Pipelining and Exceptions

- Exceptions represent another form of control dependence.
- Therefore, they create a potential branch hazard.
- Exceptions must be recognized early enough in the pipeline that subsequent instructions can be flushed before they change any permanent state.
- As long as we do that, everything else works the same as before.
- Exception-handling that always correctly identifies the offending instruction is called precise interrupts.

Superscalar Execution

- Not fundamentally different than the techniques we discussed.
- Deeper pipelines.
- Pipelining is combined with
  - superscalar execution
  - out-of-order execution
  - VLIW (very-long-instruction-word)
A modest superscalar MIPS

- what can this machine do in parallel?
- what other logic is required?

Superscalar Execution

- To execute four instructions in the same cycle, we must find four independent instructions.
- If the four instructions fetched are guaranteed by the compiler to be independent, this is a *VLIW* machine.
- If the four instructions fetched are only executed together if hardware confirms that they are independent, this is an *in-order superscalar* processor.
- If the hardware actively finds four (not necessarily consecutive) instructions that are independent, this is an *out-of-order superscalar* processor.
- What do you think are the tradeoffs?

Superscalar Scheduling

- assume in-order, 2-issue, ld-store followed by integer
  lw $6, 36($2)
  add $5, $6, $4
  lw $7, 1000($5)
  sub $9, $12, $5
- assume 4-issue, any combination (VLIW?)
  lw $6, 36($2)
  add $5, $6, $4
  lw $7, 1000($5)
  sub $9, $12, $5
  sw $5, 200($6)
  add $3, $9, $9
  and $11, $7, $6
- When does each instruction begin execution?
Dynamic Scheduling
or
Out-of-Order Scheduling

- Issues (begins execution of) an instruction as soon as all of its dependences are satisfied, even if prior instructions are stalled.

```plaintext
lw $6, 36($2)
add $5, $6, $4
lw $7, 1000($5)
sub $9, $12, $8
sw $5, 200($6)
add $3, $9, $9
and $11, $5, $6
```

Reservation Stations

- are a mechanism to allow dynamic scheduling

![Diagram of Reservation Stations]

PowerPC 604, Pentium Pro (II, III)

- Deep pipeline
- Dynamically Scheduled (out-of-order scheduling)
- Trace Cache
- Simultaneous Multithreading (HyperThreading)

Pentium 4

- Deep pipeline
- Dynamically Scheduled (out-of-order scheduling)
- Trace Cache
- Simultaneous Multithreading (HyperThreading)
Pentium 4

Modern Processors

- Pentium II, III – 3-wide superscalar, out-of-order, 14 integer pipeline stages
- Pentium 4 – 3-wide superscalar, out-of-order, simultaneous multithreading, 20+ pipe stages
- AMD Athlon, 3-wide ss, out-of-order, 10 integer pipe stages
- Alpha 21164 – 2-wide ss, in-order, 7 pipe stages
- Alpha 21264 – 4-wide ss, out-of-order, 7 pipe stages
- Intel Itanium – 3-operation VLIW, 2-instruction issue, in-order, 10-stage pipeline

Pipelining -- Key Points

- ET = Number of instructions * CPI * cycle time
- Data hazards and branch hazards prevent CPI from reaching 1.0, but forwarding and branch prediction get it pretty close.
- Data hazards and branch hazards need to be detected by hardware.
- Pipeline control uses combinational logic. All data and control signals move together through the pipeline.
- Pipelining attempts to get CPI close to 1. To improve performance we must reduce CT (superpipelining) or CPI below one (superscalar, VLIW).