a formula for computing  $f(x)$  when  $x$  is built from other elements  $a_1, a_2$  etc. based on values of  $f(a_1), f(a_2),$  etc.

1. domain –  $\mathbf{S}$ , a recursively defined set
2. codomain –  $X$
3. assignment / rule – recursively defined

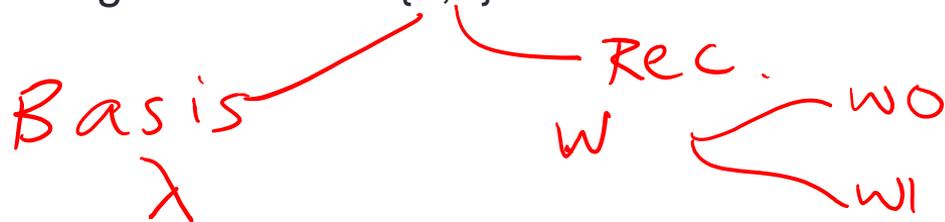
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Define a function  $f: \mathbf{S} \rightarrow X$  by

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- giving a formula for computing  $f(x)$  when  $x$  is built from other elements  $a_1, a_2$  etc. based on values of  $f(a_1)$ ,  $f(a_2)$ , etc.

Example: how do you define the length function  $l: \{0,1\}^* \rightarrow \mathbb{N}$  ?



# Reminders

- Office hours may be shifted – check Google calendar
- HW5 due Saturday night

**Next class:** more applications + variants of induction